

ANNULAR AND BOUNDARY REDUCING DEHN FILLINGS

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§0. INTRODUCTION

Surfaces of non-negative Euler characteristic, i.e., spheres, disks, tori and annuli, play a special role in the theory of 3-dimensional manifolds. For example, it is well known that every (compact, orientable) 3-manifold can be decomposed into canonical pieces by cutting it along essential surfaces of this kind [K], [M], [Bo], [JS], [Jo1]. Also, if (as in [Wu3]) we call a 3-manifold that contains no essential sphere, disk, torus or annulus *simple*, then Thurston has shown [T1] that a 3-manifold M with non-empty boundary is simple if and only if M with its boundary tori removed has a hyperbolic structure of finite volume with totally geodesic boundary. For closed 3-manifolds M , the Geometrization Conjecture [T1] asserts that M is simple if and only if M is either hyperbolic or belongs to a certain small class of Seifert fiber spaces.

Because of their importance, a good deal of attention has been directed at the question of when surfaces of non-negative Euler characteristic can be created by Dehn filling. To describe this, let M be a simple 3-manifold, with a torus boundary component $\partial_0 M$. Let α be the isotopy class of an essential simple loop (or *slope*) on $\partial_0 M$. Recall that the manifold obtained from M by α -Dehn filling is $M(\alpha) = M \cup V_\alpha$, where V_α is a solid torus, glued to M by a homeomorphism between $\partial_0 M$ and ∂V_α which identifies α with the boundary of a meridian disk of V_α . We are interested in obtaining restrictions on when $M(\alpha)$ fails to be simple. Although clearly little can be said in general about a single Dehn filling, if one considers pairs of non-simple fillings $M(\alpha)$, $M(\beta)$ then it turns out that the *distance* $\Delta(\alpha, \beta)$ between the two slopes α and β (i.e., their minimal geometric intersection number) is quite small, and hence a given M can have only a small number of non-simple fillings. More precisely, if $M(\alpha)$, $M(\beta)$ contain essential surfaces F_α, F_β of non-negative Euler characteristic, then for each of the ten possible pairs of homeomorphism classes of F_α, F_β one can obtain upper bounds on $\Delta(\alpha, \beta)$. In the present paper we deal with the case where F_α is an annulus and F_β is a disk, and prove the following theorem.

Theorem 0.1. *Let M be a simple 3-manifold such that $M(\alpha)$ is annular and $M(\beta)$ is boundary reducible. Then $\Delta(\alpha, \beta) \leq 2$.*

The assumption that M is a simple manifold can be replaced by the weaker assumption that it is boundary irreducible and anannular, see Corollary 5.5. The bound is sharp: infinitely many examples of simple 3-manifolds M with $M(\alpha)$ annular, $M(\beta)$ a solid torus, and $\Delta(\alpha, \beta) = 2$ are given in [MM2]. See also [EW].

1991 *Mathematics Subject Classification.* Primary 57N10.

¹ Partially supported by NSF grant #DMS 9626550.

² Partially supported by NSF grant #DMS 9802558.

$M(\beta) \backslash M(\alpha)$	S	D	A	T
S	1	0	2	3
D		1	2	2
A			5	5
T				8

Table 1: Upper bounds on $\Delta(\alpha, \beta)$

Theorem 0.1 completes the determination of the best possible upper bounds on $\Delta(\alpha, \beta)$ in all ten cases. These are shown in Table 1, where S , D , A and T indicate that the manifold $M(\alpha)$ or $M(\beta)$ contains an essential sphere, disk, annulus or torus, respectively. References for these bounds are: (S, S) : [GL3] (see also [BZ]); (S, D) : [Sch]; (S, A) : [Wu3]; (S, T) : [Wu1], [Oh]; (D, D) : [Wu2]; (D, T) : [GL4]; (A, A) , (A, T) and (T, T) : [Go]. Examples showing that the bounds are best possible can be found in: (S, S) : [GLi]; (D, D) : [Be] and [Ga]; (S, A) , (D, A) and (D, T) : [HM]; (S, T) : [BZ2]; (T, A) and (A, A) : [GW]; (T, T) : [T2] and [Go].

Here is a sketch of the proof of Theorem 0.1. It has been shown by Qiu [Qiu] that $\Delta \leq 3$, so we assume $\Delta = 3$, and try to get a contradiction. Let A and B be an essential annulus and an essential disk in $M(\alpha)$ and $M(\beta)$, and let P and Q be the intersection of A and B with M , respectively. Let p, q be the number of boundary components of P, Q on the torus $\partial_0 M$. Denote by K_β the core of the Dehn filling solid torus V_β in $M(\beta)$.

In Section 2 we consider the special case that $K = K_\beta$ is a *1-arch* knot, which means that it can be isotoped to a union of two arcs C_1 and C_2 , such that C_1 lies on $\partial M(\beta)$, and C_2 is disjoint from the compressing disk $B = F_\beta$ of $\partial M(\beta)$. In this case the manifold $M(\beta)$ is homeomorphic to a manifold X_C obtained by adding a 2-handle to a certain manifold X along a curve C . This changes a Dehn surgery problem to a handle addition problem, and we will use a theorem of Eudave-Muñoz to show that in this case the annulus $A = F_\alpha$ can be chosen to intersect the knot K_α at most twice, that is, $p \leq 2$.

As usual, the intersection of $P \cap Q$ defines graphs G_A, G_B on A and B . In Section 3 the “representing all types” technique developed in [GL1–GL4] is modified to suit the case that the intersection graphs have boundary edges. It will be proved that when $\Delta \geq 2$, either G_A represents all types, or G_B contains a great web. The first possibility is impossible because it would lead to a boundary reducing disk of $M(\beta)$ which has less intersection with K_β , hence G_B must contain a great web. This great web is then used in Section 4 to show that if $p \geq 3$ then the knot K_β is a 1-arch knot. Combined with the result of Section 2, this proves Theorem 0.1 in the generic case that $p \geq 3$. Finally in Section 5 the case $p \leq 2$ is ruled out, completing the proof of Theorem 0.1.

The authors would like to thank John Luecke for helpful conversations.

§1. PRELIMINARIES

Recall that a 3-manifold X is *boundary reducible* if its boundary, denoted by ∂X , is compressible in X , in which case a compressing disk of ∂X is also called a *boundary reducing disk* of X . A surface of non-positive Euler characteristic in X is *essential* if it is incompressible, ∂ -incompressible, and is not boundary parallel; a sphere (resp. disk) is essential if it is a reducing sphere (resp. boundary reducing disk.)

