

# **Program Verification**

## **Automated Test Case Generation, Part I**

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# Introduction

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- ✓ ensuring test case coverage methodically

View JML-annotated code as formal description of all anticipated runs

## ATCG Principle

- ▶ Specialize contract/code to representative selection of concrete runs
- ▶ Turn these program runs into executable test cases

# Ideas Behind Automated Test Generation

## Ideas common to systematic (automated) test generation

- ▶ **Formal** analysis of specification and/or code yields enough information to produce test cases
- ▶ Systematic algorithms give certain **coverage** guarantees
- ▶ Post conditions can be turned readily into test **oracles**
- ▶ **Mechanic reasoning** technologies achieve automation: constraint solving, deduction, symbolic execution, model finding

# Automated Test Generation Framework: Unit Tests

Test a single method or function, the **implementation under test** (IUT)

Create test case for popular JAVA unit test framework: JUNIT

## Test Cases in Unit Testing

- ▶ Initialisation of test data (**test fixture/preamble**):  
create program state from which IUT is started
- ▶ Invoke IUT
- ▶ Inspection of result: **test oracle**:  
tell whether test succeeded: PASS or FAIL

# Black box vs White Box Testing

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## Specific Pros and Cons

- ✓ White box testing can use additional information from code
- ✗ White box testing does require source code
- ✓ Black box testing does not require source code
- ✗ Black box testing can be irrelevant/insufficient for IUT

# Program States and JML Expressions

## Reminder

A given **program state**  $S$  makes a boolean JML expression **true** or **false**

## Example

Assume that `int[] arr` has value  $\{1,2\}$  in  $S$

Then “`arr.length==2 && search(arr, 1)==0`” is true in  $S$

# Program States and Test Cases

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```
int [] arr = {1,2};
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## Example

Assume that `int [] arr` has value `{1,2}` in  $S$

This state can be reached by the following preamble:

```
int [] arr = {1,2};
```

Assume we can compute such initialization code automatically

# Specification-Based Test Generation

## Generate test cases from analysing formal specification or formal model of IUT

- ▶ **Black box** technology with according pros and cons
- ▶ Many tools, commercial as well as academic:  
JMLUnit, BZ-TT, JML-TT, UniTesK, JTest,  
TestEra, Cow\_Suite, UTJML, ...
- ▶ Various specification languages: B, Z, Statecharts, JML, ...
- ▶ **Detailed formal specification/system model required**

# Test Generation Principle

View JML contract as formal description of all anticipated runs

## Specification-Based Test Generation Principle

- ▶ Specialize JML contract to representative selection of concrete runs
- ▶ Turn these program runs into executable test cases

# Contracts and Test Cases

```
/*@ public normal_behavior
   @ requires Pre;
   @ ensures  Post;
   @*/
public void m() { ... };
```

All prerequisites for intended behavior contained in **requires** clause

Unless doing robustness testing, consider unintended behavior **irrelevant**

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/*@ public normal_behavior
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All prerequisites for intended behavior contained in **requires** clause

Unless doing robustness testing, consider unintended behavior **irrelevant**

## Test Generation Principle 1

Test data must make required precondition true

# Multi-Part Contracts and Test Cases

```
/*@ public normal_behavior
  @ requires Pre1;
  @ ensures Post1;
  @ also
  @ ...
  @ also
  @ public normal_behavior
  @ requires Pren;
  @ ensures Postn;
  @*/
public void m() { ... };
```

## Test Generation Principle 2

There must be at least one test case for each operation contract

## Example

```
public class Traffic {
    private /*@ spec_public @*/ boolean red, green, yellow;
    private /*@ spec_public @*/ boolean drive, brake, halt;

    /*@ public normal_behavior
       @ requires red || yellow || green;
       @ ensures  \old(red) ==> halt      &&
       @          \old(yellow) ==> brake;
       @*/
    public boolean setAction() {
        // implementation
    }
}
```

Which test cases should be generated?

