

Recursively determined representing measures for bivariate truncated moment sequences

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Abstract

A theorem of Bayer and Teichmann [BT] implies that if a finite real multisequence $\beta \equiv \beta^{(2d)}$ has a representing measure, then the associated moment matrix M_d admits positive, recursively generated moment matrix extensions M_{d+1}, M_{d+2}, \dots . For a bivariate recursively determinate M_d , we show that the existence of positive, recursively generated extensions M_{d+1}, \dots, M_{2d-1} is sufficient for a measure. Examples illustrate that all of these extensions may be required to show that β has a measure. We describe in detail a constructive procedure for determining whether such extensions exist. Under mild additional hypotheses, we show that M_d admits an extension M_{d+1} which has many of the properties of a positive, recursively generated extension.

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1 Introduction

Let $\beta \equiv \beta^{(2d)} := \{\beta_{ij}\}_{i,j \geq 0, i+j \leq 2d}$ denote a real bivariate multisequence of degree $2d$. The Truncated Moment Problem seeks conditions on β for the existence of a positive Borel measure μ on \mathbb{R}^2 such that

$$\beta_{ij} = \int_{\mathbb{R}^2} x^i y^j d\mu \quad (i, j \geq 0, \quad i + j \leq 2d). \quad (1.1)$$

A result of [CF5] shows that β admits a finitely atomic *representing measure* μ (as in (1.1)) if and only if $M_d \equiv M_d(\beta)$, the *moment matrix* associated with β , admits a *flat extension* M_{d+k+1} , i.e., an extension to a positive semidefinite moment matrix M_{d+k+1} such that $\text{rank } M_{d+k+1} = \text{rank } M_{d+k}$. The extension of this result to general representing measures follows from a theorem of C. Bayer and J. Teichmann [BT], which implies that if β has a

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representing measure, then it has a finitely atomic representing measure (cf. [F1, Section 2], [CFM, Section 1]). At present, for a general moment matrix, there is no known concrete test for the existence of a flat extension M_{d+k+1} . In this note, for the class of bivariate *recursively determinate* moment matrices, we present a detailed analysis of an algorithm of [F1] that can be used in numerical examples to determine the existence or nonexistence of flat extensions (and representing measures). This algorithm determines the existence or nonexistence of positive, recursively generated extensions M_{d+1}, \dots, M_{2d-1} , at least one of which must be a flat extension in the case when there is a measure. Theorem 2.5 shows that there are sequences $\beta^{(2d)}$ for which the first flat extension occurs at M_{2d-1} , so all of the above extensions must be computed in order to recognize that there is a measure. This result stands in sharp contrast to traditional truncated moment theorems (concerning representing measures supported in \mathbb{R} , $[a, b]$, $[0, +\infty)$, or in a planar curve of degree 2), which express the existence of a measure in terms of tests closely related to the original moment data (cf. Remark 2.6 below and [CF1], [CF2], [CF3], [CFM], [F1]). Here we see that, at least within the framework of moment matrix extensions, we may need to go far from the original data to resolve the existence of a measure. In Theorems 2.3 and 2.13 we show that under mild additional hypotheses on M_d , the implementation of each extension step, from M_{d+j} to M_{d+j+1} , leading to a flat extension M_{d+k+1} , consists of simply verifying a matrix positivity condition.

Let $\mathcal{P}_d \equiv \mathbb{R}[x, y]_d$ denote the bivariate real polynomials of degree at most d . For $p \in \mathcal{P}_d$, $p(x, y) \equiv \sum_{i,j \geq 0, i+j \leq d} a_{ij} x^i y^j$, let $\hat{p} := (a_{ij})$ denote the vector of coefficients with respect to the basis for \mathcal{P}_d consisting of monomials in degree-lexicographic order, i.e., $1, x, y, x^2, xy, y^2, \dots, x^d, \dots, y^d$. Let $L_\beta : \mathcal{P}_{2d} \rightarrow \mathbb{R}$ denote the Riesz functional, defined by $L_\beta(\sum_{i,j \geq 0, i+j \leq 2d} a_{ij} x^i y^j) :=$

$\sum a_{ij} \beta_{ij}$. The moment matrix M_d , whose rows and columns are indexed by the monomials in \mathcal{P}_d , is defined by $\langle M_d \hat{p}, \hat{q} \rangle := L_\beta(pq)$ ($p, q \in \mathcal{P}_d$). We denote the successive rows and columns of M_d by $1, X, Y, \dots, X^d, \dots, Y^d$; thus, the entry in row $X^i Y^j$, column $X^k Y^\ell$, which we denote by $\langle X^k Y^\ell, X^i Y^j \rangle$, is equal to $\beta_{i+k, j+\ell}$. We may denote a linear combination of rows or columns by $p(X, Y) := \sum a_{ij} X^i Y^j$ for some $p \equiv \sum a_{ij} x^i y^j \in \mathcal{P}_d$; note that $p(X, Y) = M_d \hat{p}$. We say that M_d is *recursively generated* if $\ker M_d$ has the following ideal-like property:

$$p, q, pq \in \mathcal{P}_d, p(X, Y) = 0 \implies (pq)(X, Y) = 0. \quad (1.2)$$

If β has a representing measure, then M_d is positive semidefinite and recursively generated [CF5] (and in one variable these conditions are sufficient for the existence of a representing measure [CF1]). Moreover, from [BT], β actually admits a *finitely atomic* representing measure μ , which therefore has finite moments of all orders; it follows that M_d admits positive, recursively generated moment matrix extensions of all orders, namely $M_{d+1}[\mu], \dots, M_{d+k}[\mu], \dots$. Let us consider a moment matrix extension

$$M_{d+1} \equiv \begin{pmatrix} M_d & B(d+1) \\ B(d+1)^T & C(d+1) \end{pmatrix},$$

where the block $B(d+1)$ includes new moments of degree $2d+1$ (as well as old moments of degrees $d+1, \dots, 2d$), and block $C(d+1)$ consists of new moments of degree $2d+2$. We denote the columns of $B(d+1)$ by X^{d+1}, \dots, Y^{d+1} , and we say that $(M_d \ B(d+1))$ is

recursively generated if (1.2) holds in its column space, but with $p, q, pq \in \mathcal{P}_{d+1}$. M_{d+1} is positive semidefinite if and only if (i) M_d is positive semidefinite; (ii) $\text{Ran } B(d+1) \subseteq \text{Ran } M_d$ (equivalently, $B(d+1) = M_d W$ for some matrix W); (iii) $C(d+1) \succeq C^b := W^T M_d W$ (cf. [CF2]). (Here and in the sequel, for a real symmetric matrix A , we will write $A \succeq 0$ (resp. $A \succ 0$) to denote that A is positive semidefinite (resp. positive semidefinite and invertible).) If $M_{d+1} \succeq 0$, then we also have (iv) each dependence relation in $\text{Col } M_d$ (the column space of M_d) extends to $\text{Col } M_{d+1}$. In the sequel we say that M_{d+1} is an *RG extension* if properties (i), (ii), and (iv) hold and M_{d+1} is recursively generated (so, in particular, $(M_d B(d+1))$ is recursively generated). In the sequel, we provide sufficient conditions for *RG extensions*; note that to verify that an *RG extension* is positive semidefinite and recursively generated, it is only necessary to verify condition (iii).

For a general M_d , a significant difficulty in determining the existence of a flat extension M_{d+k+1} is that there may be infinitely many positive and recursively generated extensions M_{d+1} . If one such extension does not admit a subsequent flat extension, this does not preclude the possibility that some other extension does. In the sequel, we focus on the class of recursively determinate moment matrices (*RD*) introduced in [F1] (cf. [F2]). These are characterized by the property that there can be at most one positive, recursively generated extension, and there is a concrete procedure (described below) for determining the existence or nonexistence of this extension. Since such an extension is also recursively determinate, we may proceed iteratively to determine the existence or nonexistence of positive and recursively generated extensions

$$M_{d+1}, \dots, M_{2d-1}. \quad (1.3)$$

As we discuss below, the existence of the extensions in (1.3) is equivalent to the existence of a flat extension M_{d+k+1} and, in fact, one of the extensions in (1.3) is a flat extension of M_d . (If M_{j+1} is positive semidefinite, then M_j is positive semidefinite and recursively generated [CF5], so, using also [BT], it follows that (1.3) is equivalent to the existence of a positive semidefinite extension M_{2d} .)

A bivariate moment matrix M_d admits a block decomposition $M_d = (B[i, j])_{0 \leq i, j \leq d}$, where

$$B[i, j] = \begin{pmatrix} \beta_{i+j,0} & \cdots & \beta_{i,j} \\ \vdots & \ddots & \vdots \\ \beta_{j,i} & \cdots & \beta_{0,i+j} \end{pmatrix}.$$

Thus, $B[i, j]$ is constant on each cross-diagonal; we refer to this as the *Hankel property*. Note that in the extension M_{d+1} , $B(d+1) = (B[i, d+1])_{0 \leq i \leq d}$, and all of the new moments of degree $2d+1$ appear within block $B[d, d+1]$, either in column X^{d+1} (the leftmost column) or in column Y^{d+1} (on the right). Similarly, all new moments of degree $2d+2$ appear in column X^{d+1} or column Y^{d+1} of $C(d+1)$ ($= B[d+1, d+1]$). In the sequel, by a *column dependence relation* we mean a linear dependence relation of the form $X^i Y^j = r(X, Y)$, where $\deg r \leq i+j$ and each monomial term in r strictly precedes $x^i y^j$ in the degree-lexicographic order; we say that such a relation is *degree reducing* if $\deg r < i+j$. A bivariate moment matrix M_d is *recursively determinate* if there are column dependence relations of the form

$$X^n = p(X, Y) \quad (p \in \mathcal{P}_{n-1}, n \leq d) \quad (1.4)$$

and

$$Y^m = q(X, Y) \quad (q \in \mathcal{P}_m, q \text{ has no } y^m \text{ term}, m \leq d), \quad (1.5)$$

or with similar relations with the roles of p and q reversed. In the sequel, we state the main results (Theorems 2.3 and 2.5) with p and q as in (1.4)-(1.5), but these results are valid as well with the roles of p and q reversed. In Section 2 we show that if M_d is recursively determinate, then the only possible positive, recursively generated (or merely *RG*) extension is completely determined by column relations $X^{d+1} = (x^{d+1-n}p)(X, Y)$ and $Y^{d+1} = (y^{d+1-m}q)(X, Y)$.

The most important case of recursive determinacy occurs when M_d is positive and *flat*, i.e., $\text{rank } M_d = \text{rank } M_{d-1}$ (equivalently, each column of degree d can be expressed as a linear combination of columns of strictly lower degree). A fundamental result of [CF2] shows that in this case M_d admits a unique flat extension M_{d+1} (and a corresponding rank M_d -atomic representing measure). In this paper, we stay within the framework of recursive determinacy, but relax the flatness condition, and study the extent to which positive, recursively generated extensions exist.

Our main results are Theorems 2.3 and 2.13, which give sufficient conditions for *RG* extensions, and Theorem 2.5, which shows that the number of extension steps leading to a flat extension is sometimes proportional to the degree of the moment problem. Theorem 2.3 shows that if M_d is positive and recursively generated, and if all column dependence relations arise from (1.4) or (1.5) via recursiveness and linearity, then M_d admits a unique *RG* extension. In general, this extension need not be positive semidefinite (see the discussion preceding Example 1.1), but if $d = n + m - 2$, then this extension is actually a flat extension, so β admits a representing measure (Corollary 2.2). Additionally, we show in Theorem 2.13 that if M_d is positive semidefinite, recursively generated, and recursively determinate, and if all column dependence relations are degree-reducing, then M_d again admits a unique *RG* extension. However, we show in Example 2.12 that if all of the column relations are degree-reducing except that $\deg q = m$, then M_d need not even admit a block $B(d+1)$ consistent with recursiveness for $(M_d B(d+1))$. In Theorem 2.5 we show that for each d , there exists $\beta \equiv \beta^{(2d)}$, with $M_d(\beta) \in RD$, such that in the sequence of positive, recursively generated extensions, M_{d+1}, \dots, M_{2d-1} , the first flat extension is M_{2d-1} , so the determination that a measure exists takes the maximum possible number of extension steps. Moreover, at each extension step, M_{d+i} satisfies the hypotheses of Theorem 2.3, so it is guaranteed in advance that the next extension M_{d+i+1} is well-defined and recursively generated; only its positivity needs to be verified. In general, however, the existence of a positive, recursively generated extension M_{d+1} does not imply the existence of a measure. In Section 3 we answer [F1, Question 4.19] by showing that if, under the hypotheses of Theorem 2.3, M_d does admit a positive, recursively generated extension M_{d+1} , then M_{d+1} may also satisfy the conditions of Theorem 2.3, but need not admit a positive, recursively generated extension M_{d+2} , and thus M_d may fail to have a measure.