Let M be a simple 3-manifold, with a torus boundary component $\partial_0 M$. Let α, β be slopes on $\partial_0 M$ such that $M(\alpha)$ is annular and $M(\beta)$ is boundary reducible. Let A be an essential annulus in $M(\alpha)$, and let B be an essential disk in $M(\beta)$. These give rise to a punctured annulus $P = A \cap M$ and a punctured disk $Q = B \cap M$ in M , where $\partial_0 P = P \cap \partial_0 M$ consists of p copies of α , and $\partial_0 Q = Q \cap \partial_0 M$ consists of q copies of β . We assume that A, B, P, Q are chosen so that p and q are minimal. Note that p, q are positive because M is simple. Now isotope P and Q to minimize $|P \cap Q|$, the number of components of $P \cap Q$. Then no arc component of $P \cap Q$ is boundary parallel in P or Q ; no circle component of $P \cap Q$ bounds a disk in P or Q ; and each component of $\partial_0 P$ meets each component of $\partial_0 Q$ in $\Delta = \Delta(\alpha, \beta)$ points.

Ruifeng Qiu showed in [Qiu] that if M is a simple manifold, $M(\alpha)$ is annular and $M(\beta)$ is boundary reducible, then $\Delta \leq 3$. Thus to prove Theorem 0.1, we need only rule out the possibility that $\Delta = 3$. *In this paper except in Section 3, we will assume that $\Delta = 3$, and proceed to get a contradiction.* Results in Section 3 have been proved in a broader setting, so they can be used in the future.

Regarding the components of $\partial_0 P, \partial_0 Q$ as fat vertices, we get graphs G_A, G_B in A, B respectively, where the edges of G_A and G_B are the arc components of $P \cap Q$ that have at least one endpoint on $\partial_0 M$. Let $J = A$ or B . An edge of G_J is an *interior* edge if each of its endpoints lies on a vertex of G_J , and a *boundary* edge if one of its endpoints lies on a vertex of G_J and the other lies on ∂J . The faces of G_J correspond in the usual way to components of $J - \text{Int}N(G_J)$. A face of G_J is an *interior* face if it does not meet ∂J ; otherwise it is a *boundary* face. Thus the edges in the boundary of an interior face are interior edges, while the boundary of a boundary disk face contains some boundary edges. Denote by \widehat{G}_J the reduced graph of G_J , in which each parallel family of edges is replaced by a single edge.

Let u_1, \dots, u_p be the vertices of G_A , labeled successively when traveling along the Dehn filling solid torus V_α . Each u_i is given a sign according to whether V_α passes A from the positive side or negative side at this vertex. Two vertices u_i, u_j are *parallel* if they have the same sign, otherwise they are *antiparallel*. The vertices v_1, \dots, v_q of G_B are labeled and signed similarly.

If e is an edge of G_A with an endpoint on u_i , then the endpoint is labeled j if it is on $\partial u_i \cap \partial v_j$. Thus when going around ∂u_i , the labels of the edge endpoints appear as $1, 2, \dots, q$ repeated Δ times. The edge endpoints of G_B are labeled similarly.

A cycle in G_A or G_B is a *Scharlemann cycle* if it bounds a disk with interior disjoint from the graph, and all the edges in the cycle have the same pair of labels $\{i, i+1\}$ at their two endpoints, called the *label pair* of the Scharlemann cycle. A pair of edges $\{e_1, e_2\}$ is an *extended Scharlemann cycle* if there is a Scharlemann cycle $\{e'_1, e'_2\}$ such that e_i is parallel and adjacent to e'_i .

We use $N(X)$ to denote a regular neighborhood of a subset X in a given manifold.

Lemma 1.1. (Properties of G_A .)

- (1) (The Parity Rule) *An edge connects parallel vertices on G_A if and only if it connects antiparallel vertices on G_B .*
- (2) G_A does not have q parallel interior edges.
- (3) G_A contains no Scharlemann cycles.
- (4) Each label $x \in \{1, \dots, q\}$ appears at most once among the endpoints of a family \mathcal{E} of parallel edges in G_A connecting parallel vertices; in particular, \mathcal{E} contains at most $q/2$ edges.
- (5) No pair of edges are parallel on both G_A and G_B .

Proof. (1) This is on [CGLS, Page 279].

(2) If G_A contains q parallel interior edges, then the core of the Dehn filling solid torus in $M(\beta)$ would be a cable knot, in which case M contains an essential annulus, contradicting the assumption. See the proof of [GLi, Proposition 1.3].

(3) This follows from [CGLS, Lemma 2.5.2].

(4) If some label appears twice among the endpoints of a family of parallel edges connecting a pair of parallel vertices, then there is a Scharlemann cycle among this family, contradicting (3). See [CGLS, Lemma 2.6.6].

(5) By [Go, Lemma 2.1], no two interior edges are parallel on both G_A and G_B . If a pair of boundary edges are parallel on both graphs, then they cut off a disk on each of P and Q , whose union is an annulus in M , which is essential because its intersection with $\partial_0 M$ is a curve intersecting α at a single point. This contradicts the assumption that M is simple. \square

Lemma 1.2. (Properties of G_B .)

(1) *If G_B has a Scharlemann cycle, then A is a separating annulus, and p is even. Moreover, the subgraph of G_A consisting of the edges of the Scharlemann cycle and their vertices is not contained in a subdisk of A .*

(2) *If $p > 2$, then G_B has no extended Scharlemann cycle. Any two Scharlemann cycles of G_B have the same label pair.*

Proof. (1) This follows from the proof of [CGLS, Lemma 2.5.2]. It was shown that using the disk bounded by the Scharlemann cycle one can find another annulus A' in $M(\alpha)$ which has fewer intersections with the Dehn filling solid torus, and is cobordant to A , so if A were nonseparating then A' would still be essential, which would contradict the minimality of p . If the subgraph G consisting of the edges of a Scharlemann cycle and their end vertices is contained in a disk in A then $A' \cup A$ bounds a connected sum of $A \times I$ and a lens space, so A being essential implies that A' is essential, which again contradicts the minimality of p .

(2) This is [Wu3, Lemma 5.4(2) – (3)]. If G_B has an extended Scharlemann cycle or two Scharlemann cycles with distinct label pairs, then one can find another essential annulus in $M(\alpha)$ having fewer intersection with K_α , which would contradict the minimality of p . \square

§2. 1-ARCH KNOTS

Let $K = K_\beta$ be the core of the Dehn filling solid torus in $M(\beta)$. The knot K is a *1-arch* knot (with respect to B) if K is isotopic to a union of two arcs C_1 and C_2 , such that C_1 lies on $\partial M(\beta)$, and C_2 is disjoint from a compressing disk B of $\partial M(\beta)$.

Fix an orientation of K so that when traveling along K with this orientation one meets the fat vertices v_1, \dots, v_q successively. Let $K[i]$ be the point $K \cap v_i$, and for

$i \neq j$, let $K[i, j]$ be the oriented arc segment of K starting from $K[i]$ and ending at $K[j]$. Thus $K = K[i, j] \cup K[j, i]$.