# Data-Driven Test Case Generation

**Generate a test case for each possible value of each input variable**

- ✗ Combinatorial explosion (already  $2^5$  cases for our simple example)
- ✗ Infinitely many test cases for unbounded data structures
- ✗ Test cases unrelated to specification or IUT

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Restriction to test cases that satisfy precondition?

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Restriction to test cases that satisfy precondition?

Insufficient (still too many), but gives the right clue!

# Disjunctive Partitioning

```
/*@ public normal_behavior
   @ requires red || yellow || green;
   @ ensures  \old(red) ==> halt      &&
   @          \old(yellow) ==> brake;
   @*/
```

Disjunctive analysis suggests at least three test cases related to precondition

# Disjunctive Normal Form

## Disjunctive Normal Form (DNF)

Assume the requires clause has the form

$$C_1 \vee C_2 \vee \dots \vee C_n$$

where each  $C_i$  does not contain an explicit or implicit disjunction.

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For each disjunct of precondition in DNF create test case making it true

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### Example

**requires** red || yellow || green;

Gives rise to three test cases **red=true; yellow=green=false;**, etc.

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## Test Generation Principle 3

For each disjunct of precondition in DNF create test case making it true

## Importance of Establishing DNF

Implicit disjunctions must be made explicit by computing DNF:

Replace  $A \implies B$  with  $\neg A \ || \ B$ , etc.

# Test Coverage Criteria

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requires red || yellow || green;
```

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## (Specification-based) Test Coverage Criterion

How many different test cases to create that make precondition true?

- ▶ At least one (Decision Coverage)
- ▶ ...
- ▶ All (Multiple Condition Coverage)

# Consistent Test Cases

## Example (Class invariant specified in JML)

```
public class Traffic {
    /*@ public invariant (red ==> !green && !yellow) &&
       @                (yellow ==> !green && !red) &&
       @                (green ==> !yellow && !red);
    @*/
    private /*@ spec_public @*/ boolean red, green, yellow;

    /*@ public normal_behavior
       @ requires red || yellow || green;
       @ ...
    */
}
```

The program state `red=yellow=green=true;` violates the class invariant

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```

The program state `red=yellow=green=true;` violates the class invariant

If the class invariant always holds when a method is called, there is no point to generate test cases from program states violating it

# Consistent Test Cases

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public class Traffic {
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    @*/
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```

The program state `red=yellow=green=true;` violates the class invariant

## Test Generation Principle 4

Generate test cases from states that do not violate the class invariant

# Dealing with Large Datatypes (First-Order Logic)

## Example (Square root)

```
/*@ public normal_behavior
   @ requires (\exists int r; r >= 0 && r*r == n);
   @ ensures \result * \result == n;
   @*/
public static final int sqrt(int n) { ... }
```

# Dealing with Large Datatypes (First-Order Logic)

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Where is the disjunction?

# Dealing with Large Datatypes (First-Order Logic)

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## Existential quantifier as disjunction

- ▶ Existentially quantified expression ( $\exists \text{ int } r; P(r)$ )
- ▶ Rewrite as:  $P(\text{MIN\_VALUE}) \vee \dots \vee P(0) \vee \dots \vee P(\text{MAX\_VALUE})$
- ▶ Get rid of those  $P(i)$  that are false:  $P(0) \vee \dots \vee P(\text{MAX\_VALUE})$

# Equivalence Classes on Input Domains

## Example (Square root)

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public static final int sqrt(int n) { ... }
```

**Too many test cases from existential quantifier!**

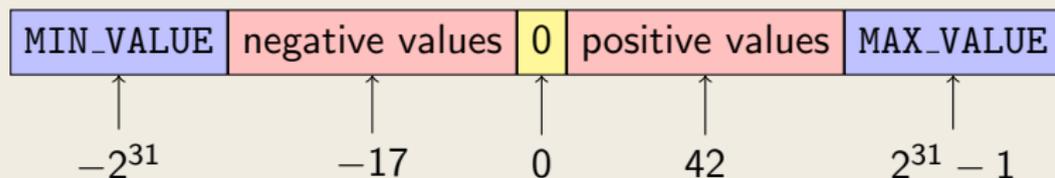
$n = 0*0;$ ,  $n = 1*1;$ , ...,  $n = \text{MAX\_VALUE}*\text{MAX\_VALUE};$