We conclude this section by reviewing and illustrating [F1, Algorithm 4.10] concerning extensions of recursively determinate bivariate moment matrices. We may assume that M_d is positive and recursively generated, for otherwise there is no representing measure. (We note that in numerical problems, positivity and recursiveness can easily be verified using elementary linear algebra.) To define block $B(d+1)$ for an extension M_{d+1} , note that blocks $B[0, d+1], \dots, B[d-1, d+1]$ consist of old moments from M_d . To define moments of degree $2d+1$ for block $B[d, d+1]$, we first use (1.4) and recursiveness to define the “left band” of columns, $X^{n+i}Y^{d+1-i-n} := (x^i y^{d+1-i-n} p)(X, Y)$ ($0 \leq i \leq d+1-i$). In block $B[d, d+1]$, certain “new moments” in column $X^n Y^{d+1-n}$ can be moved up and to the right

along cross-diagonals until they reach row X^d (the top row of $B[d, d + 1]$) in columns of the “central band,” $X^{n-1}Y^{d+2-n}, \dots, X^{d+2-m}Y^{m-1}$. These values can then be used to define $\langle X^{d+1-m}Y^m, X^d \rangle$ (the entry in row X^d , column $X^{d+1-m}Y^m$) by means of

$$\langle X^{d+1-m}Y^m, X^d \rangle := \langle (x^{d+1-m}q)(X, Y), X^d \rangle. \quad (1.6)$$

This value may be moved one position down and to the left along its cross-diagonal and then used to define $\langle X^{d+1-m}Y^m, X^{d-1}Y \rangle := \langle (x^{d+1-m}q)(X, Y), X^{d-1}Y \rangle$. We repeat this process successively to complete the definition of column $X^{d+1-m}Y^m$ in $B[d, d + 1]$ as well as the definition of the central band of columns in this block. We next complete the definition of $B[d, d + 1]$ by successively defining the “right band” of columns, $X^{d-m}Y^{m+1}, \dots, Y^{d+1}$, using

$$X^{d+1-m-i}Y^{m+i} := (x^{d+1-m-i}y^i q)(X, Y) \quad (0 \leq i \leq d + 1 - m).$$

It is necessary to check that the values in the central and right bands, as just defined, are compatible with values in the left band, and, more generally, to verify that $B(d + 1)$ is a well-defined moment matrix block. If this fails to be the case, there is no measure. If $B(d + 1)$ is well-defined, we next check that $\text{Ran } B(d + 1) \subseteq \text{Ran } M_d$, for if this is not the case, then there is no measure. Assuming the range condition is satisfied, (1.4) and (1.5) will hold in the columns of $B(d + 1)^T$ (the transpose). We then apply recursiveness and the method used just above in defining $B[d, d + 1]$ to attempt to define $C(d + 1) \equiv B[d + 1, d + 1]$. Assuming that $C(d + 1)$ is well-defined, we further check that M_{d+1} is positive and recursively generated. If any of the preceding steps fails, there is no representing measure. Our main results (Theorems 2.3 and 2.13) show that if all column relations come from (1.4) or (1.5) via recursiveness and linearity, or if (1.4) - (1.5) hold and all column dependence relations are degree-reducing, then all of the preceding steps are guaranteed to succeed, except possibly the positivity of M_{d+1} ; thus M_{d+1} is at least an RG extension.

If M_{d+1} , as just defined, is positive and recursively generated, then, since it is also recursively determinate, we may apply the above procedure successively, in attempting to define positive and recursively generated extensions M_{d+2}, M_{d+3}, \dots . Note that the central band of degree d in M_d has $n + m - d - 1$ columns, $X^{n-1}Y^{d-n+1}, \dots, X^{d+1-m}Y^{m-1}$. In each successive extension M_{d+k} , the number of columns in the central band of degree $d + k$ is $n + m - d - 1 - k$. Thus, after at most $n + m - d - 1$ extension steps, either the extension process fails, and there is no measure, or the central band disappears and there is a flat extension, at or before M_{n+m-1} , and a measure. (Note that since $n, m \leq d$, this refines our earlier assertion that a flat extension occurs at or before M_{2d-1} .) Another estimate for the number of extension steps is based on the *variety* of M_d , defined as $\mathcal{V} \equiv \mathcal{V}(M_d) := \bigcap_{r \in \mathcal{P}_d, r(X, Y) = 0} \mathcal{Z}_r$, where \mathcal{Z}_r is

the set of real zeros of r . It follows from [F1] that the number of extension steps leading to a flat extension is at most $1 + \text{card } \mathcal{V} - \text{rank } M_d$. Note also that when a measure exists, it is supported inside \mathcal{V} [CF5], so its support is a subset of the finite real variety determined by $x^n - p(x, y)$ and $y^m - q(x, y)$.

Examples are known where the RG extension M_{d+1} is not positive semidefinite (cf. [F1, Example 4.18], [CFM, Theorem 5.2], both with $n = m = d = 3$, and the example of Section 3 (below), with $d = 5, n = m = 4$). We next present an example, adapted from [F2, Example 5.2], which illustrates the algorithm in a case leading to a measure.

Example 1.1. Let $d = 3$ and consider

$$M_3 = \begin{pmatrix} 1 & 0 & 0 & 1 & 2 & 5 & 0 & 0 & 0 & 0 \\ 0 & 1 & 2 & 0 & 0 & 0 & 2 & 5 & 14 & 42 \\ 0 & 2 & 5 & 0 & 0 & 0 & 5 & 14 & 42 & 132 \\ 1 & 0 & 0 & 2 & 5 & 14 & 0 & 0 & 0 & 0 \\ 2 & 0 & 0 & 5 & 14 & 42 & 0 & 0 & 0 & 0 \\ 5 & 0 & 0 & 14 & 42 & 132 & 0 & 0 & 0 & 0 \\ 0 & 2 & 5 & 0 & 0 & 0 & 5 & 14 & 42 & 132 \\ 0 & 5 & 14 & 0 & 0 & 0 & 14 & 42 & 132 & 429 \\ 0 & 14 & 42 & 0 & 0 & 0 & 42 & 132 & 429 & c \\ 0 & 42 & 132 & 0 & 0 & 0 & 132 & 429 & c & d \end{pmatrix}.$$

We have $M_3 \succeq 0$, $M_2 \succ 0$, and $\text{rank } M_3 = 8 \iff d = 2026881 - 2844c + c^2$. When $\text{rank } M_3 = 8$, then the two column relations are

$$Y = X^3$$

and

$$Y^3 = q(X, Y),$$

where $q(x, y) := (5715 - 4c)x + 10(-1428 + c)y - 3(-2853 + 2c)x^2y + (-1422 + c)xy^2$. Let $r_1(x, y) = y - x^3$ and $r_2(x, y) = y^3 - q(x, y)$. With these two column relations in hand, Theorem 2.3 guarantees the existence of a unique RG extension M_4 . To test the positivity of M_4 , we calculate the determinant of the 9×9 matrix consisting of the rows and columns of M_4 indexed by the monomials $1, x, y, x^2, xy, y^2, x^2y, xy^2, x^2y^2$. A straightforward calculation using *Mathematica* shows that three cases arise:

- (i) $c < 1429$: here $M_4 \not\succeq 0$, so M_3 admits no representing measure;
- (ii) $c = 1429$: here M_4 is a flat extension of M_3 , so by the main result in [CF2], M_3 admits an 8-atomic representing measure;
- (iii) $c > 1429$: here M_4 is a positive RG extension of M_3 with rank 9. Although M_4 is not a flat extension of M_3 , it nevertheless satisfies the hypotheses of Theorem 2.3, so Corollary 2.4 implies that M_4 admits a flat extension M_5 , and therefore M_3 has a 9-atomic representing measure. Moreover, since the original algebraic variety $\mathcal{V} \equiv \mathcal{V}(M_3)$ associated with M_3 , $\mathcal{Z}_{r_1} \cap \mathcal{Z}_{r_2}$, can have at most 9 points (by Bézout's Theorem), it follows that $\mathcal{V} = \mathcal{V}(M_5)$. This algebraic variety must have exactly 9 points, and thus constitutes the support of the unique representing measure for M_3 .

To illustrate this case, we take the special value $c = 1430$, so that $q(x, y) \equiv -5x + 20y - 21x^2y + 8xy^2$. Let $\alpha := \frac{1}{2}\sqrt{5} - 2\sqrt{5}$ and $\gamma := \sqrt{5}\alpha$. A calculation shows that $\mathcal{V} = \{(x_i, x_i^3)\}_{i=1}^9$, where $x_1 = 0$, $x_2 = \frac{1}{2}(-1 - \sqrt{5}) \approx -1.618$, $x_3 = \frac{1}{2}(1 - \sqrt{5}) \approx -0.618$, $x_4 = -x_3 \approx 0.618$, $x_5 = -x_2 \approx 1.618$, $x_6 = -\alpha - \gamma \approx -1.176$, $x_7 = -\alpha + \gamma \approx 0.449$, $x_8 = -x_7 \approx -0.449$ and $x_9 = -x_6 \approx 1.176$. M_3 satisfies the hypothesis of Theorem 2.3 with $n = m = 3$, so we proceed to generate the RG extension M_4 . This extension is uniquely determined by imposing the column relations $X^4 = XY$, $X^3Y = Y^2$, $XY^3 = (xq)(X, Y)$, and $Y^4 = (yq)(X, Y)$ (first in $(M_3 \ B(4))$, then in $(B(4)^T \ C(4))$). A calculation shows that, as expected, these relations unambiguously define a positive moment matrix M_4 with $\text{rank } M_4 = 9 (> 8 = \text{rank } M_3)$. It follows that M_3 admits no flat extension M_4 , so we proceed

to construct the RG extension M_5 , uniquely determined by imposing the relations $X^5 = X^2Y$, $X^4Y = XY^2$, $X^3Y^2 = Y^3$, $X^2Y^3 = (x^2q)(X, Y)$, $XY^4 = (xyq)(X, Y)$, $Y^5 = (y^2q)(X, Y)$. A calculation of these columns (first in $(M_4 \ B(5))$, then in $(B(5)^T \ C(5))$), shows that, as again expected, they do fit together to unambiguously define a moment matrix M_5 . From the form of $q(x, y)$, we see that M_5 is actually a flat extension of M_4 , in keeping with the above discussion. Corresponding to this flat extension is the unique, 9-atomic, representing measure $\mu \equiv \mu_{M_5}$ as described in [CF5]. Clearly, $\text{supp } \mu = \mathcal{V}$, so μ is of the form $\mu = \sum_{i=1}^9 \rho_i \delta_{(x_i, x_i^3)}$. To compute the densities, we use the method of [CF5] and find $\rho_1 = \frac{1}{5} = 0.2$, $\rho_2 = \rho_5 = \frac{-1+\sqrt{5}}{8\sqrt{5}} \approx 0.069$, $\rho_3 = \rho_4 = \frac{1+\sqrt{5}}{8\sqrt{5}} \approx 0.181$, $\rho_6 = \rho_9 = \frac{5+3\sqrt{5}}{40\sqrt{5}} \approx 0.131$, and $\rho_7 = \rho_8 = \frac{-5+3\sqrt{5}}{40\sqrt{5}} \approx 0.019$. Thus, the existence of a representing measure for $\beta^{(6)}$ is established on the basis of the extensions M_4 and M_5 , in keeping with Theorem 2.3. Note that in this case, the actual number of extensions leading to a flat extension can be computed as either $n+m-d-1$ or as $1 + \text{card } \mathcal{V} - \text{rank } M_3$, which is consistent with our earlier discussion. \square

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2 The extension of a bivariate RD positive moment matrix

In Theorem 2.3 (below) we show that a positive recursively determinate moment matrix $M_d \equiv M_d(\beta)$, each of whose column dependence relations is recursively generated by a relation of the form

$$X^n = p(X, Y), \quad (p \in \mathcal{P}_{n-1}) \quad (2.1)$$

or

$$Y^m = q(X, Y), \quad (q \in \mathcal{P}_m), \quad (2.2)$$

(where $n, m \leq d$ are fixed and q has terms $x^u y^v$ with $v < m$), always admits a unique RG extension

$$M_{d+1} \equiv \begin{pmatrix} M_d & B(d+1) \\ B(d+1)^T & C(d+1) \end{pmatrix}.$$

The main step towards Theorem 2.3 is the following result, which shows that M_d (as above) admits an extension block $B(d+1)$ that is consistent with the structure of a positive, recursively generated moment matrix extension M_{d+1} .

Theorem 2.1. *Suppose the bivariate moment matrix $M_d(\beta)$ is positive and recursively generated, with column dependence relations generated entirely by (2.1) and (2.2) via recursiveness and linearity. Then there exists a unique moment matrix block $B(d+1)$ such that $(M_d \ B(d+1))$ is recursively generated and $\text{Ran } B(d+1) \subseteq \text{Ran } M_d$.*

The hypothesis implies that the column dependence relations in M_d are precisely those of the form

$$X^{n+i}Y^j = (x^i y^j p)(X, Y) \quad (i, j \geq 0, i+j+n \leq d) \quad (2.3)$$

and

$$X^k Y^{m+l} = (x^k y^l q)(X, Y) \quad (k, l \geq 0, k+l+m \leq d). \quad (2.4)$$

In particular, the degree d columns $X^d, \dots, X^n Y^{d-n}$ are recursively determined in terms of columns of strictly lower degree. Since, by (2.4), each column $X^{d-m-k} Y^{m+k}$ ($0 \leq k \leq d-m$) may be expressed as a linear combination of columns to its left, it follows that if $n \leq d-m+1$, then M_d is flat. Since a flat positive moment matrix admits a unique positive, recursively generated extension (cf. [CF2]), we may assume that not every column of degree d is recursively determined, i.e., $n > d-m+1$, or

$$n + m > d + 1. \quad (2.5)$$

We may denote

$$X^n = p(X, Y) \equiv \sum_{r,s \geq 0, r+s \leq n-1} a_{rs} X^r Y^s \quad (2.6)$$

and

$$\begin{aligned} Y^m &= q(X, Y) \equiv \sum_{u,v \geq 0, u+v \leq m, v < m} b_{uv} X^u Y^v \\ &\equiv \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} X^a Y^b + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} X^c Y^e. \end{aligned} \quad (2.7)$$

Thus, in any positive and recursively generated (or merely RG) extension M_{d+1} , certain columns of $B(d+1)$ are recursively determined. On the left of $B(d+1)$, there is a band of columns,

$$\begin{aligned} X^{n+f} Y^{d+1-n-f} &:= (x^f y^{d+1-n-f} p)(X, Y) \\ &\equiv \sum_{r,s \geq 0, r+s \leq n-1} a_{rs} X^{r+f} Y^{s+d+1-n-f} \quad (0 \leq f \leq d+1-n) \end{aligned} \quad (2.8)$$

each of which is well-defined as a linear combination of columns of M_d . On the right of $B(d+1)$ there is another band of recursively determined columns,

$$\begin{aligned} X^{d+1-m-g} Y^{m+g} &:= (x^{d+1-m-g} y^g q)(X, Y) \\ &\equiv \sum_{u,v \geq 0, u+v \leq m, v < m} b_{uv} X^{u+d+1-m-g} Y^{v+g} \quad (0 \leq g \leq d+1-m). \end{aligned} \quad (2.9)$$

If $\deg q = m$, the sum in (2.9) may involve columns from the middle band, $X^{u+d+1-m-g} Y^{v+g}$ ($u+v = m, u+d+1-m-n < g < m-v$), which has not yet been defined, so some care is needed in implementing (2.9).