Lemma 2.1. *If G_A contains q parallel boundary edges, then $K = K_\beta$ is a 1-arch knot.*

Proof. Without loss of generality we may assume that the interior endpoints of the parallel boundary edges e_1, \dots, e_q are successively labeled $1, \dots, q$. Let D be the disk on P cut off by e_1 and e_q . Then D can be extended into the Dehn filling solid torus $N(K)$ to get a disk D' in $M(\beta)$ such that $\partial D' = e'_1 \cup K[1, q] \cup e'_q \cup C_1$, where e'_i is an arc on B containing e_i , connecting $K[i]$ to the endpoint of e_i on ∂A , and $C_1 = D \cap \partial A$ lies on $\partial M(\beta)$. Now K is isotopic to $C_1 \cup (e'_1 \cup K[q, 1] \cup e'_q)$ via the disk D' . Let $C_2 = e'_1 \cup K[q, 1] \cup e'_q$. After a slight isotopy one can make B disjoint from C_2 , as desired. \square

Lemma 2.2. *Suppose \widehat{G}_A has a vertex u of valency 4, such that one of the four edges of \widehat{G}_A incident to u is a boundary edge, and the two edges adjacent to it are interior edges. Then either G_B contains a Scharlemann cycle, or $K = K_\beta$ is a 1-arch knot.*

Proof. Let $\widehat{e}_1, \widehat{e}_2, \widehat{e}_3, \widehat{e}_4$ be the four edges of \widehat{G}_A incident to u , and assume that \widehat{e}_2 is a boundary edge. By Lemmas 2.1 and 1.1(2) we may assume that each \widehat{e}_i represents at most $q - 1$ parallel edges of G_A . Now each label appears at most twice among the endpoints at u of edges represented by \widehat{e}_2 or \widehat{e}_4 , hence all labels appear on endpoints at u of edges represented by \widehat{e}_1 or \widehat{e}_3 . Suppose $\widehat{e}_1, \widehat{e}_3$ connect u to u' and u'' , respectively. If both u' and u'' are antiparallel to u , then by the parity rule each vertex v on G_B is incident to an edge connecting it to a parallel vertex, with label u at its endpoint at v . By [CGLS, Lemmas 2.6.3 and 2.6.2] this implies that G_B contains a great u -cycle, hence a Scharlemann cycle, and we are done. Also, notice that u' and u'' cannot both be parallel to u , otherwise by Lemma 1.1(4) each of \widehat{e}_1 and \widehat{e}_3 represents at most $q/2$ edges, and since each of \widehat{e}_2 and \widehat{e}_4 represents at most $q - 1$ edges, this would contradict the fact that the total valency of u in G_A is $\Delta q = 3q$. (This also takes care of the case that $\widehat{e}_1 = \widehat{e}_3$ is a loop at u .) Therefore, we may assume that u' is parallel to u , and u'' is antiparallel to u . Since the total number of edges represented by $\widehat{e}_1 \cup \widehat{e}_2 \cup \widehat{e}_3$ is more than $2q$, we can choose $2q$ successive edges at u , forming a subgraph as shown in Figure 2.1. One can now use [Wu2, Lemma 2.2] and the proof of [Wu2, Lemma 3.4] to show that there is a disk D in $M(\beta)$ with $\partial D = K[1, q] \cup \alpha_1 \cup C_1 \cup \alpha_2$, where α_1 and α_2 are arcs on the compressing disk B connecting $K[1]$ and $K[q]$ to ∂B , and C_1 is an arc on $\partial M(\beta)$. As in the proof of Lemma 2.1, this implies that K is a 1-arch knot. \square

We need the following result of Eudave-Mun̄oz in the proof of Proposition 2.4. If C is a simple loop on the boundary of a 3-manifold X , denote by X_C the manifold obtained by adding a 2-handle to X along the curve C .

Lemma 2.3. *Let X be an irreducible, orientable 3-manifold with ∂X compressible, and C a simple closed curve on ∂X such that $\partial X - C$ is incompressible. Suppose X_C contains an essential annulus A' . Then it contains an essential annulus A which intersects the attached 2-handle in at most two disks. Furthermore, if A' is nonseparating, then A can be chosen to be disjoint from the attached 2-handle.*

Proof. This is essentially [Eu, Theorem 1]. The theorem there says that under the above assumption, either one can find A to be disjoint from the attached 2-handle,

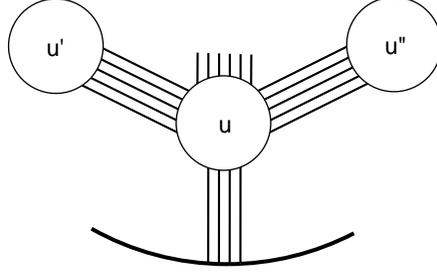


Figure 2.1

or after sliding the cocore σ of the attached 2-handle over itself to get a 1-complex τ , one can find an essential annulus A which intersects τ at a single point. Moreover, if A' is nonseparating, then A is disjoint from τ (see also [Jo2, Sch]). Sliding τ back to σ , we see that A is isotopic to an annulus intersecting σ at most twice. See also the remarks after the statement of Theorem 2 in [Eu]. \square

Proposition 2.4. *If $K = K_\beta$ is a 1-arch knot in $M(\beta)$, then $p \leq 2$, and A is a separating annulus in $M(\alpha)$.*

Proof. The first part of the proof here is the same as that in the proof of [Wu2, Proposition 1]. Suppose K is isotopic to $C_1 \cup C_2$ as in the definition of 1-arch knot. Let Y be the manifold obtained by adding a 1-handle H_1 to $M(\beta)$ along two disks centered at ∂C_1 , and let C be a simple closed curve on ∂Y obtained by taking the union of $C_1 \cap \partial Y$ and an arc on ∂H_1 . Let K' be the union of C_2 and the core of the 1-handle H_1 . Then after adding a 2-handle H_2 to Y along C the 1-handle and the 2-handle cancel each other and we get a manifold M' homeomorphic to the original manifold $M(\beta)$, with the knot K identified to K' ; hence we have a homeomorphism of pairs $(M(\beta), K) \cong (M', K')$. Let W (denoted by Q in [Wu2]) be the manifold obtained from Y by Dehn surgery on K' along the slope α . Then $M(\alpha)$ is homeomorphic to the manifold W_C obtained by adding the 2-handle H_2 to W along the curve C . It was shown in [Wu2, Lemmas 1.2 and 1.3] that ∂W is compressible, and $\partial W - C$ is incompressible in W when $\Delta \geq 2$.

Since $M(\beta)$ is ∂ -reducible, by [Sch] the manifold $M(\alpha) = W_C$ is irreducible. This implies that W is irreducible because a reducing sphere in a manifold always remains a reducing sphere after 2-handle additions. Therefore we can apply Lemma 2.3 and conclude that there is an essential annulus A in $M(\alpha) = W_C$ intersecting the attached 2-handle H_2 in $n \leq 2$ disks; moreover, if A is nonseparating, then it is disjoint from H_2 . Our goal is to show that A also intersects the knot K_α in $M(\alpha)$ in two or zero points, respectively.