# Equivalence Classes on Input Domains

## Example (Square root)

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/*@ public normal_behavior
   @ requires (\exists int r; r >= 0 && r*r == n);
   @ ensures \result * \result == n;
   @*/
public static final int sqrt(int n) { ... }
```

Partition large/infinite domains in finitely many equivalence classes



... and create test case for only one representative of each

## Example (Square root)

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/*@ public normal_behavior
   @ requires (\exists int r; r >= 0 && r*r == n);
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```

Choice of `r=MAX_VALUE` exhibits defective spec for overflow

# Boundary Values

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Choice of `r=MAX_VALUE` exhibits defective spec for overflow

## Test Generation Principle 5

Include boundary values of ordered domains as class representatives

# Boundary Values

## Example (Square root)

```
/*@ public normal_behavior
   @ requires (\exists int r; r >= 0 && r*r == n)
   @           && n <= MAX_VALUE;
   @ ensures \result * \result == n;
   @*/
public static final int sqrt(int n) { ... }
```

Choosing exact boundary value for  $n$  amounts to computing result

Computing exact boundary values can be difficult or impossible!

## Test Generation Principle 5

Include boundary values of ordered domains as class representatives

# Implicit Disjunctions, Part I

## Example (Binary search, target not found)

```
/*@ public normal_behavior
  @ requires (\forall int i; 0 < i && i < array.length
             ==> array[i-1] <= array[i]);
  @         (\forall int i; 0 <= i && i < array.length
             ==> array[i] != target);
  @ ensures \result == -1;
  @*/
int search( int array[], int target ) { ... }
```

# Implicit Disjunctions, Part I

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/*@ public normal_behavior
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int search( int array[], int target ) { ... }
```

No disjunction in precondition!?

# Implicit Disjunctions, Part I

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             @           ==> array[i] != target);
  @ ensures  \result == -1;
  @*/
int search( int array[], int target ) { ... }
```

We can freely choose array in precondition!

## Test Generation Principle 6

Values of variables without explicit quantification can be freely chosen

# Data Generation Principles

## Test Generation Principle 6

Values of variables without explicit quantification can be freely chosen

### Systematic enumeration of values by data generation principle

Assume declaration: `int[] ar;`, then the array `ar` is

1. either the null array: `int[] ar = null;`
2. or the empty array of type `int`: `int[] ar = new int[0];`
3. or an `int` array with one element
  - 3.a `int[] ar = { MIN_VALUE };`
  - 3.b ...
  - 3.ω `int[] ar = { MAX_VALUE };`
4. or an `int` array with two elements ...
- n. or an `int` array with  $n$  elements ...

# Combining the Test Generation Principles

## Example (Binary search, target found)

```
requires (\exists int i; 0 <= i && i < array.length
          && array[i] == target)      &&
(\forall int i; 0 < i && i < array.length
  ==> array[i-1] <= array[i]);
```

## Apply test generation principles

- ▶ Use data generation principle for `int` arrays
- ▶ Choose equivalence classes and representatives of `int`, `int []`:

`int []` empty array, singleton, two elements

`int` 0, 1

- ▶ Generate all test cases that make precondition true

# Combining the Test Generation Principles

## Example (Binary search, target found)

```
requires (\exists int i; 0 <= i && i < array.length
          && array[i] == target)      &&
(\forall int i; 0 < i && i < array.length
  ==> array[i-1] <= array[i]);
```

- ▶ empty array: precondition cannot be made true, no test case
- ▶ singleton array, target must be only array element

```
array = { 0 }; target = 0;
```

```
array = { 1 }; target = 1;
```

- ▶ two-element **sorted** array, target occurs in array, four tests

```
array = { 0,0 }; target = 0;
```

```
array = { 0,1 }; target = 0;
```

etc.

# Implicit Disjunctions, Part II

## Example (Copy)