The proof of Theorem 2.1 entails two main steps, which we prove in detail in Section 4: the construction of the block $B(d+1)$, and the verification of the inclusion $\text{Ran } B(d+1) \subseteq \text{Ran } M_d$. Assuming that we have already built a unique block $B(d+1)$ consistent with the existence of a positive, recursively generated extension

$$M_{d+1} \equiv \begin{pmatrix} M_d & B(d+1) \\ B(d+1)^T & C(d+1) \end{pmatrix},$$

we next use this to construct a unique block $C(d+1)$ consistent with the existence of an RG extension.

Corollary 2.2. *If M_d satisfies the hypotheses of Theorem 2.1, then there exists a unique moment matrix block $C \equiv C(n+1)$ consistent with the structure of an RG extension M_{d+1} .*

Proof. In any RG extension M_{d+1} the column relations (2.8) and (2.9) must hold. The proof of Theorem 2.1 shows that these relations define a unique moment matrix block $B \equiv B(d+1)$ consistent with positivity and recursiveness. To define $C \equiv C(n+1)$, we may formally repeat the proof of Theorem 2.1 concerning the well-definedness and uniqueness of block $B(d+1)$, but applying the argument with M_d replaced with $B(d+1)^T$, and $B[d, d+1]$ replaced with $C \equiv B[d+1, d+1]$. In brief, we use $B(d+1)^T$ and (2.8) to define the left recursive band in C . We then define column $X^{d+1-m}Y^m$ by applying (2.9) successively, starting in row X^{n+1} , so that this column is Hankel with respect to the central band, which we are completing simultaneously. We then use (2.9) to successively define the remaining columns on the right. Lemma 4.3 can be used to show that the left band is internally Hankel, and an adaptation of the argument in Lemma 4.4 can be used to show that column $X^{d+1-m}Y^m$ is Hankel with respect to the left and central blocks. Finally, the argument of Lemma 4.5 can be adapted to show that the right band is also Hankel. \square

By combining Theorem 2.1 with Corollary 2.2, we immediately obtain the first of our main results, which follows.

Theorem 2.3. *If M_d is positive, with column relations generated entirely by (2.1) and (2.2) via recursiveness and linearity, then M_d admits a unique RG extension M_{d+1} , i.e., $\text{Ran } B(n+1) \subseteq \text{Ran } M_d$, (2.8)-(2.9) hold in $\text{Col } M_{d+1}$, and M_{d+1} is recursively generated.*

Corollary 2.4. *If M_d satisfies the hypotheses of Theorem 2.3 and $d = n + m - 2$, then M_d admits a flat moment matrix extension M_{d+1} (and β admits a rank M_d -atomic representing measure).*

Proof. Each column in the left band is, from (2.8), a linear combination of columns of strictly lower degree. Since $d = n + m - 2$, there is no central band in the construction of $B(d+1)$ in Theorem 2.1 and of $C(d+1)$ in Corollary 2.2. It thus follows from (2.9) that each column in the right band is also a linear combination of columns of strictly lower degree, so M_{d+1} is a flat extension. \square

To illustrate Corollary 2.4 in the simplest case, let $n = m = d = 2$ and suppose that M_2 satisfies the hypotheses of Theorem 2.3. It follows from [CF4] that M_2 admits a representing measure if and only if the equations $x^2 - p(x, y) = 0$ and $y^2 - q(x, y) = 0$ have at least 4 common real zeros. Corollary 2.4 implies that the latter ‘‘variety condition’’ is superfluous; indeed, from Corollary 2.4, there *is* a representing measure, so [CF4] implies that the system *must* have at least 4 ($= \text{rank } M_2$) common real zeros.

Note that if $M_d(\beta)$ satisfies the hypothesis of Theorem 2.3, then the existence or nonexistence of a representing measure for β will be established in at most $d - 1$ extension steps (after which the central band would vanish and every column of M_{2d-1} would be recursively determined). The next result shows that for every $d \geq 2$, there exists $M_d(\beta)$, satisfying the conditions of Theorem 2.3, for which the determination that a representing measure exists entails the maximum number of extension steps, each of which falls within the scope of Theorem 2.3.

Theorem 2.5. *For $d \geq 1$, there exists a moment matrix M_d , satisfying the conditions of Theorem 2.3, for which the extension algorithm determines successive positive, recursively generated extensions M_{d+1}, \dots, M_{2d-1} , and for which the first flat extension occurs at M_{2d-1} . Moreover, each extension M_{d+i} satisfies the conditions of Theorem 2.3, so to continue the sequence it is only necessary to verify that the RG extension M_{d+i+1} is positive semidefinite.*

Remark 2.6. To illustrate the significance of Theorem 2.5, let us compare it to the following result of [CF6, Theorem 1.2]: If $M_d(\beta)$ is a bivariate moment matrix with a column relation $p(X, Y) = 0$ (*deg* $p \leq 2$), then β has a measure if and only if M_d is positive, recursively generated, and $\text{rank } M_d \leq \text{card } \mathcal{V}(M_d)$. In this result, we see that the existence of a measure can be determined directly from the data by establishing the positivity, rank, and variety of M_d . By contrast, in Theorem 2.5 we see that it may be necessary to extend M_d to M_{2d-1} in order to establish that a measure exists. In this sense, within the framework of moment matrices, we see that the general case of the truncated moment problem cannot be solved in “closed form.” We may therefore seek to go beyond the framework of moment matrices. Recall that for $\beta \equiv \beta^{(2d)}$, L_β is *positive* if $p \in \mathcal{P}_{2d}$, $p|_{\mathbb{R}^d} \geq 0 \implies L_\beta(p) \geq 0$. In [CF7] we showed that β admits a representing measure if and only if L_β admits a positive extension $L : \mathcal{P}_{2d+2} \rightarrow \mathbb{R}$. Thus, as an alternative to constructing all of the extensions M_{d+1}, \dots, M_{2d+1} , in principle it would suffice to test the positivity of the Riesz functional corresponding to M_{d+1} . Unfortunately, at present there is no known concrete test for positivity of Riesz functionals (except in special cases, cf. [CF7], [FN1], [FN2]), so the moment matrix extension algorithm remains the most viable approach to resolving the existence of a representing measure in the bivariate *RD* case.

For the proof of Theorem 2.5, we require some preliminaries. For $d \geq 1$, suppose x_1, \dots, x_d are distinct and y_1, \dots, y_d are distinct. Let $P(x, y) := (x - x_1) \cdots (x - x_d)$, $Q(x, y) := (y - y_1) \cdots (y - y_d)$, and set $\mathcal{Z}_{P,Q} := \{(x_i, y_j)\}_{1 \leq i, j \leq d}$, the common zeros of P and Q . Let J be an ideal in $\mathbb{R}[x, y]$ with real variety $\mathcal{V} \equiv \mathcal{V}(J) := \{(x, y) \in \mathbb{R}^2 : s(x, y) = 0 \ \forall s \in J\}$. Let $I(\mathcal{V}) = \{f \in \mathbb{R}[x, y] : f|_{\mathcal{V}} \equiv 0\}$. In general, $I(\mathcal{V}(J))$ may be strictly larger than J [CLO]. However, for $J := (P, Q)$ (with P and Q as above), we will show below (Proposition 2.11) that each element of $I(\mathcal{V}(J))$ admits a “degree-bounded” representation which displays it as a member of J ; in particular, J is a *real ideal* in the sense of [M]. Although this result may well be known, we could not find a reference, so we include a proof for the sake of completeness. First, we need three auxiliary results.

Lemma 2.7. *(The Division Algorithm in $\mathbb{R}[x_1, \dots, x_n]$ [CLO, Section 2.3, Theorem 3]) Fix a monomial order $>$ on $\mathbb{Z}_{\geq 0}^n$ and let $F = (f_1, \dots, f_s)$ be an ordered s -tuple of polynomials in $\mathbb{R}[x_1, \dots, x_n]$. Then every $f \in \mathbb{R}[x_1, \dots, x_n]$ can be written as*

$$f = a_1 f_1 + \cdots + a_s f_s + r,$$

where $a_i \in \mathbb{R}[x_1, \dots, x_n]$, and either $r = 0$ or r is a linear combination, with coefficients in \mathbb{R} , of monomials, none of which is divisible by any of the leading terms in f_1, \dots, f_s .

Furthermore, if $a_i f_i \neq 0$, then we have $\text{multideg}(f) \geq \text{multideg}(a_i f_i)$.

Lemma 2.8. [S, p. 67] *For $N \geq 1$ let v_1, \dots, v_N be distinct points in \mathbb{R}^2 , and consider the multivariable Vandermonde matrix $V_N := (v_i^\alpha)_{1 \leq i \leq N, \alpha \in \mathbb{Z}_+^2, |\alpha| \leq N-1}$, of size $N \times \frac{N(N+1)}{2}$. Then the rank of V_N equals N .*

Corollary 2.9. Let $\mathbf{x} \equiv \{x_1, \dots, x_m\}$ and $\mathbf{y} \equiv \{y_1, \dots, y_n\}$ be sets of distinct real numbers, and consider the grid $\mathbf{x} \times \mathbf{y} := \{(x_i, y_j)\}_{1 \leq i \leq m, 1 \leq j \leq n}$ consisting of $N := mn$ distinct points in \mathbb{R}^2 . Then the generalized Vandermonde matrix $V_{\mathbf{x} \times \mathbf{y}}$, obtained from V_N by removing all columns indexed by monomials divisible by x^m or y^n , is invertible.

Proof. The columns of V_N are indexed by the monomials in x and y of degree at most N , listed in degree-lexicographic order. The size of V_N is $N \times \frac{N(N+1)}{2}$, and by Lemma 2.8 we know that its rank is N . We will show that $V_{\mathbf{x} \times \mathbf{y}}$ has exactly N columns, and that each column that was removed from V_N to produce $V_{\mathbf{x} \times \mathbf{y}}$ is a linear combination of other columns in V_N . Toward the first assertion, assume without loss of generality that $m \leq n$, let $k := n - m$ (so that $m + k = n$), and observe that the columns of $V_{\mathbf{x} \times \mathbf{y}}$ are indexed by the following monomials:

$$\begin{aligned} & 1, \\ & \quad x, y, \\ & \quad \quad x^2, xy, y^2, \\ & \quad \quad \quad \dots, x^{m-1}, \dots, y^{m-1}, \\ & \quad \quad \quad x^{m-1}y, \dots, xy^{m-1}, y^m, \\ & \quad \quad \quad x^{m-1}y^2, \dots, xy^m, y^{m+1}, \\ & \quad \quad \quad x^{m-1}y^3, \dots, xy^{m+1}, y^{m+2}, \\ & \quad \quad \quad \dots, \\ & \quad \quad \quad x^{m-1}y^k, \dots, xy^{m+k-2}, y^{m+k-1}, \\ & \quad \quad \quad x^{m-1}y^{k+1}, \dots, xy^{m+k-2} \\ & \quad \quad \quad \dots, \\ & \quad \quad \quad x^{m-1}y^{n-1}. \end{aligned}$$

The number of monomials is then $(1 + 2 + \dots + m) + mk + [(m-1) + (m-2) + \dots + 2 + 1] = \frac{m(m+1)}{2} + mk + \frac{(m-1)m}{2} = m^2 + mk = m(m+k) = mn$. It follows that $V_{\mathbf{x} \times \mathbf{y}}$ has exactly $N \equiv mn$ columns.

To prove the second assertion, observe that the polynomials $P := (x - x_1) \cdots (x - x_m)$ and $Q := (y - y_1) \cdots (y - y_n)$ vanish identically on $\mathbf{x} \times \mathbf{y}$, and therefore the columns of V_N indexed by multiples of x^m or y^n are linear combinations of columns preceding them in degree-lexicographic order.

By combining the preceding two assertions, it follows that $V_{\mathbf{x} \times \mathbf{y}}$, having size N and rank N , must be invertible. \square

The following result is a special case of Alon's Combinatorial Nullstellensatz [A]; for completeness, we give a proof based on Corollary 2.9.

Corollary 2.10. Let $G \equiv \mathbf{x} \times \mathbf{y}$ be a grid as in Corollary 2.9, let $N := mn$, and let $p \in \mathbb{R}[x, y]$ be such that $\deg_x p < m$ and $\deg_y p < n$. Assume also that $p|_G \equiv 0$. Then $p \equiv 0$.

Proof. We wish to apply Corollary 2.9. From the hypotheses, it is straightforward to verify that p does not contain any monomials divisible by x^m or y^n , so \hat{p} , properly extended with

zeros to indicate the absence of relevant monomials, can be regarded as a vector in \mathbb{R}^N , the domain of the generalized Vandermonde matrix V_G in Corollary 2.9. Since, by assumption, $p(x_i, y_j) = 0$ for all $1 \leq i \leq m$ and $1 \leq j \leq n$, it follows that $V_G \hat{p} = 0$. Since V_G is invertible (by Corollary 2.9), we must have $\hat{p} = 0$, so $p \equiv 0$, as desired. \square

Proposition 2.11. *Let $P(x, y) := (x - x_1) \cdots (x - x_d)$ and let $Q(x, y) := (y - y_1) \cdots (y - y_d)$. If $\rho := \text{multideg}(f) \geq d$ and $f|_{\mathcal{V}((P, Q))} \equiv 0$, then there exists $u, v \in \mathcal{P}_{\rho-d}$ such that $f = uP + vQ$.*

Proof. Let $\mathcal{V} := \mathcal{V}((P, Q))$. By Lemma 2.7, we can write $f = uP + vQ + r$, where $\text{multideg}(uP) \leq \rho$ and $\text{multideg}(vQ) \leq \rho$. It follows that $u, v \in \mathcal{P}_{\rho-d}$ and that $r|_{\mathcal{V}} \equiv 0$. Moreover, r is a linear combination, with coefficients in \mathbb{R} , of monomials, none of which is divisible by any of the leading terms in P and Q , that is, they are not divisible by x^d and y^d . Therefore, r satisfies the hypotheses of Corollary 2.10 with $m = n = d$. By Corollary 2.10, $r \equiv 0$. Thus, $f = uP + vQ$, as desired. \square

Proof of Theorem 2.5. At several points of the proof we will use the fact that if a moment matrix M_k admits a representing measure ν and $f \in \mathcal{P}_k$, then $f|_{\text{supp } \nu} \equiv 0$ if and only if $f(X, Y) = 0$ in \mathcal{C}_{M_k} [CF2, Proposition 3.1]. Let x_1, \dots, x_d and y_1, \dots, y_d be sets of distinct real numbers, and let $G := \mathbf{x} \times \mathbf{y} \equiv (x_i, y_j)_{1 \leq i, j \leq d}$ denote the corresponding grid. Let μ denote a measure whose support is precisely equal to G and let $M_d := M_d[\mu]$. Let $P(x, y) := (x - x_1) \cdots (x - x_d)$ and let $Q(x, y) := (y - y_1) \cdots (y - y_d)$. Since $P|_G \equiv 0$ and $Q|_G \equiv 0$, then $P(X, Y) = 0$ and $Q(X, Y) = 0$ in \mathcal{C}_{M_d} , whence $X^d = p(X)$ and $Y^d = q(Y)$ for certain $p, q \in \mathcal{P}_{d-1}$ satisfying $P(x, y) \equiv x^d - p(x)$ and $Q(x, y) \equiv y^d - q(y)$; thus, M_d is recursively determinate. We first show that the only column dependence relations in M_d arise from the above relations via linearity, so that M_d falls within the scope of Theorem 2.3. If $\deg f = d$ and $f(X, Y) = 0$ in $\text{Col } M_d$, then $f|_G \equiv 0$, so Proposition 2.11 implies that there exists scalars u and v such that $f = uP + vQ$. Thus, $f(X, Y) = uP(X, Y) + vQ(X, Y)$. Further, if $\deg f < d$ and $f(X, Y) = 0$, then since $f|_G \equiv 0$, it follows from Corollary 2.10 that $f \equiv 0$ (whence $M_{d-1} \succ 0$). Thus, M_d satisfies the conditions of Theorem 2.3.