We assume $n = 2$, the cases $n = 0$ or 1 are similar. Let D_1, D_2 be the disks $A \cap H_2$, and let F be the twice punctured annulus $A - \text{Int}(D_1 \cup D_2)$ in W . A meridian disk D of the 1-handle H_1 gives rise to a nonseparating essential annulus $D_0 = D \cap X$ in the manifold $X = Y - \text{Int}N(K') = W - \text{Int}N(K_\alpha)$. Let $F_0 = F \cap X$. Form intersection graphs G_D and G_F in the usual way, i.e, G_F has $F \cap N(K_\alpha)$ as fat vertices, G_D has a single vertex $D \cap N(K_\beta)$, and the edges of G_D and G_F are

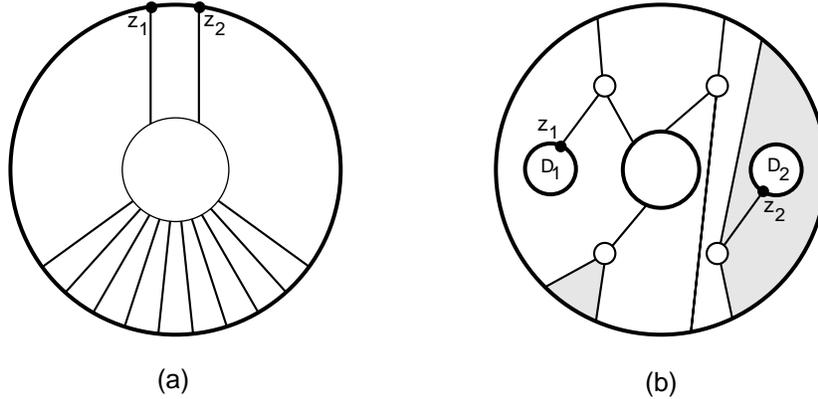


Figure 2.2

the arc components of $D_0 \cap F_0$ which has at least one endpoint on the fat vertices. See Figure 2.2. Choose D and F so that D_0 intersects F_0 minimally. Then each fat vertex of G_F has valency $\Delta = 3$, and the only vertex x of G_D has valency $3t$, where t is the number of vertices of G_F . Note that, since A is essential in $M(\alpha)$, we have $t \geq p$. As usual, there are no trivial loops. Hence each edge of G_D connects x to ∂D .

Each of ∂D_1 and ∂D_2 intersects ∂D at a single point, which we denote by z_1, z_2 , as indicated by the dark dots in Figure 2.2(a) and (b). They divide ∂D into two arcs α_1, α_2 , one of which, say α_1 , lies in $N(C)$, which is the attaching region of the two handle H_2 above. Hence the interior of α_1 is disjoint from ∂F . It follows that all the endpoints of the edges of G_D on ∂D lie on the arc α_2 , as shown in Figure 2.2(a).

Now suppose G_F has $t \geq 3$ vertices. Since each of ∂D_1 and ∂D_2 is adjacent to at most one edge, there is a vertex of G_F with two edges connecting it to the same component of $\partial A \subset \partial F$. An outermost such vertex has a pair of edges a_1, a_2 on G_F , cutting off a region B_1 on F_0 which is either a disk, or a once punctured disk containing one of $\partial D_1, \partial D_2$, as shown by the two shaded regions in Figure 2.2(b). (If B_1 contains both ∂D_i , choose another outermost vertex.) They also cut off a disk B_2 on D_0 which, by the property in the last paragraph, has boundary disjoint from α_1 . Therefore, $B_1 \cup B_2$ is either an annulus or a once punctured annulus in X , with one boundary component γ a curve on $\partial N(K)$, another a curve on $\partial X - \partial N(K)$ disjoint from C , and a possible third curve parallel to C . After capping off the last component by a disk in the attached 2-handle H_2 , the surface becomes an annulus in X_C . However, since $X_C = M$, and since the boundary component γ of the annulus on the torus $\partial N(K) = \partial_0 M$ is an essential curve, (essential because γ is the union of an arc in α and an arc in β , and α intersects β minimally,) this contradicts the fact that M is ∂ -irreducible and anannular.

When $n = 0$, since $t \geq p > 0$, there is a pair of edges which are parallel on both graphs G_D and G_F . As shown above, this would give rise to an essential annulus in M , which would contradict the simplicity of M . Hence this case does not happen.

In particular, this and Lemma 2.3 show that $M(\alpha)$ cannot contain a nonseparating annulus. \square

§3. REPRESENTING TYPES

Denote by $\mathbf{q} = \{1, \dots, q\}$ the set of labels of the vertices of G_B . We have the concept of a \mathbf{q} -type etc. from [GL1]. An interior face of G_A represents a \mathbf{q} -type τ if it is a disk and represents τ in the sense of [GL1]. We say G_A represents τ if some interior face of G_A represents τ .

Theorem 3.1. *G_A does not represent all \mathbf{q} -types.*

Proof. See [GL4, Proof of Theorem 2.2]. The proof works for any essential surface F in $M(\alpha)$ (in [GL4] F was a torus). A set of representatives of all \mathbf{q} -types contains a set \mathcal{D} of interior faces of G_F which can be used to surger Q tubed along the annuli corresponding to the corners of the faces in \mathcal{D} , contradicting the minimality of q . \square

A *web* in G_B is a non-empty connected subgraph Λ of G_B such that all the vertices of Λ have the same sign, and such that there are at most p edge endpoints at vertices of Λ which are not endpoints of edges in Λ . Note that a web may have boundary edges.

Let U be a component of $B - N(\Lambda)$ that meets ∂B . Then $D = B - U$ is a *disk bounded by Λ* . Λ is a *great web* if there is a disk bounded by Λ such that Λ contains all the edges of G_B that lie in D .

Remark. If there are no boundary edges, then these definitions coincide with those in [GL2, Section 2]. The following is the analog in our present setting of [GL2, Theorem 2.3].

Theorem 3.2. *Suppose $\Delta \geq 2$. Let L be a subset of \mathbf{q} , and τ be a non-trivial L -type such that*

- (i) *all elements of $C(\tau)$ have the same sign, and*
- (ii) *all elements of $A(\tau)$ have the same sign.*

If $G_A(L)$ does not represent τ then G_B contains a web Λ such that the set of vertices of Λ is a subset of either $C(\tau)$ or $A(\tau)$.

Proof. Regard $G_A(L)$ as a graph in S^2 , by capping off the boundary components of A with two additional fat vertices v_1, v_2 .

Define a directed graph $\Gamma = \Gamma(\tau)$ as follows. The vertices of Γ are the fat vertices of $G_A(L)$ plus v_1, v_2 , together with dual vertices of $G_A(L)$ (one in the interior of each face of $G_A(L)$.) The edges of Γ join each dual vertex to the fat vertices in the boundary of the corresponding face. The edges of Γ are oriented as follows: If an edge e has an endpoint on a vertex of $G_A(L)$, then it is oriented according to the type τ , (as in [GL1], where Γ is denoted by $\Gamma(T)^*$); if e has an endpoint on v_1 or v_2 , orient e so that no dual vertex in a boundary face of $G_A(L)$ is a sink or source of Γ . See Figure 3.1.