```
/*@ public normal_behavior
   @ requires src != null && dst != null;
   @ ensures ...
   @*/
static void java.util.Collections.copy (List src, List dst)
```

# Implicit Disjunctions, Part II

## Example (Copy)

```
/*@ public normal_behavior
   @ requires src != null && dst != null;
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   @*/
static void java.util.Collections.copy (List src, List dst)
```

## Aliasing and Exceptions

In JAVA object references `src`, `dst` can be **aliased**, ie, `src==dst`

- ▶ Admission of aliasing often unintended in contract

Forgotten protection against runtime exceptions

```
src.length <= dst.length
```

# Implicit Disjunctions, Part II

## Example (Copy)

```
/*@ public normal_behavior
   @ requires src != null && dst != null;
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   @*/
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## Test Generation Principle 7

Generate separate test cases that enforce aliasing and raising exceptions

# The Postcondition as Test Oracle

## Oracle Problem in Automated Testing

How to determine automatically whether a test run succeeded?

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How to determine automatically whether a test run succeeded?

The “ensures” clause of a JML contract tells exactly this provided that “requires” clause is true for given test case

## Test Generation Principle 1

Test data must make required precondition true

## Test Generation Principle 8

Use “ensures” clauses (postconditions) of JML contracts as test oracles

# Executable JML Expressions

How to determine whether a JML expression is true in a program state?

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## Example

```
\exists int i; 0 <= i && i < ar.length && ar[i] == target
```

is of the form

```
\exists int i; guard(i) && test(i)
```

- ▶ guard() is JAVA guard expression with fixed upper/lower bound
- ▶ test() is executable JAVA expression

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- ▶ `guard()` is JAVA guard expression with fixed upper/lower bound
- ▶ `test()` is executable JAVA expression

## Guarded existential JML quantifiers as Java (Example)

```
for (int i = 0; 0 <= i && i < ar.length; i++) {  
    if (ar[i]==target) { return true; }  
} return false;
```

# Executable JML Expressions

How to determine whether a JML expression is true in a program state?

## Example

`\exists int i; 0 <= i && i < ar.length && ar[i] == target`  
is of the form

`\exists int i; guard(i) && test(i)`

- ▶ `guard()` is JAVA guard expression with fixed upper/lower bound
- ▶ `test()` is executable JAVA expression

## Guarded existential JML quantifiers as Java (General)

```
for (int i = lowerBound; guard(i); i++) {  
    if (test(i)) { return true; }  
} return false;
```

# Executable JML Expressions

How to determine whether a JML expression is true in a program state?

## Example

`\exists int i; 0 <= i && i < ar.length && ar[i] == target`  
is of the form

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- ▶ `guard()` is JAVA guard expression with fixed upper/lower bound
- ▶ `test()` is executable JAVA expression

## Guarded JML quantifiers as Java

- ▶ Universal quantifiers treated similarly (exercise)
- ▶ Alternative JML syntax for quantifiers ok as well:

`\exists int i; guard(i) ; test(i)`

# Summary

- ▶ Black box vs White box testing
- ▶ Black box testing ~ Specification-based Test Generation
- ▶ Systematic test case generation from JML contracts guided by Test Generation Principles
- ▶ Only generate test cases that make precondition true
- ▶ Each operation contract and each disjunction in precondition gives rise to a separate test case
- ▶ Coverage criteria, decision coverage
- ▶ Large/infinite datatypes represented by class representatives
- ▶ Values of free variables supplied by Data Generation Principle
- ▶ Create separate test cases for potential aliases and exceptions
- ▶ Postconditions of contract provide test oracle
- ▶ Turn pre- and postconditions into executable JAVA code

# What Next?

## Remaining Problems of ATCG

1. How to automate specification-based test generation?
2. Generated test cases have no relation to implementation

# What Next?

## Remaining Problems of ATCG

1. How to automate specification-based test generation?
2. Generated test cases have no relation to implementation

1. Tools jml-junit and jtest discussed in Exercises
2. Code-based test generation that uses symbolic execution of IUT