Since M_d has the finitely atomic representing measure μ , M_d admits successive positive, recursively generated extensions $M_{d+1}[\mu], M_{d+2}[\mu], \dots$, so clearly these are the unique successive positive, recursively determined extensions of M_d ; let $M_{d+k} := M_{d+k}[\mu]$ ($1 \leq k \leq d-1$). We seek to show that each of M_{d+1}, \dots, M_{2d-1} falls within the scope of Theorem 2.3 and that the first flat extension in this sequence occurs with $\text{rank } M_{2d-1} = \text{rank } M_{2d-2}$. We first give a concrete description of $\ker M_{d+k}$. Since $M_{d-1} \succ 0$, if $r \in \mathcal{P}_{d+k}$ with $\hat{r} \in \ker M_{d+k}$, then $\deg r = d + j$ for $0 \leq j \leq k$. Since μ is a representing measure for M_{d+k} , it follows that $r|_{\text{supp } \mu} \equiv 0$. Proposition 2.11 now implies that there exist $u, v \in \mathcal{P}_j$ such that $r = uP + vQ$ (with P and Q defined above in the description of μ). Thus $\ker M_{d+k}$ is indexed by the recursively determined columns; precisely, $\ker M_{d+k}$ is the span of all of the columns $x^s y^t (\widehat{x^d - p})$ and $x^s y^t (\widehat{y^d - q})$ ($s, t \geq 0, s + t \leq k$). Thus, M_{d+k} satisfies the conditions of Theorem 2.3. In passing from M_{d+k-1} to M_{d+k} there are $d + k + 1$ new columns, of which $2(k + 1)$ are recursively determined, and since these correspond (as just above) to elements of $\ker M_{d+k}$, we have $\text{rank } M_{d+k} = \text{rank } M_{d+k-1} + (d + k + 1) - 2(k + 1) = \text{rank } M_{d+k-1} + d - k - 1$. Thus the first flat extension occurs when $k = d - 1$, in passing from M_{2d-2} to M_{2d-1} . \square

We continue with an example which shows that Theorem 2.1 is no longer valid if we permit column dependence relations in M_d in addition to those in (2.3) - (2.4).

Example 2.12. We define M_3 by setting $\beta_{00} = \beta_{20} = \beta_{02} = 1$; $\beta_{11} = \beta_{30} = \beta_{21} = \beta_{03} = 0$; $\beta_{12} = \beta_{40} = 2$; $\beta_{31} = \beta_{13} = 0$; $\beta_{22} = 5$, $\beta_{04} = 22$; $\beta_{50} = -1$, $\beta_{41} = -2$, $\beta_{32} = 13$, $\beta_{23} = 3$, $\beta_{14} = \frac{894}{13}$, $\beta_{05} = \frac{336}{13}$; $\beta_{60} = 178$, $\beta_{51} = 139$, $\beta_{42} = 159$, $\beta_{33} = \frac{1657}{13}$, $\beta_{24} = \frac{4298}{13}$, $\beta_{15} = r$, $\beta_{06} = \gamma \equiv \frac{443272376768 - 2742712830r - 4826809r^2}{41327767}$. Thus, we have

$$M_3 = \begin{pmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 2 & 0 \\ 0 & 1 & 0 & 0 & 0 & 2 & 2 & 0 & 5 & 0 \\ 0 & 0 & 1 & 0 & 2 & 0 & 0 & 5 & 0 & 22 \\ 1 & 0 & 0 & 2 & 0 & 5 & -1 & -2 & 13 & 3 \\ 0 & 0 & 2 & 0 & 5 & 0 & -2 & 13 & 3 & \frac{894}{13} \\ 1 & 2 & 0 & 5 & 0 & 22 & 13 & 3 & \frac{894}{13} & \frac{336}{13} \\ 0 & 2 & 0 & -1 & -2 & 13 & 178 & 139 & 159 & \frac{1657}{13} \\ 0 & 0 & 5 & -2 & 13 & 3 & 139 & 159 & \frac{1657}{13} & \frac{4298}{13} \\ 2 & 5 & 0 & 13 & 3 & \frac{894}{13} & 159 & \frac{1657}{13} & \frac{4298}{13} & r \\ 0 & 0 & 22 & 3 & \frac{894}{13} & \frac{336}{13} & \frac{1657}{13} & \frac{4298}{13} & r & \gamma \end{pmatrix}. \quad (2.10)$$

It is straightforward to check that M_3 is positive, recursively generated, and recursively determinate, with $M_2 \succ 0$, $\text{rank } M_3 = 7$ and column dependence relations

$$X^3 = p(X, Y) := 40 \cdot 1 - 24X + 4Y - 53X^2 - 2XY + 13Y^2, \quad (2.11)$$

$$X^2Y = t(X, Y) := 35 \cdot 1 - 22X - Y - 46X^2 + 3XY + 11Y^2, \quad (2.12)$$

and

$$Y^3 = q(X, Y) := d_1 \cdot 1 + d_2X + d_3Y + d_4X^2 + d_5XY + d_6Y^2 + d_7XY^2, \quad (2.13)$$

where $d_1 = \frac{3(487658 - 1651r)}{1447}$, $d_2 = \frac{3(-342075 + 1157r)}{1447}$, $d_3 = \frac{2(-2131598 + 6591r)}{18811}$, $d_4 = \frac{-2000094 + 6773r}{1447}$, $d_5 = \frac{2338519 - 6591r}{18811}$, $d_6 = \frac{2(-316575 + 1079r)}{1447}$, $d_7 = \frac{-48015 + 169r}{1447}$. Thus, M_3 satisfies all of the hypotheses of Theorem 2.1, except that (2.12) is an ‘‘extra’’ dependence relation (not a linear combination of the relations defined in (2.11) and (2.13)). We claim that M_3 does not admit a moment matrix extension block $B(4)$ such that $(M_3 \ B(4))$ is recursively generated. Indeed, if such a block existed, then in the column space of $(M_3 \ B(4))$ we would have $X^3Y = (yp)(XY) := 40Y - 24XY + 4Y^2 - 53X^2Y - 2XY^2 + 13Y^3$ and also $X^3Y = (xt)(X, Y) := 35X - 22X^2 - XY - 46X^3 + 3X^2Y + 11XY^2$. A calculation shows that $\langle (yp)(X, Y) - (xt)(X, Y), XY^2 \rangle = \frac{-49462 + 169r}{13}$, so for $r \neq \frac{49462}{169}$, X^3Y is not well-defined. Thus, the conclusions of Theorem 2.1 do not hold for M_3 (and thus there is no representing measure). \square

By contrast with the preceding example, we next show that if $M_d \in RD$, with all column dependence relations of strictly lower degree, then M_d does admit an RG extension.

Theorem 2.13. *Suppose M_d is positive and recursively generated, and satisfies (2.3)-(2.4). If each column relation in M_d can be expressed as $X^iY^j = r(X, Y)$ with $\text{deg } r < i + j$, then M_d admits a unique RG extension.*

We present the proof of Theorem 2.13 in Section 5. Finally, we note that in applying the algorithm, Theorem 2.3 or Theorem 2.13 may apply at some extension steps, but not at others. Consider [F1, Example 4.15], which concerns a recursively determinate M_5 with $n = m = d = 5$, $\deg p = 5$, $\deg q = 4$. The moment matrix M_5 satisfies the hypotheses of Theorem 2.3 (with the roles of p and q reversed). The RG extension M_6 is positive semidefinite and satisfies the hypotheses of Theorem 2.3. The RG extension M_7 is also positive semidefinite, but has a new column relation, $X^3Y^4 = r(X, Y)$ ($\deg r = 6$), that is not recursively determined from $X^5 = p(X, Y)$ or $Y^5 = q(X, Y)$. Thus, Theorem 2.3 does not apply to M_7 , nor does Theorem 2.13 (since $\deg p = 5 = n$). Nevertheless, in this case, when the algorithm is applied to M_7 , a flat extension M_8 (and a measure) results.

3 An extension sequence that fails at the second stage

Recall that in the most important case of recursive determinacy, a positive, flat M_d admits unique positive, recursively generated extensions of all orders, $M_{d+1}, \dots, M_{d+k}, \dots$, leading to a unique representing measure. Further, in all of the examples of [CF3], [CFM] and [F1], when a positive, recursively generated, recursively determinate M_d fails to have a representing measure, it is because it fails to admit a positive, recursively generated extension M_{d+1} . These results suggest the question as to whether a positive, recursively generated, recursively determinate M_d which admits a positive, recursively generated M_{d+1} necessarily admits positive, recursively generated extensions of all orders (and thus a representing measure) [F1, Question 4.19]. In this section we provide a negative answer to this question. In the sequel we construct a positive, recursively generated, recursively determinate $M_4(\beta^{(8)})$ which admits a positive, recursively generated extension M_5 , but such that M_5 fails to admit a positive, recursively generated extension M_6 . It then follows from the Bayer-Teichmann Theorem that $\beta^{(8)}$ has no representing measure.

We define M_4 by defining its component blocks in the decomposition

$$M_4 = \begin{pmatrix} M_3 & B(4) \\ B(4)^T & C(4) \end{pmatrix}. \quad (3.1)$$

We begin by setting $\beta_{00} = \beta_{20} = \beta_{02} = \beta_{22} = 1$, $\beta_{40} = \beta_{04} = \beta_{42} = \beta_{24} = 2$, $\beta_{60} = \beta_{06} = 5$, and all other moments up to degree 6 set to 0, so that

$$M_3 = \begin{pmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 2 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 2 \\ 1 & 0 & 0 & 2 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 2 & 0 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 & 0 & 5 & 0 & 2 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 2 & 0 & 2 \\ 0 & 1 & 0 & 0 & 0 & 0 & 2 & 0 & 2 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 & 0 & 2 & 0 & 5 \end{pmatrix}. \quad (3.2)$$

We next set

$$B(4) = \begin{pmatrix} 2 & 0 & 1 & 0 & 2 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 5 & 0 & 2 & 0 & 2 \\ 0 & 2 & 0 & 2 & 0 \\ 2 & 0 & 2 & 0 & 5 \\ a & b & 0 & 0 & 0 \\ b & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & g \\ 0 & 0 & 0 & g & h \end{pmatrix}, \quad (3.3)$$

where $\beta_{70} = a$, $\beta_{61} = b$, $\beta_{16} = g$, $\beta_{07} = h$, and all other degree 7 moments equal 0. Let

$$p(x, y) := ax^3 + bx^2y + 3x^2 - by - 2ax - 1 \quad (3.4)$$

and

$$q(x, y) := gxy^2 + hy^3 + 3y^2 - 2hy - gx - 1, \quad (3.5)$$

so that in the column space of $(M_3 \ B(4))$, we have the relations

$$X^4 = p(X, Y) \quad (3.6)$$

and

$$Y^4 = q(X, Y), \quad (3.7)$$

and $\text{rank} \begin{pmatrix} M_3 & B(4) \end{pmatrix} = 13$.

We complete the definition of a recursively determinate M_4 by extending the relations (3.6) and (3.7) to the columns of $(B(4)^T \ C(4))$, leading to

$$C(4) = \begin{pmatrix} 13 + a^2 + b^2 & ab & 5 & 0 & 4 \\ ab & 5 & 0 & 4 & 0 \\ 5 & 0 & 4 & 0 & 5 \\ 0 & 4 & 0 & 5 & gh \\ 4 & 0 & 5 & gh & 13 + g^2 + h^2 \end{pmatrix}. \quad (3.8)$$

Since $M_3 \succ 0$ (positive and invertible), we see that $M_4 \succeq 0$ with rank 13 if and only if $\Delta(4) \equiv C(4) - B(4)^T M_3^{-1} B(4) \succ 0$. In view of (3.6) and (3.7), this is equivalent to the positivity of the compression of $\Delta(4)$ to rows and columns indexed by X^3Y , X^2Y^2 , XY^3 , i.e.,

$$[\Delta(4)]_{\{X^3Y, X^2Y^2, XY^3\}} \equiv \begin{pmatrix} 1 - b^2 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 - g^2 \end{pmatrix} \succ 0. \quad (3.9)$$

Thus, if b and g satisfy $1 - b^2 > 0$ and $1 - g^2 > 0$, then M_4 is positive, recursively generated, and recursively determinate, with $\text{rank} \ M_4 = 13$, so M_4 satisfies the hypotheses of Theorem 2.1.