By Glass' index formula (see [Gl]) applied to Γ , we have

$$\sum_{\text{vertices}} I(v) + \sum_{\text{faces}} I(f) = \chi(S^2) = 2.$$

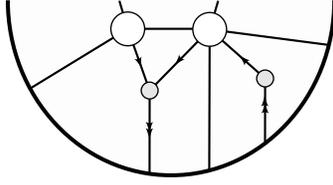


Figure 3.1

Assume $G_A(L)$ does not represent τ . Then no dual vertex in Γ is a sink or source. Hence

$$\sum_{v \text{ dual}} I(v) \leq 0.$$

Let

$$\begin{aligned} c(\tau) &= \# \text{ clockwise switches } (= \# \text{ anticlockwise switches}) \text{ of } \tau \\ c(v_i) &= \# \text{ clockwise switches } (= \# \text{ anticlockwise switches}) \text{ at } v_i, \quad i = 1, 2 \end{aligned}$$

For v a vertex of $G_A(L)$, we have

$$I(v) = 1 - \Delta c(\tau).$$

Also,

$$I(v_i) = 1 - c(v_i), \quad i = 1, 2.$$

Therefore, the number of switch edges, including all switch boundary edges, is at least

$$\begin{aligned} \sum I(f) &\geq 2 + p(\Delta c(\tau) - 1) + (c(v_1) - 1) + (c(v_2) - 1) \\ &= p(\Delta c(\tau) - 1) + c(v_1) + c(v_2) \end{aligned} \quad (*)$$

Since the number of switch edge endpoints is twice the number of switch edges, this is also a lower bound for the number of (say) clockwise switch edge endpoints. The total number of clockwise switches is $p\Delta c(\tau) + c(v_1) + c(v_2)$, so the number of clockwise switches that are not endpoints of clockwise switch edges is at most p .

Since $\Delta \geq 2$, the right hand side of (*) is positive. Let Λ be a component of the subgraph of G_B consisting of the edges corresponding to the clockwise switch edges of $G_A(L)$. Then at most p edge endpoints at vertices of Λ do not belong to edges of Λ . Thus Λ is a web, as described. \square

The following is the analog of [GL2, Theorem 2.5].

Theorem 3.3. *Suppose $\Delta \geq 2$. Let Λ be either (i) a web in G_B , or (ii) the empty set. In case (i), let D be a disk bounded by Λ , and in case (ii), let $D = B$. Let L be the set of vertices of $G_B - \Lambda$ that lie in D . Then either G_B contains a great web or $G_A(L)$ represents all L -types.*

Proof. Basically, this follows from the proof of [GL2, Theorem 2.5]. We indicate briefly how this goes.

We prove the result by induction on $|L|$.

Let τ be an L -type. We show that if $G_A(L)$ does not represent τ then G_B contains a great web. There are two cases.

CASE 1. τ is trivial. Proceed as in [GL2, Proof of Theorem 2.5]. Let $\widehat{\Lambda}$ be a component of the subgraph of G_B consisting of vertices J , all interior edges with both endpoints on vertices in J , and all boundary edges with one endpoint on a vertex in J .

(a) $\widehat{\Lambda}$ is a web. Argue as in [GL2], with “faces” meaning “interior faces”.

(b) $\widehat{\Lambda}$ is not a web. Again the argument in [GL2] remains valid. More precisely, since $\widehat{\Lambda}$ is not a web, there are more than p edges of G_B connecting a vertex of $\widehat{\Lambda}$ to an antiparallel vertex. Let Σ be the subgraph of G_A consisting of the vertices of G_A together with those edges. Note that these are interior edges of G_A , connecting parallel vertices. Applying Euler’s formula to Σ , a graph in A , gives

$$V - E + \sum \chi(f) = 0.$$

Therefore

$$\sum \chi(f) = E - V > p - p = 0.$$

Hence Σ has a disk face, which must be an interior face. This face then contains a face of $G_A(L)$ representing the trivial type.

CASE 2. τ is non-trivial. Here the argument in [GL2] goes through essentially without change, (using Theorem 3.2), where we always interpret “face” as “interior face”. In particular, [GL2, Lemmas 2.4 and 2.6] carry over in this way. \square

Theorems 3.1 and 3.3 (in case (ii)) imply:

Corollary 3.4. *If $\Delta \geq 2$, then G_B contains a great web. \square*

§4. THE GENERIC CASE

Let Λ be a great web in G_B given by Corollary 3.4, and let D be a disk bounded by Λ with the property in the definition of a great web. Let x be a label of the vertices of G_A , and let Λ_x be the subgraph of Λ consisting of all vertices of Λ and all edges in Λ with an endpoint labeled x . Let V be the number of vertices of Λ . A *ghost endpoint* of Λ is an endpoint, at a vertex of Λ , of an edge of G_B which does not belong to Λ . A *ghost endpoint* of Λ_x is a ghost endpoint of Λ labeled x . (It is called a ghost label in [GL2].) By the definition of a web, Λ has at most p ghost endpoints.

By a *monogon* we mean a disk face with one edge in its boundary, and by a *bigon* we mean a disk face with two edges in its boundary.

Lemma 4.1. (Cf. [GL2, Lemma 4.2]) *If Λ_x has at least $3V - 2$ edges then Λ_x contains a bigon in D .*

Proof. Let Ω be the graph in S^2 obtained from Λ_x by regarding ∂B as a vertex. Then Ω has $V + 1$ vertices, E edges (= number of edges of Λ_x), and the faces of Ω are the faces of Λ_x in D together with an additional face f_0 . Note that f_0 is not a monogon. Suppose Λ_x contains no bigon in D .

First suppose f_0 is not a bigon. Then $2E \geq 3F$, where $F = \sum \chi(f)$ summed over all faces of Ω . Also,

$$(V + 1) - E + F = 2.$$

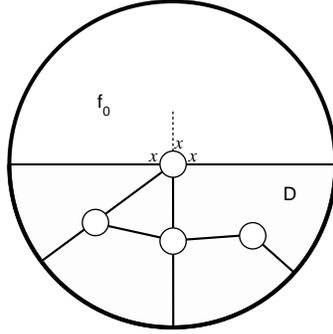


Figure 4.1

Hence

$$1 = V - E + F \leq V - E + \frac{2E}{3} = V - \frac{E}{3},$$

giving $3 \leq 3V - E$, i.e, $E \leq 3V - 3$, contrary to assumption.

Now suppose f_0 is a bigon; see Figure 4.1. Then Λ_x has at most one ghost endpoint. Therefore $E \geq 3V - 1$. Also, Since Λ_x has no bigon in D , we have

$$2E \geq 3F - 1.$$

Hence, as before,

$$1 = V - E + F \leq V - E + \frac{2E + 1}{3} = V - \frac{E}{3} + \frac{1}{3}.$$

Therefore $3 \leq 3V - E + 1$, implying $E \leq 3V - 2$, a contradiction. \square

Remark. One can show that the conclusion of Lemma 4.1 still holds if we only assume that Λ_x has at least $3V - 3$ edges, but Lemma 4.1 will suffice for our purposes.

Lemma 4.2. Λ_x contains a bigon in D for at least $2p/3$ labels x .

