# CS:5810 Formal Methods in Software Engineering 

## Case Study: Hotel Lock System

## Acknowledgments

These notes are based on an Alloy example in the book:
[Jack06] Daniel Jackson. Software abstractions - Logic, Language, and Analysis. The MIT press, 2006

## The Task

- Model in Alloy the disposable card key system used in most hotels for locking and unlocking guest rooms
- The system uses recordable locks, which prevent previous guests from entering a room once its has been re-assigned
- We will model both static and dynamic aspects of the system


## Problem Description [Jack06]

> "[...] the hotel issues a new key to the next occupant, which recodes the lock, so that previous keys will no longer work.

The lock is a simple, stand-alone unit [...] with a memory holding the current key combination.

A hardware device [...] [within the lock] generates a sequence of pseudorandom numbers."

## Problem Description [Jack06]

"The lock is opened either by the current key combination, or by its successor;

If a key with the successor is inserted, the successor is made to be the current combination, so that the old combination will no longer be accepted.

This scheme requires no communication between the front desk and the door lock."

## Problem Description [Jack06]

"By synchronizing the front desk and the door locks initially, and by using the same pseudorandom generator,
the front desk can keep its records of the current combinations in step with the doors themselves."

## Signatures and Fields

## Signatures: Key, Room, Guest, FrontDesk

- Key refers to the key combination stored in the magnetic strip of the card
- FrontDesk stores at any time a mapping
- between each room and its most recent key combination (if any), and
- between each room and its current guest


## Signatures and Fields

Room refers to the room lock
Each room (lock) has

- an associated set of possible keys, and
- exactly one current key at a time

Each key belongs to at most one room
Each guest has zero or more keys at any time

## Signatures and Fields

module hotel
open util/ordering [Key] as KO

## Signatures and Fields

```
module hotel
open util/ordering [Key] as KO
sig Key {}
sig Room {
    keys: set Key,
    var currentKey: Key
}
sig Guest {
    var keys: set Key
}
one sig FrontDesk {
    var lastKey: Room -> lone Key,
    var occupant: Room -> Guest
}
```

```
fun sig FDlastKey : Room -> Key {
    FrontDesk.lastKey
}
fun sig FDoccupant: Room -> Guest {
    FrontDesk.occupant
}
```


## Room Constraint

## Each key belongs to at most one room <br> fact \{ <br> all k: Key | lone (Room <: keys).k <br> \}

```
module hotel
open util/ordering [Key] as KO
sig Key {}
sig Room {
    keys: set Key,
    var currentKey: Key
}
sig Guest { var keys: set Key }
one sig FrontDesk {
    var lastKey: Room -> lone Key,
    var occupant: Room -> Guest
}
fun sig FDlastKey : Room -> Key {
    FrontDesk.lastKey
}
fun sig FDoccupant: Room -> Guest
    FrontDesk.occupant
}
```


## New Key Generation

Given a key k and a set ks of keys,
nextKey returns the smallest key (in the key ordering) in ks that follows k
fun nextKey [k: Key, ks: set Key]: set Key
\{
KO/min [KO/nexts[k] \& ks]
\}

## Initial State

```
module examples/hotel
open util/ordering [Key] as KO
sig Key {}
sig Room {
    keys: set Key,
    var currentKey: Key
}
sig Guest {
}
                        ) No constraints
the record of each room's key
at the front desk is synchronized
with the current combination
of the lock itself
one sig FrontDesk {
    var lastKey: Room -> lone Key,
    var occupant: Room -> Guest
}
```


## Hotel Operations: Initial State

```
pred init {
    -- no guests have keys
    no Guest.keys
    -- the roster at the front desk shows
    -- no room as occupied
    no FDoccupant
    -- the record of each room's key at the
    -- front desk is synchronized with the
    -- current combination of the lock itself
    all r: Room
    r.FDlastKey = r.currentKey
}
```


## Hotel Operations: Guest Entry

```
pred entry [ g: Guest, r: Room, k: Key ]
```

- Preconditions:
- The key used to open the lock is one of the keys the guest is holding
- Pre and Post Conditions:
- The key on the card
- either matches the lock's current key, and the lock remains unchanged (not a new guest), or
- matches its successor, and the lock is advanced (new guest)
- Frame conditions:
- no changes to the state of other rooms, to the set of keys held by guests, or to the records at the front desk


## Hotel Operations: Guest Entry

```
pred entry [ g: Guest, r: Room, k: Key ] {
    -- preconditions
    -- the key used to open the lock is one of the keys
    -- held by the guest
    k in g.keys
    -- pre and post conditions
    (-- not a new guest
    k = r.currKey and r.currKey' = r.currKey
    or
    -- new guest
    k = nextKey[r.currKey, r.keys] and r.currKey' = k
    )
    -- frame conditions
    noFrontDeskChange
    noRoomChange[Room - r]
    noGuestChange[Guest]
}
```