We next seek to extend M_4 to a positive and recursively generated M_5 . In view of (3.6) and (3.7), this can only be accomplished by defining

$$X^5 := (xp)(X, Y) \quad (3.10)$$

and

$$Y^5 := (yq)(X, Y). \quad (3.11)$$

Theorem 2.1 implies that the resulting $B(5)$ is well-defined and satisfies $\text{Ran } B(5) \subseteq \text{Ran } M_4$, so there exists W satisfying $B(5) = M_4W$. A calculation now shows that if we define $C(5)$ via (3.10) and (3.11) (as we must to preserve recursiveness), then $M_5 \succeq 0$ if and only if

$$\Delta(5) \equiv C(5) - B(5)^T W = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{-1+2b^2}{-1+b^2} & bg & 0 & 0 \\ 0 & 0 & bg & \frac{-1+2g^2}{-1+g^2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \succeq 0.$$

Thus, using nested determinants, and since $b^2 < 1$, we see that M_5 is positive and recursively generated, with rank $M_5 = 15$ if and only if

$$b^2 < \frac{1}{2} \quad (3.12)$$

and

$$1 - 2b^2 - 2g^2 + 3b^2g^2 + b^4g^2 + b^2g^4 - b^4g^4 > 0. \quad (3.13)$$

For example, setting $b = g = \frac{1}{4}$, the expression in (3.13) equals $\frac{49951}{65536} (> 0)$, so it follows that M_5 is positive and recursively generated, with rank $M_5 = 15$, whence M_5 satisfies the conditions of Theorem 2.1.

With these values for b and g (or using other appropriate values), we next attempt to define a positive and recursively generated extension M_6 . This can only be done by defining $X^6 := (x^2p)(X, Y)$ and $Y^6 := (y^2q)(X, Y)$. Theorem 2.1 implies that the resulting $B(6)$ is well-defined and that there is a matrix V such that $B(6) = M_5V$. Further, $C(6)$ is uniquely defined via the preceding column relations. M_6 as thus defined is recursively generated (by construction), but we will show that it need not be positive. Indeed, a calculation shows that $\Delta(6) \equiv C(6) - B(6)^T V$ is identically 0 except perhaps for the element in the row and column indexed by X^3Y^3 (the row 4, column 4 element), which is equal to

$$\frac{(1 - 3b^2 + b^4 - ab^2g + ab^4g + bh - 2b^3h)(-1 - ag + 3g^2 + 2ag^3 - g^4 + bg^2h - bg^4h)}{-1 + 2b^2 + 2g^2 - 3b^2g^2 - b^4g^2 - b^2g^4 + b^4g^4}.$$

Note that the denominator of the preceding expression is the negative of the expression in (3.13), and is thus strictly negative. Thus M_6 is positive if and only if

$$\eta := (1 - 3b^2 + b^4 - ab^2g + ab^4g + bh - 2b^3h)(-1 - ag + 3g^2 + 2ag^3 - g^4 + bg^2h - bg^4h) \leq 0. \quad (3.14)$$

With $b = g = \frac{1}{4}$, we have

$$\eta = \frac{(-836 + 15a - 224h)(836 + 224a - 15h)}{1048576}.$$

If we choose a and h so that $\eta = 0$, then M_6 is a flat extension of M_5 , and $\beta \equiv \beta^{(8)}$ has a 15-atomic representing measure. If we choose a and h so that $\eta < 0$, then M_6 is positive with rank 16, and since, in Corollary 2.4, $n = m = 4$ and $d = 6$, it follows that M_6 has a flat extension M_7 . However, if we choose a and h so that $\eta > 0$ (e.g., with $h = 0$ and $a > \frac{836}{15}$), then M_6 is not positive, whence β has no representing measure.

4 Proof of Theorem 2.1

The proof of Theorem 2.1 entails two main steps: (i) the construction of the block $B(d+1)$ from the column relations (2.1) and (2.2) so that $(M_d B(d+1))$ is recursively generated; and (ii) the verification that $\text{Ran } B(d+1) \subseteq \text{Ran } M_d$.

STEP (i): Step (i) will follow from a series of five auxiliary results (Lemmas 4.1 - 4.5). To begin the formal definition of $B(d+1)$, note that blocks $B[0, d+1], \dots, B[d-1, d+1]$ are completely defined in terms of moments in M_d . Indeed, for $0 \leq i \leq d+1$, $0 \leq j \leq d-1$, and $h, k \geq 0$ with $h+k=j$, the component of $B[j, d+1]$ in row $X^h Y^k$ and column $X^i Y^{d+1-i}$, which we denote by $\langle X^i Y^{d+1-i}, X^h Y^k \rangle$, must equal $\beta_{i+h, d+1-i+k}$. Note also that for $i \geq n$, the above component is alternately defined by (2.8), so we must show that the two definitions agree.

Lemma 4.1. *For $0 \leq f \leq d+1-n$ and $i, j \geq 0$ with $i+j \leq d-1$, the entry in column $X^{n+f} Y^{d+1-n-f}$, row $X^i Y^j$, as defined by (2.8), coincides with the moment inherited from M_d by moment matrix structure, $\beta_{n+f+i, d-n-f+j+1}$.*

Proof. Consider first the case when $d-n-f \geq 0$. From (2.8), we have

$$\begin{aligned} X^{n+f} Y^{d+1-n-f} &:= (x^f y^{d+1-n-f} p)(X, Y) \\ &\equiv \sum_{r,s \geq 0, r+s \leq n-1} a_{rs} X^{r+f} Y^{s+d+1-n-f} \quad (0 \leq f \leq d+1-n), \end{aligned}$$

so

$$\langle X^{n+f} Y^{d+1-n-f}, X^i Y^j \rangle = \sum a_{rs} \langle X^{r+f} Y^{s+d+1-n-f}, X^i Y^j \rangle.$$

Since $r+f+s+d+1-n-f \leq d$, $s+d+1-n-f \geq 1$ and $i+j \leq d-1$, using the moment matrix structure of the blocks of M_d we may express the last sum as

$$\sum a_{rs} \langle X^{r+f} Y^{s+d-n-f}, X^i Y^{j+1} \rangle.$$

Now (2.3) implies that in M_d the later expression is equal to

$$\langle X^{n+f} Y^{d-n-f}, X^i Y^{j+1} \rangle = \beta_{n+f+i, d-n-f+j+1}.$$

For the case $f = d-n+1$ and $i+j \leq d-1$,

$$\begin{aligned} \langle X^{d+1}, X^i Y^j \rangle &= \sum a_{rs} \langle X^r Y^s X^{d+1-n}, X^i Y^j \rangle \\ &= \sum a_{rs} \langle X^r Y^s X^{d-n}, X^{i+1} Y^j \rangle \\ &= \langle X^d, X^{i+1} Y^j \rangle \\ &= \beta_{d+i+1, j}. \end{aligned}$$

□

We have just verified that in the left recursive band, in blocks of degree at most $d-1$, each column element coincides with the corresponding “old” moment from M_d . Old moments are also used to *define* the central (nonrecursive) band of columns in blocks of degree at most $d-1$. We next use these left and central bands, together with (2.9), to show that the column elements in the right recursive band, in blocks of degree at most $d-1$, also agree with corresponding old moments.

Lemma 4.2. For $0 \leq k \leq d+1-m$, $i, j \geq 0$, $i+j \leq d-1$, column $X^{d+1-m-k}Y^{m+k}$, as defined by (2.9), satisfies $\langle X^{d+1-m-k}Y^{m+k}, X^iY^j \rangle = \beta_{i+d+1-m-k, m+k+j}$.

Proof. The proof is by induction on k . For $k=0$, we show that $\langle X^{d+1-m}Y^m, X^iY^j \rangle = \beta_{i+d+1-m, m+j}$. From (2.9), we have

$$\begin{aligned} \langle X^{d+1-m}Y^m, X^iY^j \rangle &= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m+a}Y^b, X^iY^j \rangle \\ &+ \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{c+d+1-m}Y^e, X^iY^j \rangle. \end{aligned}$$

Since $a+b < m$, then in M_d ,

$$\langle X^{d+1-m+a}Y^b, X^iY^j \rangle = \beta_{d+1-m+a+i, b+j}.$$

Since $e < m$, $\langle X^{c+d+1-m}Y^e, X^iY^j \rangle$ is in either the left or central band, and thus equals the old moment $\beta_{c+d+1-m+i, e+j}$. Now

$$\langle X^{d+1-m}Y^m, X^iY^j \rangle = \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \beta_{d+1-m+a+i, b+j} + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \beta_{c+d+1-m+i, e+j}.$$

In M_d , the latter expression equals

$$\begin{aligned} \sum \alpha_{ab} \langle X^{d-m+a}Y^b, X^{i+1}Y^j \rangle + \sum \gamma_{ce} \langle X^{c+d-m}Y^e, X^{i+1}Y^j \rangle &= \langle X^{d-m}Y^m, X^{i+1}Y^j \rangle \\ &= \beta_{d-m+i+1, m+j}, \end{aligned}$$

as desired. We next assume the result is true for $0, \dots, k-1$. Consider first the case when $k < d+1-m$. We have

$$\begin{aligned} \langle X^{d+1-m-k}Y^{m+k}, X^iY^j \rangle &= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m-k+a}Y^{b+k}, X^iY^j \rangle \\ &+ \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{c+d+1-m-k}Y^{e+k}, X^iY^j \rangle. \end{aligned}$$

The term $\langle X^{d+1-m-k+a}Y^{b+k}, X^iY^j \rangle$ is a component of M_d , and thus equals the corresponding moment. Since $e+k \leq m+(k-1)$, $X^{c+d+1-m-k}Y^{e+k}$ is, by induction, a column for which the elements of row-degree $i+j$ are old moments. Thus,

$$\langle X^{d+1-m-k}Y^{m+k}, X^iY^j \rangle = \sum \alpha_{ab} \beta_{d+1-m-k+a+i, b+k+j} + \sum \gamma_{ce} \beta_{c+d+1-m-k+i, e+k+j}.$$

In M_d , the last expression equals

$$\begin{aligned} \sum \alpha_{ab} \langle X^{d+a-m-k}Y^{b+k}, X^{i+1}Y^j \rangle \\ + \sum \gamma_{ce} \langle X^{c+d-m-k}Y^{e+k}, X^{i+1}Y^j \rangle &= \langle X^{d-m-k}Y^{m+k}, X^{i+1}Y^j \rangle \\ &= \beta_{d-m-k+i+1, m+k+j}. \end{aligned}$$

Finally, we consider the case $k = d + 1 - m$. We have

$$\begin{aligned} \langle Y^{d+1}, X^i Y^j \rangle &= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^a Y^{b+d+1-m}, X^i Y^j \rangle \\ &\quad + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^c Y^{e+d+1-m}, X^i Y^j \rangle. \end{aligned}$$

Since $e < m$, then $c \geq 1$, so $X^c Y^{e+d+1-m}$ is to the left of Y^{d+1} , i.e., $c = d + 1 - m - k'$ for $k' = d + 1 - m - c < k$. Thus, by induction,

$$\langle Y^{d+1}, X^i Y^j \rangle = \sum \alpha_{ab} \beta_{a+i, b+d+1-m+j} + \sum \gamma_{ce} \beta_{c+i, e+d+1-m+j}.$$

In M_d , the last expression equals

$$\begin{aligned} \sum \alpha_{ab} \langle X^a Y^{b+d-m}, X^i Y^{j+1} \rangle + \sum \gamma_{ce} \langle X^c Y^{e+d-m}, X^i Y^{j+1} \rangle &= \langle Y^d, X^i Y^{j+1} \rangle \\ &= \beta_{i, d+j+1}, \end{aligned}$$

as desired. \square

To complete the definition of $B(d+1)$ we must define $B[d, d+1]$. Within this proposed block, we first use (2.8) to define the left recursive band, $X^{d+1}, \dots, X^n Y^{d+1-n}$. Note that between the end of the left band, $X^n Y^{d+1-n}$, and the beginning of the right band, $X^{n+1-m} Y^m$, there is a central band of $n + m - d - 2$ columns; set $\delta := n + m - d - 1$. In row X^d , each of the components in the central columns, $\langle X^{n-1} Y^{d+2-n}, X^d \rangle, \dots, \langle X^{d+2-m} Y^{m-1}, X^d \rangle$, corresponds via a cross-diagonal to a component of column $X^n Y^{d+1-n}$ (whose value is known from (2.8)), i.e.,

$$\langle X^{n-j} Y^{d+1-n+j}, X^d \rangle = \langle X^n Y^{d+1-n}, X^{d-j} Y^j \rangle \quad (1 \leq j \leq m + n - d - 2).$$

We may thus use (2.9) to define $\langle X^{d+1-m} Y^m, X^d \rangle$, and we extend the latter value along the central-band section of the cross-diagonal to which it belongs. Next, in row $X^{d-1} Y$, we use this value with (2.9) to define $\langle X^{d+1-m} Y^m, X^{d-1} Y \rangle$, and we extend this value along the central-band section of its cross-diagonal. Proceeding in this way, we completely define column $X^{d+1-m} Y^m$ and insure that it is Hankel with respect to the central band. Finally, we use (2.9) to define column $X^{d-m} Y^{m+1}$, and, successively, $X^{d-m-1} Y^{m+2}, \dots, Y^{d+1}$. This completes the definition of a proposed block $B[d, d+1]$. However, to ensure that it is well-defined as a moment block, we must check that for a cross-diagonal which intersects columns $X^n Y^{d+1-n}$ and $X^{d+1-m} Y^m$, the components of the cross-diagonal in these columns agree in value, i.e., the values arising from (2.8) are consistent with those arising from (2.9). More generally, we need to show that the block we have defined is constant on cross-diagonals.

To show that $B[d, d+1]$ is well-defined and Hankel, we begin with the following general result concerning adjacent columns that are recursively determined from the same column dependence relation. Suppose in $Col M_d$ there is a dependence relation $X^c Y^e = p(X, Y)$, where $c + e = d$ and $p(x, y) \equiv \sum_{a,b \geq 0, a+b \leq d-1} \alpha_{ab} x^a y^b \in \mathcal{P}_{d-1}$. Then the elements of $Col M_d$ defined by

$$X^{c+1} Y^e \equiv (xp)(X, Y) := \sum_{a,b \geq 0, a+b \leq d-1} \alpha_{ab} X^{a+1} Y^b$$

and

$$X^c Y^{e+1} \equiv (yp)(X, Y) := \sum_{a,b \geq 0, a+b \leq d-1} \alpha_{ab} X^a Y^{b+1}$$

are Hankel with respect to each other, as follows.