Proof. (Cf. [GL2, Theorem 4.3]). By Lemma 4.1, if Λ_x does not contain a bigon in D then Λ_x has at most $3V - 3$ edges. Since the vertices of Λ_x are all parallel, by the parity rule no edge of Λ_x has both endpoints labeled x , so among the endpoints of edges of Λ_x , at most $3V - 3$ are labeled x . Since $\Delta = 3$, this means that Λ_x has at least 3 ghost endpoints. Since the total number of ghost endpoints in Λ is at most p by the definition of a great web, there can be at most $p/3$ such labels x . Hence for at least $2p/3$ labels x , Λ_x does contain a bigon in D . \square

Note that a boundary bigon in Λ_x gives $p + 1$ parallel boundary edges in G_B .

In the remainder of this section, we assume that $p \geq 3$.

Lemma 4.3. *For some label x , Λ_x contains a boundary bigon.*

Proof. A bigon face of Λ_x in D is either a boundary bigon or an interior bigon. The latter is either an order 2 Scharlemann cycle in G_B , or contains an extended Scharlemann cycle. The second is impossible by Lemma 1.2(2). When p is odd, the first is also impossible (Lemma 1.2(1)), and when p is even, any two Scharlemann cycles have the same label pair (Lemma 1.2(2)). Hence, by Lemma 4.2, the number of labels x such that Λ_x contains a boundary bigon is at least

$$\begin{cases} 2p/3 \geq 2 \times 3/3 = 2, & p \text{ odd;} \\ 2p/3 - 2 \geq 2 \times 4/3 - 2 = 2/3, & p \text{ even.} \end{cases}$$

Hence there is at least one label x with the stated property. \square

Corollary 4.4. (a) *Every vertex of G_A has a boundary edge incident to it.*

(b) *G_A has a vertex with two non-parallel boundary edges.*

Proof. (a) By Lemma 4.3, Λ_x contains a boundary bigon for some x , which gives rise to $p + 1$ parallel boundary edges in G_B . Hence each label of $\mathbf{p} = \{1, \dots, p\}$ appears at the endpoint of some boundary edge of G_B , and the result follows.

(b) By Lemma 1.1(5), the two boundary edges in a boundary bigon of Λ_x are nonparallel on G_A . \square

Lemma 4.5. *\widehat{G}_A has no vertex of valency at most 3.*

Proof. G_A has at most $q - 1$ parallel interior edges by Lemma 1.1(2), and at most $q - 1$ parallel boundary edges by Lemma 2.1, Proposition 2.4, and the assumption that $p \geq 3$. Since the total valency of each vertex of G_A is $3q$, the result follows. \square

Corollary 4.6. *\widehat{G}_A has no vertex with two boundary edges going to the same component of ∂A .*

Proof. Consider an outermost such vertex, with E the corresponding subdisk of A . Doubling E along the two boundary edges in question and applying [CGLS, Lemma 2.6.5] gives a vertex in the interior of E of valency at most 3, contradicting Lemma 4.5. \square

Lemma 4.7. *\widehat{G}_A has a vertex v of valency 4, such that no two boundary edges of \widehat{G}_A at v have endpoints adjacent on ∂v .*

Proof. By Corollaries 4.4(b) and 4.6, \widehat{G}_A has at least one vertex with two boundary edges going to different components of ∂A . Cut A along all such pairs of edges; we get a certain number (≥ 1) of disk regions. If there are no vertices in the interior of any of these regions, then every vertex of \widehat{G}_A satisfies the conclusion of the lemma; (recall that there is no vertex of valency ≤ 3 by Lemma 4.5). So consider a region with a non-zero number of vertices in its interior; see Figure 4.2(a). Note that each vertex v in the interior of the region is incident to exactly one boundary edge, hence we need only show that some v has valency 4.

If the number of vertices in the interior of the region is 1 or 2, the result is obvious. So suppose there are at least 3 such vertices. Delete v_2 and all edges incident to it, and push v_1 inwards and attach a boundary edge to it, as shown in Figure 4.2(b). Applying assertion (*) in the proof of Lemma 2.6.5 in [CGLS] to the resulting graph, we conclude that there is a vertex $v \neq v_1$ of valency at most

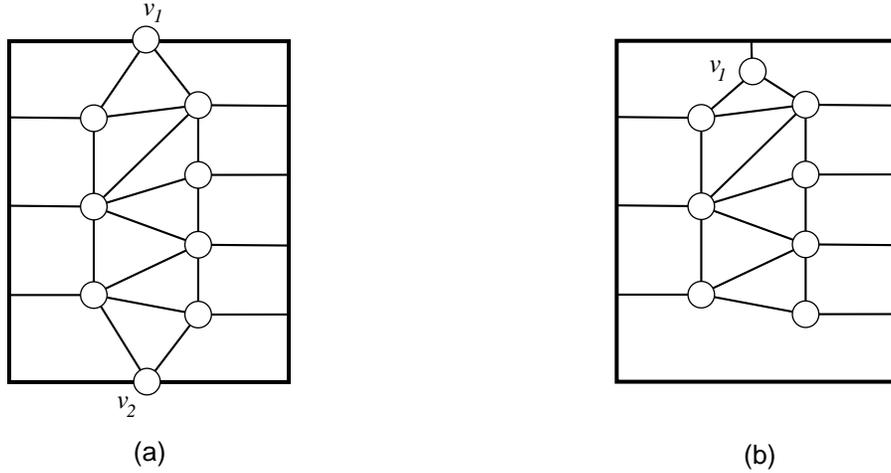


Figure 4.2

3. Since there is at most one edge joining v to v_2 , v has valency at most 4 (hence exactly 4) in the original graph \widehat{G}_A . \square

Proposition 4.8. *Theorem 0.1 is true if $p \geq 3$.*

Proof. Let v be the vertex of valency 4 given by Lemma 4.7. Since we have assumed $p \geq 3$, by Lemma 2.2 and Proposition 2.4, G_B contains a Scharlemann cycle. Suppose the Scharlemann cycle has label pair $\{1, 2\}$. Then by Lemma 1.2(1), p is even, hence $p \geq 4$, and the edges of the Scharlemann cycle are not contained in a disk on A . Thus in \widehat{G}_A there are two edges connecting v_1 to v_2 , as shown in Figure 5.1(b). They separate the two boundary components of the annulus A , so no other vertex is incident to two edges going to different boundary components of A . It follows from Corollary 4.6 that the only possible vertices of \widehat{G}_A with two boundary edges are v_1 and v_2 . Since there are no Scharlemann cycles on any other label pair, and no extended Scharlemann cycles (Lemma 1.2(2)), the only labels x for which Λ_x has a bigon in D are 1 and 2. Hence by Lemma 4.2, we have $2p/3 \leq 2$, i.e. $p \leq 3$. But we have just shown that $p \geq 4$, which is a contradiction. \square

§5. THE CASE THAT $p \leq 2$

After Proposition 4.8, it remains to consider the case that the graph G_A on the annulus A has at most two vertices. In this section we will consider this remaining case, and complete the proof of Theorem 0.1. As before, we assume that $\Delta = 3$.