## Hotel Operations: Check-out

pred checkOut [ g: Guest ]

- Preconditions:
- the guest occupies one or more rooms
- Postconditions:
- the guest's rooms become available
- Frame conditions:
- Nothing changes but the occupant relation


## Hotel Operations: Check-out

```
pred checkOut [g: Guest] {
    -- preconditions
    -- the guest occupies one or more rooms
    some FDoccupant.g
    -- postconditions
    -- the guest's rooms become available
    FDoccupant' = FDoccupant - (Room -> g)
    -- frame conditions
    FDlastKey' = FDlastKey
    noRoomChange[Room]
    noGuestChange[Guest]
}
```


## Hotel Operations: Check-in

## pred checkIn [ g: Guest, r: Room, k: Key ]

- Preconditions:
- the room is available
- the input key is the successor of the last key in the sequence associated to the room
- Postconditions:
- the guest holds the input key and becomes the new occupant of the room
- the input key becomes the room's current key
- Frame conditions:
- Nothing changes but the occupant relation and the guest's relations


## Hotel Operations: Check-in

```
pred checkIn [ g: Guest, r: Room, k: Key ] {
    preconditions
    -- the room has no current occupant
    no r.FDoccupant
    -- the input key is the successor of the last key in
    -- the sequence associated to the room
    k = nextKey[r.FDlastKey, r.keys]
    postconditions
    -- the guest becomes the new occupant of the room
    FDoccupant' = FDoccupant + (r -> g)
    -- the guest holds the input key
    g.keys' = g.keys + k
    -- the input key becomes the room's current key
    FDlastKey' = FDlastKey ++ (r -> k)
-- frame conditions
    noRoomChange[Room]
    noGuestChange[Guest - g] }
```


## Trace Generation

- The first time step satisfies the initialization conditions
- Any pair of consecutive time steps are related by
- an entry operation, or
- a check-in operation, or
- a check-out operation


## Trace Generation

```
fact Traces {
    init
    always
    some g: Guest, r: Room, k: Key |
        entry[g, r, k] or
        checkin[g, r, k] or
        checkout[g]
}
```


## Analysis

Let's check if unauthorized entries are possible:

- If a guest $g$ enters room $r$ at time $t$, and the front desk records show $r$ as occupied at that time, then $g$ must be a recorded occupant of $r$.

```
assert noBadEntry {
    always all r: Room, g: Guest, k: Key |
        let o = r.FDoccupant |
            (entry[g, r, k] and some o) implies g in o
```

\}

## Analysis

## check noBadEntry for 3

but 1 Room, 2 Guest, 5 steps

- It is enough to check for problem already with just 2 guests and 2 rooms
- steps's scope must be at least 5 because at least 4 steps are needed to execute each operation once
- There is a counterexample (see file dynamic/hote11-e1ec.a1s)


## 0: Initial State



- initially, the current key of Room is Key0
- this is also reflected in the front desk's record


## 1: CheckIn Operation



- Guest1 checks in to Room and receives key Key1
- the occupancy roster at the front desk is updated accordingly
- Key1 is recorded as the last key assigned to Room


## 2: CheckOut Operation



- Guest1 checks out, and the occupancy roster is cleared
- Since Guest1 never entered, Room's current key is still Key0
- Guest1 still holds Key1


## 3: CheckIn Operation



- Guest0 checks into Room and receives key Key2
- the occupancy roster at the front desk is updated accordingly
- Key2 is recorded as the last key assigned to Room


## 4: Entry Operation



- Guest1 presents Key1 to the lock of Room
- Since Room's current key is still Key0, Guest1 is admitted


## Necessary Restriction

There must be no intervening operation between a guest's check-in and room entry

```
pred noInterveningOps {
    always
        all g: Guest, r: Room, k: Key |
            checkIn[g, r, k] implies after entry[g, r, k]
}
```


## Conditional Assertion

Make assertion under noInterveningOps assumption
assert noBadEntry \{

## noInterveningops implies

```
always al1 r: Room, g: Guest, k: Key |
    let o = r.FrontDesk.occupant |
    (entry[g, r, k] and some o) implies g in o
```

\}

## Analysis

- We check once again:
check noBadEntry for 3
but 2 Room, 2 Guest, 5 steps
- No counter-example (see file dynamic/hote12-e1ec.als)
- For greater confidence, we increase the scope:
check noBadEntry for 5
but 3 Room, 3 Guest, 20 steps
- No counterexamples