Lemma 4.3. For $i, j \geq 0$, $i + j \leq d$, $j > 0$,

$$\langle X^{c+1} Y^e, X^i Y^j \rangle = \langle X^c Y^{e+1}, X^{i+1} Y^{j-1} \rangle$$

Proof. We have

$$\langle X^{c+1} Y^e, X^i Y^j \rangle = \sum_{a,b \geq 0, a+b \leq d-1} \alpha_{ab} \langle X^{a+1} Y^b, X^i Y^j \rangle,$$

and since each row and column in the last sum has degree at most d , relative to M_d we may rewrite this sum as

$$\sum_{a,b \geq 0, a+b \leq d-1} \alpha_{ab} \langle X^a Y^{b+1}, X^{i+1} Y^{j-1} \rangle = \langle X^c Y^{e+1}, X^{i+1} Y^{j-1} \rangle.$$

This completes the proof. \square

It follows immediately from Lemma 4.3 that the left recursive band in $B[d, d+1]$ is constant on cross-diagonals. We next check that if an element of a column in the non-recursive central band can be reached on a cross-diagonal which intersects both columns $X^n Y^{d+1-n}$ (at the edge of the left recursive band) and $X^{d+1-m} Y^m$ (at the edge of the right recursive band), then the values obtained from both of these columns agree. This is the substance of the following lemma.

Lemma 4.4. For $0 \leq k \leq 2d + 1 - n - m$,

$$\langle X^{d+1-m} Y^m, X^{d-k} Y^k \rangle = \langle X^n Y^{d+1-n}, X^{d-\delta-k} Y^{\delta+k} \rangle. \quad (4.1)$$

Proof. The proof is by induction on k . We begin with the base case, $k = 0$, and seek to show that $\langle X^{d+1-m} Y^m, X^d \rangle = \langle X^n Y^{d+1-n}, X^{d-\delta} Y^\delta \rangle$ (recall that $\delta := n + m - d - 1$). Using (2.9), we may express $\langle X^{d+1-m} Y^m, X^d \rangle$ as

$$\sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m+a} Y^b, X^d \rangle + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{d+1-e} Y^e, X^d \rangle. \quad (4.2)$$

Note that $\langle X^{d+1-m+a} Y^b, X^d \rangle$ is a component of M_d ; further, since $e < m$, $\langle X^{c+d+1-m} Y^e, X^d \rangle$ is the endpoint of a cross-diagonal that lies entirely in the left and central bands, and is thus constant. Therefore, we may rewrite (4.2) as

$$\sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^d, X^{d+1-m+a} Y^b \rangle + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{d+1}, X^{d-e} Y^e \rangle$$

$$\begin{aligned}
&= \sum \alpha_{ab} \langle \sum a_{rs} X^{d-n+r} Y^s, X^{d+1-m+a} Y^b \rangle + \sum \gamma_{ce} \langle \sum a_{rs} X^{d-n+r+1} Y^s, X^{d-e} Y^e \rangle \\
&= \sum a_{rs} \sum \alpha_{ab} \langle X^{d-n+r} Y^s, X^{d+1-m+a} Y^b \rangle + \sum a_{rs} \sum \gamma_{ce} \langle X^{d-n+r+1} Y^s, X^{d-e} Y^e \rangle \\
&= \sum a_{rs} (\sum \alpha_{ab} \langle X^{d+1-m+a} Y^b, X^{d-n+r} Y^s \rangle + \sum \gamma_{ce} \langle X^{d-e} Y^e, X^{d-n+r+1} Y^s \rangle) \\
&= \sum a_{rs} \langle \sum \alpha_{ab} X^{d-m+a} Y^b + \sum \gamma_{ce} X^{d-e} Y^e, X^{d-n+r+1} Y^s \rangle \\
&= \sum a_{rs} \langle X^{d-m} Y^m, X^{d-n+r+1} Y^s \rangle.
\end{aligned}$$

Since $\delta = m + n - d - 1$, in M_d the last sum is equal to

$$\sum a_{rs} \langle X^r Y^{d+1-n+s}, X^{d-\delta} Y^\delta \rangle = \langle X^n Y^{d+1-n}, X^{d-\delta} Y^\delta \rangle,$$

which completes the proof of the base case.

We assume now that (4.1) holds for $0, \dots, k-1$, with $k-1 < 2d+1-n-m$. To establish (4.1) for k , we consider first the case $d-k \geq n$. Let us write $\kappa := \langle X^{d+1-m} Y^m, X^{d-k} Y^k \rangle$ as

$$\begin{aligned}
\kappa &= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m+a} Y^b, X^{d-k} Y^k \rangle \\
&+ \sum_{c,e \geq 0, c+e=m, e < m, d+1-e \geq n} \gamma_{ce} \langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle \\
&+ \sum_{c,e \geq 0, c+e=m, e < m, d+1-e \geq n, d+1-e < n} \gamma_{ce} \langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle. \tag{4.3}
\end{aligned}$$

Note that the components in the first sum of (4.3) lie in M_d . In the third sum, since $d+1-e < n$, column $X^{d+1-e} Y^e$ is in the middle band, and the component $\gamma := \langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle$ lies on a cross-diagonal σ strictly above the cross-diagonal for κ . Either because σ does not intersect column $X^{d+1-m} Y^m$, or by induction if it does, we see that γ has the same value as $\langle X^n Y^{d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle$ (on the same cross-diagonal). Thus (4.3) can

be expressed as

$$\begin{aligned}
\kappa &= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d-k} Y^k, X^{d+1-m+a} Y^b \rangle \\
&+ \sum_{c,e \geq 0, c+e=m, e < m, d+1-e \geq n} \gamma_{ce} \langle X^n X^{d+1-e-n} Y^e, X^{d-k} Y^k \rangle \\
&+ \sum_{c,e \geq 0, c+e=m, e < m, d+1-e \geq n, d+1-e < n} \gamma_{ce} \langle X^n Y^{d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle \\
&= \sum \alpha_{ab} \langle X^n X^{d-k-n} Y^k, X^{d+1-m+a} Y^b \rangle \\
&+ \sum \gamma_{ce} \langle X^n X^{d+1-e-n} Y^e, X^{d-k} Y^k \rangle \\
&+ \sum \gamma_{ce} \langle X^n Y^{d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle \\
&= \sum \alpha_{ab} \sum a_{rs} \langle X^{r+d-k-n} Y^{s+k}, X^{d+1-m+a} Y^b \rangle \\
&+ \sum \gamma_{ce} \sum a_{rs} \langle X^{r+d+1-e-n} Y^{s+e}, X^{d-k} Y^k \rangle \\
&+ \sum \gamma_{ce} \sum a_{rs} \langle X^r Y^{s+d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle \\
&= \sum a_{rs} \left(\sum \alpha_{ab} \langle X^{r+d-k-n} Y^{s+k}, X^{d+1-m+a} Y^b \rangle \right. \\
&+ \sum \gamma_{ce} \langle X^{r+d+1-e-n} Y^{s+e}, X^{d-k} Y^k \rangle \\
&+ \left. \sum \gamma_{ce} \langle X^r Y^{s+d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle \right). \tag{4.4}
\end{aligned}$$

Using the symmetry of M_d in the first and third inner sums of the last expression, we may rewrite this expression as

$$\begin{aligned}
&\sum a_{rs} \left(\sum \alpha_{ab} \langle X^{d+1-m+a} Y^b, X^{r+d-k-n} Y^{s+k} \rangle + \sum \gamma_{ce} \langle X^{r+d+1-e-n} Y^{s+e}, X^{d-k} Y^k \rangle \right. \\
&+ \left. \sum \gamma_{ce} \langle X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)}, X^r Y^{s+d+1-n} \rangle \right). \tag{4.5}
\end{aligned}$$

In the second inner sum of (4.5), $\langle X^{r+d+1-e-n} Y^{s+e}, X^{d-k} Y^k \rangle$ is a component of M_d and thus equals the moment $\beta_{r+d+1-e-n+d-k, s+e+k}$. Since $X^{d-k+r-n} Y^{s+k}$ is a row of degree at most $d-1$, this moment coincides with $\langle X^{c+d+1-m} Y^e, X^{d-k+r-n} Y^{s+k} \rangle$ from the left band of $B[d+r+s-n, n+1]$. Further, in the third inner sum of (4.5),

$$\langle X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)}, X^r Y^{s+d+1-n} \rangle$$

is also a component of M_d , equal to $\beta_{r+d-k-(n-(d+1-e)), s+k+n-(d+1-e)+d+1-n}$, and this moment coincides with $\langle X^{d+1-e} Y^e, X^{r+d-k-n} Y^{s+k} \rangle$ from the middle band in $B[d+r+s-n, d+1]$. Thus, the expression in (4.5) can be written as

$$\begin{aligned}
&\sum a_{rs} \left(\sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m+a} Y^b, X^{r+d-k-n} Y^{s+k} \rangle \right. \\
&+ \sum_{c,e \geq 0, c+e=m, e < m, d+1-e \geq n} \gamma_{ce} \langle X^{d+1-m+c} Y^e, X^{r+d-k-n} Y^{s+k} \rangle \\
&+ \left. \sum_{c,e \geq 0, c+e=m, e < m, d+1-e < n} \gamma_{ce} \langle X^{d+1-m+c} Y^e, X^{r+d-k-n} Y^{s+k} \rangle \right), \tag{4.6}
\end{aligned}$$

which equals

$$\sum a_{rs} \langle X^{d+1-m} Y^m, X^{r+d-k-n} Y^{s+k} \rangle. \quad (4.7)$$

Since $X^{r+d-k-n} Y^{s+k}$ is a row of degree at most $d-1$, Lemma 4.2 implies that the expression in (4.7) equals

$$\begin{aligned} \sum a_{rs} \beta_{d+1-m+r+d-k-n, m+s+k} &= \sum a_{rs} \beta_{r+d-\delta-k, s+d+1-n+\delta+k} \\ &= \sum a_{rs} \langle X^r Y^{s+d+1-n}, X^{d-\delta-k} Y^{\delta+k} \rangle \\ &= \langle X^n Y^{d+1-n}, X^{d-\delta-k} Y^{\delta+k} \rangle. \end{aligned}$$

This completes the proof of the induction step for (4.1) when $d-k \geq n$.

We next treat the case when $d-k < n$, which implies $\delta+k \geq m$. We have

$$\begin{aligned} \langle X^n Y^{d+1-n}, X^{d-\delta-k} Y^{\delta+k} \rangle &= \sum a_{rs} \langle X^r Y^{d+1-n+s}, X^{d-\delta-k} Y^{\delta+k} \rangle \\ &= \sum a_{rs} \langle X^{d-\delta-k} Y^m Y^{\delta+k-m}, X^r Y^{d+1-n+s} \rangle \\ &= \sum a_{rs} \left(\sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{a+d-\delta-k} Y^{b+\delta+k-m}, X^r Y^{d+1-n+s} \rangle \right. \\ &\quad \left. + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{c+d-\delta-k} Y^{e+\delta+k-m}, X^r Y^{d+1-n+s} \rangle \right). \end{aligned} \quad (4.8)$$

Note for future reference that all of the matrix components that appear in (4.8) come from M_d .

We now consider

$$\begin{aligned} \langle X^{d+1-m} Y^m, X^{d-k} Y^k \rangle &= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m+a} Y^b, X^{d-k} Y^k \rangle \\ &\quad + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{d+1-m+c} Y^e, X^{d-k} Y^k \rangle \\ &= \sum \alpha_{ab} \langle X^{d-k} Y^k, X^{d+1-m+a} Y^b \rangle \\ &\quad + \sum \gamma_{ce} \langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle \end{aligned} \quad (4.9)$$

(using symmetry of M_d in the first sum). Since $k - (n - (d - k))$, $d + 1 - m + a - n + d - k$ ($= a + d - \delta - k$), and $b + n - (d - k)$ are all nonnegative, by applying the block-Hankel property of M_d to the first sum in (4.9), we may rewrite the expression in (4.9) as

$$\begin{aligned} &\sum \alpha_{ab} \langle X^n Y^{k-(n-(d-k))}, X^{d+1-m+a-n+d-k} Y^{b+n-(d-k)} \rangle \\ &+ \sum \gamma_{ce} \langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle \end{aligned} \quad (4.10)$$

$$\begin{aligned} &= \sum \alpha_{ab} \sum a_{rs} \langle X^r Y^{s+k-(n-(d-k))}, X^{d+1-m+a-n+d-k} Y^{b+n-(d-k)} \rangle \\ &+ \sum \gamma_{ce} \langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle, \end{aligned} \quad (4.11)$$

and all of the matrix components in the first double sum of (4.11) are from M_d . Comparing the components in the first double sums of (4.8) and (4.11), we have

$$\begin{aligned}\langle X^{a+d-\delta-k}Y^{b+\delta+k-m}, X^rY^{d+1-n+s} \rangle &= \beta_{a+d-\delta-k+r, b+\delta+k-m+d+1-n+s} \\ &= \beta_{r+d+1-m+a-n+d-k, s+k-(n-(d-k))+b+n-(d-k)} \\ &= \langle X^rY^{s+k-(n-(d-k))}, X^{d+1-m+a-n+d-k}Y^{b+n-(d-k)} \rangle,\end{aligned}$$

so the first double sums of (4.8) and (4.11) are equal.