Lemma 5.1. *If $p \leq 2$, then $p = 2$, and the two vertices of G_A are antiparallel.*

Proof. First assume $p = 1$. Then A is a nonseparating annulus in $M(\alpha)$. The reduced graph \widehat{G}_A consists of one vertex, at most one loop, and at most two boundary edges. By Lemma 1.1(4) the number of endpoints of loops is at most q . Since the total valency of the vertex is $3q$, there exist q parallel boundary edges. By Lemma

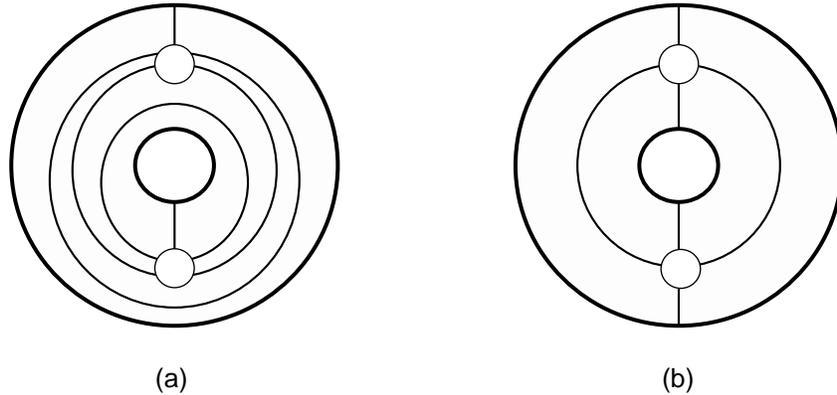


Figure 5.1

2.1 K_β is a 1-arch knot. However, by Proposition 2.4 in this case $M(\alpha)$ contains no nonseparating annulus, a contradiction.

Now assume $p = 2$ and the two vertices of G_A are parallel. Then again A is nonseparating in $M(\alpha)$. The reduced graph \widehat{G}_A is a subgraph of one of the two graphs shown in Figure 5.1, depending on whether or not \widehat{G}_A has a loop. Since the two vertices are parallel, by Lemma 1.1(4) each family of parallel interior edges contains at most $q/2$ edges, hence in both cases there is a family of at least q parallel boundary edges. As above, this implies that K_β is a 1-arch knot, hence contradicts Proposition 2.4 and the fact that A is nonseparating. \square

We may now assume that $p = 2$ and the two vertices of G_A are antiparallel. Suppose W is a submanifold of $M(\alpha)$ containing K_α . We use $\partial_i W$ to denote the closure of $\partial W - \partial M(\alpha)$, which is the frontier of W in $M(\alpha)$, and call it the *interior boundary*.

If D is a disk embedded (improperly) in $M(\alpha)$ such that $D \cap W = D \cap \partial_i W$ is a single arc C on the boundary of D , and $\partial D - C$ lies on $\partial M(\alpha)$, then the pair $(W \cup N(D), K_\alpha)$ is homeomorphic to (W, K_α) , with $\partial_i(W \cup N(D))$ identified to $\partial_i W$ cut along the arc C . This observation will be useful in the proof of Lemma 5.4.

For the purpose of this section, we define an *extremal component* of a subgraph Λ of G_B to be a component Λ_0 such that there is an arc γ cutting B into B_1 and B_2 , with $B_1 \cap \Lambda = \Lambda_0$.

Lemma 5.2. *If $p = 2$ and the two vertices of G_A are antiparallel, then each vertex of G_B is incident to a boundary edge. In particular, each face of G_B is a disk.*

Proof. The reduced graph \widehat{G}_A is a subgraph of that shown in Figure 5.1(a) or (b). In case (b), each of the interior edges of \widehat{G}_A represents at most q edges of G_A , hence each label appears at most four times at endpoints of interior edges. It follows that each vertex of G_B is incident to at least two boundary edges.

In case (a), consider the edge endpoints at a vertex v of G_A . Let s be the number of boundary edges at v , and let t be the number of loops based at v . Observe that

if $s < q$ but $s + 2t > q$ then some label would appear twice among the endpoints of the parallel loops, which would contradict Lemma 1.1(4). If $s + 2t \leq q$, then the two nonloop edges of \widehat{G}_A would represent $3q - (s + 2t) \geq 2q$ edges, which would contradict Lemma 1.1(2). Therefore we must have $s \geq q$, which implies that each vertex of G_B is incident to a boundary edge.

If some face f of G_B is not a disk, then the vertices inside of a nontrivial loop in f would have no boundary edges, which would contradict the above conclusion. \square

Lemma 5.3. *Suppose $p = 2$ and the two vertices of G_A are antiparallel. Then there is a vertex v_0 of \widehat{G}_B with the following properties.*

- (1) v_0 has valency 2 or 3 in \widehat{G}_B , and belongs to a single boundary edge e of \widehat{G}_B .
- (2) If the valency of v_0 is 3, then the face opposite to the boundary edge is an interior face.
- (3) One of the two faces of \widehat{G}_B containing e intersects ∂B in a single arc.

Proof. By Lemma 5.2 each vertex of G_B belongs to a boundary edge. Consider an extremal component C of G_B , and let \widehat{C} be its reduced graph. Let \widehat{C}' be its corresponding component in \widehat{G}_B . Note that \widehat{C}' and \widehat{C} are almost identical, except that one of the vertices v' of \widehat{C}' may have two parallel boundary edges, in which case \widehat{C} can be obtained from \widehat{C}' by amalgamating these two edges together.

Note that \widehat{C} must have at least two vertices, for otherwise, since C is extremal, the vertex would have 6 parallel boundary edges in G_B , two of which would also be parallel on G_A because \widehat{G}_A has at most four boundary edges, which would then contradict Lemma 1.1(5). Hence each vertex of C is incident to at least one interior edge and one boundary edge, so the valency of each vertex of \widehat{C} is at least 2. Modify \widehat{C} as follows. If some vertex v of \widehat{C} satisfies condition (1) but not (2), add a boundary edge to v in the face opposite to the boundary edge at v . Having done this for all v , we get a graph \widehat{C}'' , which is still a reduced graph, with at least one boundary edge incident to each vertex. Now using (*) in the proof of [CGLS, Lemma 2.6.5] and arguing directly when \widehat{C}'' has only two or three vertices, we see that \widehat{C}'' contains at least two vertices, each of which has valency 2 or 3 in \widehat{C}'' and belongs to a single boundary edge of \widehat{C}'' . At least one of these two vertices, say v_0 , is not the vertex v' above, hence it has property (1) when considered as a vertex in \widehat{G}_B . By the definition of \widehat{C}'' , v_0 automatically has property (2). To prove (3), notice that if both faces containing e intersect ∂B in more than one arc, then C would not be an extremal component. \square

Lemma 5.4. *If $p = 2$, and the two vertices of G_A are antiparallel, then ∂M is a union of tori.*

Proof. Let W_0 be a regular neighborhood of $A \cup K_\alpha$. Since K_α intersects A in two points of different signs, $\partial_i W_0$ has two components F_b, F_w , each being a twice punctured torus. The annulus A cuts W_0 into two components W_0^b and W_0^w (with $W_0^b \supset F_b$), which will be called the black region and the white region, respectively. If E is a disk face of G_B or more generally a disk in $M(\alpha)$, then E is said to be *black* (resp. *white*) if $E \cap W_0$ lies in the black (resp. white) region.