Let us write the rightmost sum in (4.11) as

$$\begin{aligned}&\sum_{c,e \geq 0, c+e=m, e < m, d+1-e \geq n} \gamma_{ce} \langle X^n X^{d+1-e-n} Y^e, X^{d-k} Y^k \rangle \\ &+ \sum_{c,e \geq 0, c+e=m, e < m, d+1-e < n} \gamma_{ce} \langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle.\end{aligned}\quad (4.12)$$

In the second sum of (4.12), since $d+1-e < n$, the component $\langle X^{d+1-e} Y^e, X^{d-k} Y^k \rangle$ (from the middle band) has the same value as the component $\langle X^n Y^{d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle$ on the same cross-diagonal. (This is because the cross-diagonal is strictly above that for $\langle X^{d+1-m} Y^m, X^{d-k} Y^k \rangle$, so the conclusion follows by definition or induction.) We may now write the expression in (4.12) as

$$\begin{aligned}&\sum_{c,e \geq 0, c+e=m, e < m, d+1-e \geq n} \gamma_{ce} \sum a_{rs} \langle X^{r+d+1-e-n} Y^{s+e}, X^{d-k} Y^k \rangle \\ &+ \sum_{c,e \geq 0, e < m, c+e=m, d+1-e < n} \gamma_{ce} \langle X^n Y^{d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle \\ = &\sum_{c,e \geq 0, e < m, c+e=m, d+1-e \geq n} \gamma_{ce} \sum a_{rs} \langle X^{r+d+1-e-n} Y^{s+e}, X^{d-k} Y^k \rangle \\ &+ \sum_{c,e \geq 0, c+e=m, d+1-e < n} \gamma_{ce} \sum a_{rs} \langle X^r Y^{s+d+1-n}, X^{d-k-(n-(d+1-e))} Y^{k+n-(d+1-e)} \rangle.\end{aligned}\quad (4.13)$$

All of the matrix components in (4.13) are from M_d , so (4.13) can be expressed as

$$\sum a_{rs} \sum_{c+e=m} \beta_{r+d+1-e-n+d-k, s+e+k}.$$

It is straightforward to check that this double sum coincides with the second double sum in (4.8) (whose matrix components also come entirely from M_d). This completes the proof that the second double sums in (4.8) and (4.11) have the same value, so the expressions in (4.8) and (4.11) are equal, which completes the proof of the induction when $d-k < n$. Thus, the induction is complete. \square

We have shown above that in $B[d, d+1]$ the columns $X^{d+1}, \dots, X^{d+1-m} Y^m$ are well-defined and Hankel with respect to one another. Using (2.9), we also successively defined columns $X^{d-m} Y^{m+1}, \dots, Y^{d+1}$. We next show that the columns $X^{d-m+1} Y^m, \dots, Y^{d+1}$ are Hankel with respect to each other, so that all of $B[d, d+1]$ has the Hankel property.

Lemma 4.5. For $0 \leq s \leq d+1-m$ and $i, j \geq 0$ with $i+j = d$ and $j > 0$, we have

$$\langle X^{d+1-m-s}Y^{m+s}, X^iY^j \rangle = \langle X^{d-m-s}Y^{m+s+1}, X^{i+1}Y^{j-1} \rangle.$$

Proof. The proof is by induction on $s \geq 0$. For $s = 0$, we have $\langle X^{d+1-m}Y^m, X^iY^j \rangle$

$$= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m+a}Y^b, X^iY^j \rangle + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{d+1-e}Y^e, X^iY^j \rangle. \quad (4.14)$$

In the first sum, each component is from M_d . In the second sum, column $X^{d+1-e}Y^e$ is strictly to the left of $X^{d+1-m}Y^m$, so it is Hankel with respect to its right successor, $X^{d-e}Y^{e+1}$. We may thus rewrite the expression in (4.14) as

$$\sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d-m+a}Y^{b+1}, X^{i+1}Y^{j-1} \rangle + \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{d-e}Y^{e+1}, X^{i+1}Y^{j-1} \rangle$$

$$= \langle X^{d-m}Y^{m+1}, X^{i+1}Y^{j-1} \rangle.$$

Assume now that the Hankel property holds through $s-1$ and consider

$$\begin{aligned} \langle X^{d+1-m-s}Y^{m+s}, X^iY^j \rangle &= \sum_{a,b \geq 0, a+b \leq m-1} \alpha_{ab} \langle X^{d+1-m+a-s}Y^{b+s}, X^iY^j \rangle \\ &+ \sum_{c,e \geq 0, c+e=m, e < m} \gamma_{ce} \langle X^{d+1-e-s}Y^{e+s}, X^iY^j \rangle. \end{aligned} \quad (4.15)$$

As above, in the first sum, each component is from M_d ; in the second sum, each column $X^{d+1-e-s}Y^{e+s}$ is to the left of $X^{d+1-m-s}Y^{m+s}$, so the Hankel property holds for this column by induction. We may thus write the expression in (4.15) as

$$\begin{aligned} \sum_{a,b} \alpha_{ab} \langle X^{d-m+a-s}Y^{b+s+1}, X^{i+1}Y^{j-1} \rangle + \sum_{c+e=m, e < m} \gamma_{ce} \langle X^{d-e-s}Y^{e+s+1}, X^{i+1}Y^{j-1} \rangle \\ = \langle X^{d-m-s}Y^{m+s+1}, X^{i+1}Y^{j-1} \rangle, \end{aligned}$$

which completes the proof by induction. \square

STEP (ii): The preceding results show that under the hypotheses of Theorem 2.1, there exists a unique block $B(d+1)$ that is consistent with recursiveness in $(M_d \ B(d+1))$. To prove Theorem 2.1, we must also show that $\text{Ran } B(d+1) \subseteq \text{Ran } M_d$. The following lemma is a step toward this end; it shows that the rows of $(M_d \ B(d+1))$ of the form $X^{n+f}Y^g$ ($f, g \geq 0$, $n+f+g \leq d$) are recursively determined from row X^n .

Lemma 4.6. For $i, j \geq 0, i+j \leq d+1$ and for $f, g \geq 0, n+f+g \leq d$,

$$\langle X^iY^j, X^{n+f}Y^g \rangle = \sum_{r,s \geq 0, r+s \leq n-1} a_{rs} \langle X^iY^j, X^{r+f}Y^{s+g} \rangle. \quad (4.16)$$

Proof. Since M_d is real symmetric, it follows from (2.8) that (4.16) holds for $i+j \leq d$. We may thus assume that $j = d+1-i$. Consider first the case when $n+f+g < d$. In the subcase when $i \leq d$, it follows from the presence of old moments in $B[n+f+g, d+1]$ that

$$\langle X^iY^{d+1-i}, X^{n+f}Y^g \rangle = \beta_{i+n+f, d+1-i+g},$$

and in M_d we have

$$\begin{aligned}
\beta_{i+n+f,d+1-i+g} &= \langle X^{n+f}Y^{g+1}, X^iY^{d-i} \rangle \\
&= \sum_{r,s \geq 0, r+s \leq n-1} a_{rs} \langle X^{r+f}Y^{s+g+1}, X^iY^{d-i} \rangle \\
&= \sum a_{rs} \langle X^iY^{d-i}, X^{r+f}Y^{s+g+1} \rangle \quad (\text{by symmetry in } M_d) \\
&= \sum a_{rs} \beta_{i+r+f, d-i+s+g+1} \\
&= \sum a_{rs} \langle X^iY^{d+1-i}, X^{r+f}Y^{s+g} \rangle \\
&\quad (\text{by moment matrix structure in } B(d+1)).
\end{aligned} \tag{4.17}$$

For the subcase when $i = d + 1$, we first note that $\langle X^{d+1}, X^{n+f}Y^g \rangle = \beta_{d+1+n+f,g} = \langle X^d, X^{n+f}Y^g \rangle = \langle X^{n+f}Y^g, X^d \rangle$, and we then proceed beginning as in (4.17).

We next consider the case $n + f + g = d$, and we seek to show that

$$\langle X^iY^{d+1-i}, X^{n+f}Y^g \rangle = \sum a_{rs} \langle X^iY^{d+1-i}, X^{r+f}Y^{s+g} \rangle. \tag{4.18}$$

We begin by showing that (4.18) holds if the column X^iY^{d+1-i} is recursively determined from (2.8), i.e., $i \geq n$. In this case, we have $0 \leq i \leq d + 1 - n$, so

$$\begin{aligned}
\langle X^iY^{d+1-i}, X^{n+f}Y^g \rangle &= \sum a_{rs} \langle X^rY^sX^{i-n}Y^{d+1-i}, X^{n+f}Y^g \rangle \\
&= \sum a_{rs} \langle X^{n+f}Y^g, X^rY^sX^{i-n}Y^{d+1-i} \rangle \\
&= \sum a_{uv} \sum a_{rs} \langle X^uY^vX^fY^g, X^rY^sX^{i-n}Y^{d+1-i} \rangle \\
&= \sum a_{uv} \sum a_{rs} \langle X^rY^sX^{i-n}Y^{d+1-i}, X^{u+f}Y^{v+g} \rangle \\
&= \sum a_{uv} \langle X^iY^{d+1-i}, X^{u+f}Y^{v+g} \rangle.
\end{aligned}$$

Thus

$$\langle X^iY^{d+1-i}, X^{n+f}Y^g \rangle = \sum a_{uv} \langle X^iY^{d+1-i}, X^{u+f}Y^{v+g} \rangle,$$

which is equivalent to (4.18).

Returning to the proof of (4.18), we next assume that column X^iY^{d+1-i} is not recursively determined, i.e., $d + 1 - m < i < n$. By the Hankel condition in $B(d + 1)$, we have

$$\begin{aligned}
\langle X^iY^{d+1-i}, X^{n+f}Y^g \rangle &= \langle X^nY^{d+1-n}, X^{i+f}Y^{n-i+g} \rangle \\
&= \sum a_{rs} \langle X^rY^{s+d+1-n}, X^{i+f}Y^{n-i+g} \rangle \\
&= \sum a_{rs} \beta_{r+i+f, s+d+1-i+g} \quad (\text{in } M_d) \\
&= \sum a_{rs} \langle X^iY^{d+1-i}, X^{r+f}Y^{s+g} \rangle \\
&\quad (\text{since } r + f + s + g < n + f + g = d).
\end{aligned}$$

Note that if (4.16) holds for a collection of columns, then it holds for linear combinations of those columns. Thus, using the preceding cases and (2.9), we see that (4.16) holds, successively, for $X^{d+1-m}Y^m, \dots, Y^{d+1}$, which completes the proof. \square

The following result shows that the rows of $(M_d \ B(d+1))$ of the form $X^f Y^{m+g}$ ($f, g \geq 0, m+f+g \leq d$) are recursively determined from row Y^m .

Lemma 4.7. For $i, j \geq 0, i+j \leq d+1$ and for $f, g \geq 0, m+f+g \leq d$,

$$\langle X^i Y^j, X^f Y^{m+g} \rangle = \sum_{u,v \geq 0, u+v \leq m, v < m} b_{uv} \langle X^i Y^j, X^{u+f} Y^{v+g} \rangle. \quad (4.19)$$

Proof. Since M_d is real symmetric and recursively generated, its rows are also recursively generated from (2.1) and (2.2), so (4.19) holds if $i+j \leq d$. We may now assume $j = d+1-i$, and we first consider the case $m+f+g < d$ and the subcase $i \leq d$. Since $f+g+m < d$, using old moments we see that

$$\begin{aligned} \langle X^i Y^{d+1-i}, X^f Y^{m+g} \rangle &= \beta_{i+f, d+1-i+g+m} \\ &= \langle X^f Y^{m+g+1}, X^i Y^{d-i} \rangle \quad (\text{in } M_d) \\ &= \sum b_{uv} \langle X^{u+f} Y^{v+g+1}, X^i Y^{d-i} \rangle \\ &= \beta_{i+u+f, d-i+v+g+1} \quad (\text{in } M_d) \\ &= \sum b_{uv} \langle X^i Y^{d+1-i}, X^{u+f} Y^{v+g} \rangle \\ &\quad (\text{since } u+v+f+g < d). \end{aligned}$$

The subcase when $i = d+1$ proceeds as above, but starting with $\langle X^{d+1}, X^f Y^{m+g} \rangle = \beta_{d+1+f, m+g} = \langle X^d, X^{f+1} Y^{m+g} \rangle = \langle X^{f+1} Y^{m+g}, X^d \rangle$. For the case $m+f+g = d$, we first consider the subcase when $i \geq n$, so $X^i Y^{d+1-i}$ is in the left recursive band. We have

$$\begin{aligned} \langle X^i Y^{d+1-i}, X^f Y^{m+g} \rangle &= \langle X^n X^{i-n} Y^{d+1-i}, X^f Y^{m+g} \rangle \\ &= \sum a_{rs} \langle X^{r+i-n} Y^{s+d+1-i}, X^f Y^{m+g} \rangle \\ &= \sum a_{rs} \sum b_{uv} \langle X^{r+i-n} Y^{s+d+1-i}, X^{u+f} Y^{v+g} \rangle \\ &\quad (\text{by row recursiveness in } M_d) \\ &= \sum b_{uv} \langle X^i Y^{d+1-i}, X^{n+f} Y^{v+g} \rangle. \end{aligned}$$

In the next subcase, we consider a column in the center band, of the form $X^{d+1-i} Y^i$ with $d+1-n < i < m$. In this case, (4.19) is equivalent to

$$\langle X^{d+1-i} Y^i, X^f Y^{m+g} \rangle = \sum_{u,v \geq 0, u+v \leq m, v < m} b_{uv} \langle X^{d+1-i} Y^i, X^{u+f} Y^{v+g} \rangle. \quad (4.20)$$

Note that the component $\langle X^{d+1-i} Y^i, X^f Y^{m+g} \rangle$ lies on a cross-diagonal that reaches column $X^{d+1-m} Y^m$, so since $B(d+1)$ is well-defined, we have

$$\begin{aligned} \langle X^{d+1-i} Y^i, X^f Y^{m+g} \rangle &= \langle X^{d+1-m} Y^m, X^{f+m-i} Y^{g+i} \rangle \\ &= \sum_{u,v \geq 0, u+v \leq m, v < m} b_{uv} \langle X^{u+d+1-m} Y^v, X^{f+m-i} Y^{g+i} \rangle. \quad (4.21) \end{aligned}$$

For the subcase when $u+v < m$, in M_d we have

$$\begin{aligned} \langle X^{u+d+1-m} Y^v, X^{f+m-i} Y^{g+i} \rangle &= \beta_{u+d+1+f-i, v+g+i} \\ &= \langle X^{d+1-i} Y^i, X^{u+f} Y^{v+g} \rangle \quad (\text{since } u+f+v+g \leq d-1). \end{aligned}$$

For the subcase when $u + v = m$, there are three further subcases in showing that

$$\langle X^{d+1-i}Y^i, X^{u+f}Y^{v+g} \rangle = \langle X^{u+d+1-m}Y^v, X^{f+m-i}Y^{g+i} \rangle. \quad (4.22)$$

For $v = i$, (4.22) is clear. For $v < i$, the Hankel property in $B[d, d + 1]$ implies

$$\begin{aligned} \langle X^{d+1+u-m}Y^v, X^{m+f-i}Y^{g+1} \rangle &= \langle X^{d+1+u-m-(i-v)}Y^{v+(i-v)}, X^{m+f-i+(i-v)}Y^{g+i-(i-v)} \rangle \\ &= \langle X^{d+1-i}Y^i, X^{u+f}Y^{g+v} \rangle. \end{aligned}$$