Suppose D is a compressing disk of F_b in $M(\alpha) - \text{Int}W_0$. If ∂D is a nonseparating curve on F_b , then after adding the 2-handle $N(D)$ to W_0 , the surface F_b becomes an annulus. If ∂D is separating on F_b , then it is not parallel to a boundary curve of F_b because ∂F_b is parallel to ∂A and A is incompressible in $M(\alpha)$; thus ∂D must cut

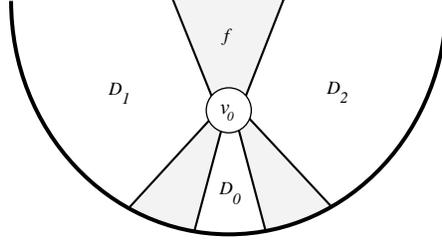


Figure 5.2

F_b into a once punctured torus and a thrice punctured sphere, and after adding the 2-handle $N(D)$ the surface F_b becomes the union of a torus S_1 and an annulus S_2 . Since M is simple, S_1 either is boundary parallel or bounds a solid torus, and S_2 must be boundary parallel, because ∂S_2 is parallel to ∂A and A is incompressible, which implies that S_2 is incompressible. In any case, we have shown that if F_b is compressible in $M(\alpha) - \text{Int}W_0$ then there is a component C_b of $M(\alpha) - \text{Int}W_0$ such that $C_b \cap W_0 = F_b$ and $C_b \cap \partial M(\alpha)$ is either an annulus or the union of an annulus and a torus. Similarly for F_w . In particular, if both F_b and F_w are compressible in $M(\alpha) - \text{Int}W_0$, then $\partial M(\alpha)$ is a union of tori, and we are done. From now on, we will assume that F_w is incompressible in $M(\alpha) - \text{Int}W_0$ and show that this will lead to a contradiction. Note that the assumption implies that G_B has no interior white face: For, by Lemma 5.2 all faces of G_B are disks, and since G_A has no trivial loops, the boundary of an interior face is always an essential curve on $\partial_i W_0$; hence an interior white face would give rise to a compressing disk of F_w in $M(\alpha) - \text{Int}W_0$.

Let v_0 be a vertex of G_B given by Lemma 5.3. By Lemma 5.3(1), v_0 has valency 2 or 3 in \widehat{G}_B . First assume that v_0 has valency 2 in \widehat{G}_B . Then the interior edge e of \widehat{G}_B incident to v_0 must represent exactly two edges of G_B : It cannot represent more than two edges, otherwise there would be two interior faces of different colors, contradicting the fact that G_B has no interior white face. It cannot represent only one edge of G_B , otherwise v_0 would have five parallel boundary edges, which would contradict Lemma 1.1(5) because \widehat{G}_A has at most four boundary edges. Thus the part of G_B near v_0 is as shown in Figure 5.2, where f is the interior (black) face bounded by the two edges represented by e .

Now assume that v_0 has valency 3 in \widehat{G}_B . Then by Lemma 5.3(2) the face f of \widehat{G}_B opposite to the boundary edge at v_0 is an interior face. Thus f is a black face, and each of the interior edges of \widehat{G}_B incident to v_0 represents only one edge of G_B as otherwise there would be a white interior face. Hence again the part of G_B near v_0 is as shown in Figure 5.2.

Consider the white boundary faces D_0, D_1, D_2 as shown in Figure 5.2. By Lemma 5.3(3), we may assume that D_1 intersects ∂B in a single arc. Let $C_i, i = 0, 1$, be the arc $D_i \cap F_w$. Then C_0, C_1 are essential arcs on F_w . Moreover, since C_0 intersects a meridian of K_α exactly once, while C_1 intersects it at least twice, they are nonparallel. Recall that $\partial_i(W_0 \cup N(D_0 \cup D_1))$ is obtained from $\partial_i W_0$ by cutting along $C_0 \cup C_1$. Since C_0 and C_1 are nonparallel, they cut F_w into one or

two annuli, which must be boundary parallel because we have assumed that F_w is incompressible and M is simple. It follows that the whole surface F_w is boundary parallel. Now a meridian disk of $N(K_\alpha)$ in the white region corresponds to a disk in $M(\alpha)$ intersecting the curve K_α in a single point, which gives rise to an essential annulus in M , contradicting the fact that M is anannular. \square

Proof of Theorem 0.1. By Proposition 4.8, we may assume that $p \leq 2$. By Lemmas 5.1 and 5.4, ∂M is a union of tori. Since $M(\beta)$ is ∂ -reducible, either it is reducible or it is a solid torus. In the first case the result follows from [Wu3, Theorem 5.1]. So we assume that $M(\beta)$ is a solid torus. In particular, $\partial M(\alpha)$ is a single torus T . The boundary of the annulus A cuts T into two annuli A_1, A_2 . If some $A \cup A_i$ is an essential torus in $M(\alpha)$ then $M(\alpha)$ is toroidal, so the result follows from [GL4]. If each $A \cup A_i$ is inessential, then it bounds a solid torus (note that it cannot be boundary parallel, otherwise A would be boundary parallel). It follows that $M(\alpha)$ is a Seifert fiber space with orbifold a disk with two singular points. It was shown in [MM1, Theorem 1.2] that if $M(\alpha)$ is a Seifert fiber space and $M(\beta)$ is a solid torus then $\Delta \leq 1$. This completes the proof of Theorem 0.1. \square

In the proof of Theorem 0.1, we assumed that the manifold M is simple. However, the conditions that M is irreducible and atoroidal can be removed from the assumptions.

Corollary 5.5. *Suppose M is anannular and boundary irreducible. If $M(\alpha)$ is annular and $M(\beta)$ is boundary reducible, then $\Delta(\alpha, \beta) \leq 2$.*

Proof. First assume that M is irreducible but toroidal. Since M is anannular, by the canonical splitting theorem of Jaco-Shalen-Johannson (see [JS, p. 157]) there is a set of essential tori \mathcal{T} cutting M into a manifold M' such that each component of M' is either a Seifert fiber space or a simple manifold. If the component X containing the boundary torus $\partial_0 M$ is Seifert fibered, then it contains an essential annulus consisting of Seifert fibers, with both boundary components on $\partial_0 M$, so M would be annular, contradicting our assumption. So assume X is simple. Since $M(\beta)$ is boundary reducible, by looking at a boundary reducing disk B which has minimal intersection with \mathcal{T} , one can see that $X(\beta)$ must be boundary reducible. Similarly one can show that $X(\alpha)$ is either boundary reducible or annular. Applying Theorem 0.1 and [Wu2, Theorem 1] to X , we have $\Delta \leq 2$.

If M is reducible, split along a maximal set of reducing spheres to get an irreducible manifold M' . By an innermost circle argument one can show that $M'(\alpha)$ is annular and $M'(\beta)$ is boundary reducible, so the result follows from that for irreducible manifolds. \square

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