For $v > i$ we have, similarly,

$$\begin{aligned} \langle X^{d+1-i}Y^i, X^{u+f}Y^{v+g} \rangle &= \langle X^{d+1-i-(v-i)}Y^{i+v-i}, X^{u+f+v-i}Y^{v+g-(v-i)} \rangle \\ &= \langle X^{d+1+u-m}Y^v, X^{m+f-i}Y^{g+i} \rangle. \end{aligned}$$

Since (4.19) holds in M_d and in all columns of the left and center bands, it now follows, using (2.9) successively, that it holds for columns in the right recursive band, which completes the proof. \square

We are now prepared to prove that $\text{Ran } B(d + 1) \subseteq \text{Ran } M_d$. It follows immediately from (2.8) that each column in the left recursive band of $B(d + 1)$ belongs to $\text{Ran } M_d$. In view of (2.9), to establish range inclusion, it suffices to show that each central-band column of $B(d + 1)$ belongs to $\text{Ran } M_d$. Let \mathcal{S} denote the set of recursively determined columns of M_d , i.e.,

$$\mathcal{S} = \{X^n, X^{n+1}, X^n Y, \dots, X^d, \dots, X^n Y^{d-n}, \dots, Y^m, X Y^m, Y^{m+1}, \dots, X^{d-m} Y^m, \dots, Y^d\}.$$

Let \mathcal{B} denote the basis for $\text{Col } M_d$ (the column space of M_d) consisting of those columns of M_d which do not belong to \mathcal{S} . Let $M_{\mathcal{B}}$ denote the compression of M_d to the rows and columns indexed by \mathcal{B} . Since $M_d \succeq 0$, we also have $M_{\mathcal{B}} \succeq 0$. Let $X^i Y^{d+1-i}$ ($d + 1 - m < i < n$) denote a central-band column of $B(d + 1)$, and let $v_i \equiv [X^i Y^{d+1-i}]_{\mathcal{B}}$ denote the compression of $X^i Y^{d+1-i}$ to the rows of \mathcal{B} . There exists a unique vector of coefficients $(c_{ab}^{(i)})_{X^a Y^b \in \mathcal{B}}$ such that

$$v_i = \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} [X^a Y^b]_{\mathcal{B}},$$

i.e., for each $X^u Y^v \in \mathcal{B}$,

$$\langle X^i Y^{d+1-i}, X^u Y^v \rangle = \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} \langle X^a Y^b, X^u Y^v \rangle. \quad (4.23)$$

To complete the proof that $\text{Ran } B(d + 1) \subseteq \text{Ran } M_d$, it suffices to prove that $X^i Y^{d+1-i} = \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} X^a Y^b$, which, in view of (4.23), follows from the next result.

Lemma 4.8. *For each $X^c Y^e \in \mathcal{S}$,*

$$\langle X^i Y^{d+1-i}, X^c Y^e \rangle = \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} \langle X^a Y^b, X^c Y^e \rangle. \quad (4.24)$$

Proof. We may assume without loss of generality that $n \leq m$, so the elements of \mathcal{S} may be arranged in degree-lexicographic order as $X^n, \dots, Y^m, \dots, X^d, \dots, Y^d$. We will prove (4.24) by induction on the position number of row $X^c Y^e \in \mathcal{S}$ within the degree-lexicographic ordering. For row X^n ($c = n, e = 0$), Lemma 4.6 implies that

$$\langle X^i Y^{d+1-i}, X^n \rangle = \sum_{r,s \geq 0, r+s \leq n-1} a_{rs} \langle X^i Y^{d+1-i}, X^r Y^s \rangle. \quad (4.25)$$

Since $r + s < n$, $X^r Y^s \in \mathcal{B}$, so the sum in (4.25) may be expressed as

$$\sum a_{rs} \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} \langle X^a Y^b, X^r Y^s \rangle = \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} \langle X^a Y^b, X^n \rangle$$

(using Lemma 4.6 again). Assume now that (4.24) holds for all rows $X^C Y^E \in \mathcal{S}$ with order position up to $k - 1$, and consider $X^c Y^e \in \mathcal{S}$ with position k . Either $c \geq n$ or $e \geq m$; we present the argument for the case $e \geq m$ (the other case is simpler). We have $e = m + g$ for some $g \geq 0$. From Lemma 4.7, we have

$$\langle X^i Y^{d+1-i}, X^c Y^{m+g} \rangle = \sum_{u,v \geq 0, u+v \leq m, v < m} b_{uv} \langle X^i Y^{d+1-i}, X^{c+u} Y^{g+v} \rangle.$$

Now $X^{c+u} Y^{g+v}$ is either a basis vector, or, since $v < m$, it precedes $X^c Y^{m+g}$ in the ordering of \mathcal{S} . Thus, by definition (for the basis rows) and by induction (for the non-basis rows), the preceding sum is equal to

$$\begin{aligned} &= \sum b_{uv} \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} \langle X^a Y^b, X^{c+u} Y^{g+v} \rangle = \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} \langle X^a Y^b, \sum b_{uv} X^{c+u} Y^{g+v} \rangle \\ &= \sum_{X^a Y^b \in \mathcal{B}} c_{ab}^{(i)} \langle X^a Y^b, X^c Y^e \rangle \end{aligned}$$

(by another application of Lemma 4.7). □

The proof of Theorem 2.1 is now complete.

5 Proof of Theorem 2.13

For the proof of Theorem 2.13, we require a preliminary result concerning a general moment matrix.

Lemma 5.1. *Suppose M_{d+1} satisfies $\text{Ran } B(d+1) \subseteq \text{Ran } M_d$. If $p \in \mathcal{P}_d$ and $p(X, Y) = 0$ in $\text{Col } M_d$, then $p(X, Y) = 0$ in $\text{Col } M_{d+1}$.*

Proof. Since M_d is real symmetric, we have $p(X, Y) = 0$ in the row space of M_d , and we first show that $p(X, Y) = 0$ holds in the row space of $(M_d \ B(d+1))$. Let $\rho := \deg p$ and suppose $p(x, y) \equiv \sum_{r,s \geq 0, r+s \leq \rho} a_{rs} x^r y^s$. Then for $i, j \geq 0$ with $i + j \leq d$, we have

$$\sum_{r,s} \alpha_{rs} \langle X^i Y^j, X^r Y^s \rangle = 0. \quad (5.26)$$

Consider a column of degree $d + 1$, $X^u Y^{d+1-u}$ ($0 \leq u \leq d + 1$). We seek to show that

$$\sum_{r,s} \alpha_{rs} \langle X^u Y^{d+1-u}, X^r Y^s \rangle = 0. \quad (5.27)$$

By the range inclusion, we have a dependence relation in $Col (M_d \ B(d + 1))$ of the form

$$X^u Y^{d+1-u} = \sum_{a,b \geq 0, a+b \leq d} c_{ab}^{(u)} X^a Y^b. \quad (5.28)$$

Thus,

$$\begin{aligned} \sum_{r,s} \alpha_{rs} \langle X^u Y^{d+1-u}, X^r Y^s \rangle &= \sum_{r,s} \alpha_{rs} \sum_{a,b \geq 0, a+b \leq d} c_{ab}^{(u)} \langle X^a Y^b, X^r Y^s \rangle \\ &= \sum_{a,b} c_{ab}^{(u)} \sum_{r,s} \alpha_{rs} \langle X^a Y^b, X^r Y^s \rangle = 0 \quad (\text{by (5.26)}). \end{aligned}$$

Now, $p(X, Y) = 0$ in the row space of $(M_d \ B(d + 1))$, so $p(X, Y) = 0$ in $Col \begin{pmatrix} M_d \\ B(d + 1)^T \end{pmatrix}$. \square

Proof of Theorem 2.13. It follows from the proof of Theorem 2.5 that M_d admits a unique extension M_{d+1} which satisfies $Ran \ B(d + 1) \subseteq Ran \ M_d$ and such that (2.8)-(2.9) hold in $Col \ M_{d+1}$. It remains only to prove that M_{d+1} is recursively generated. Since M_d is recursively generated, it suffices to consider a dependence relation in $Col \ M_{d+1}$ of degree d , say

$$X^i Y^{d-i} = \sum_{g,h \geq 0, g+h \leq d-1} c_{gh} X^g Y^h \quad (5.29)$$

(where $0 \leq i \leq d$), and to show that

$$X^{i+1} Y^{d-i} = \sum_{g,h \geq 0, g+h \leq d-1} c_{gh} X^{g+1} Y^h \quad (5.30)$$

and

$$X^i Y^{d-i+1} = \sum_{g,h \geq 0, g+h \leq d-1} c_{gh} X^g Y^{h+1}. \quad (5.31)$$

Suppose first that $i \geq n$, so that $X^i Y^{d-i}$ lies in the left band. Then from (2.3) we also have

$$X^i Y^{d-i} = \sum_{r+s \leq n-1} a_{rs} X^{i-n+r} Y^{s+d-i}. \quad (5.32)$$

Thus, in M_d we have the column relation of degree at most $d - 1$,

$$\sum_{g+h \leq d-1} c_{gh} X^g Y^h = \sum_{r+s \leq n-1} a_{rs} X^{i-n+r} Y^{s+d-i}.$$

Since M_d is recursively generated, it follows that in $Col \ M_d$ we also have

$$\sum_{g+h \leq d-1} c_{gh} X^{g+1} Y^h = \sum_{r+s \leq n-1} a_{rs} X^{i-n+r+1} Y^{s+d-i}.$$

Lemma 5.1 implies that the last equation also holds in $Col M_{d+1}$, where, from (2.9), the right-hand sum represents $X^{i+1}Y^{d-i}$; this establishes (5.30). We omit the proof of (5.31), which is similar. The case when $d-i \geq m$, so that X^iY^{d-i+1} is in the right band, is handled in an entirely analogous fashion, so we also omit the proof of this case.

We next consider the case when $d-m < i < n$, so that column X^iY^{d-i} in (5.29) is in the central band. To establish (5.30), it suffices to verify that

$$\langle X^{i+1}Y^{d-i}, X^kY^j \rangle = \sum_{g,h \geq 0, g+h \leq d-1} c_{gh} \langle X^{g+1}Y^h, X^kY^j \rangle \quad (k, j \geq 0, k+j \leq d+1). \quad (5.33)$$

The case when $k+j < d$ is easy, using (5.29) and the old moments in block $B[k+j, d+1]$. We consider next the case $k+j = d$ and the subcase when $k \geq n$. In this subcase, $\langle X^{i+1}Y^{d-i}, X^kY^{d-k} \rangle$ belongs to a cross-diagonal of $B[d, d+1]$ that intersects column X^nY^{d+1-n} , so from the definition of $B[d, d+1]$ in the proof of Theorem 2.1, we have

$$\begin{aligned} \langle X^{i+1}Y^{d-i}, X^kY^{d-k} \rangle &:= \langle X^nY^{d+1-n}, X^{k-(n-i-1)}Y^{d-k+n-i-1} \rangle \\ &= \sum a_{rs} \langle X^rY^{s+d+1-n}, X^{k-(n-i-1)}Y^{d-k+n-i-1} \rangle. \end{aligned} \quad (5.34)$$

Now, we have

$$\begin{aligned} \sum_{g,h \geq 0, g+h \leq d-1} c_{gh} \langle X^{g+1}Y^h, X^kY^{d-k} \rangle &= \sum c_{gh} \langle X^kY^{d-k}, X^{g+1}Y^h \rangle \\ &= \sum c_{gh} \sum_{rs} a_{rs} \langle X^{r+k-n}Y^{s+d-k}, X^{g+1}Y^h \rangle \\ &= \sum a_{rs} \sum c_{gh} \langle X^{g+1}Y^h, X^{r+k-n}Y^{s+d-k} \rangle \\ &= \sum a_{rs} \sum c_{gh} \langle X^gY^h, X^{r+k-n+1}Y^{s+d-k} \rangle \quad (\text{in } M_d) \\ &= \sum a_{rs} \langle X^iY^{d-i}, X^{r+k-n}Y^{s+d-k} \rangle \\ &= \sum a_{rs} \langle X^{r+k-n}Y^{s+d-k}, X^iY^{d-i} \rangle \\ &= \sum a_{rs} \langle X^rY^{s+d-k+(k-n+1)}, X^{i+(k-n+1)}Y^{d-i-(k-n+1)} \rangle. \end{aligned}$$

This last expression agrees with (5.34), so (5.30) is established for this subcase. The proof of this subcase for (5.31) is very similar, so we omit the details. In the subcase when $k < n$, then $d-k \geq m$, and we see that $\langle X^{i+1}Y^{d-i}, X^kY^{d-k} \rangle$ belongs to a cross-diagonal of $B[d, d+1]$ that intersects column $X^{d+1-m}Y^m$. Since $\deg q < m$, the proof of this subcase is entirely analogous to that above, but using (2.9) for the definition of $X^{d+1-m}Y^m$.

Finally, we consider the case $k+j = d+1$. As above, we will treat the subcase of (5.30) when $k \geq n$ in detail and omit the proofs of the other subcases of (5.30) and (5.31), which are similar. Since $k \geq n$, then, as above, we have

$$\begin{aligned} \langle X^{i+1}Y^{d-i}, X^kY^{d+1-k} \rangle &:= \langle X^nY^{d+1-n}, X^{k-(n-i-1)}Y^{d+1-k+n-i-1} \rangle \\ &= \sum a_{rs} \langle X^rY^{s+d+1-n}, X^{k-(n-i-1)}Y^{d-k+n-i} \rangle. \end{aligned} \quad (5.35)$$

Now,

$$\begin{aligned}
& \sum_{g,h \geq 0, g+h \leq d-1} c_{gh} \langle X^{g+1} Y^h, X^k Y^{d+1-k} \rangle \\
&= \sum c_{gh} \langle X^k Y^{d+1-k}, X^{g+1} Y^h \rangle \\
&\quad (\text{since } \begin{pmatrix} M_d & B(d+1) \end{pmatrix} \text{ is the transpose of } \\
&\quad \begin{pmatrix} M_d \\ B(d+1)^T \end{pmatrix}) \\
&= \sum c_{gh} \sum_{rs} a_{rs} \langle X^{r+k-n} Y^{s+d+1-k}, X^{g+1} Y^h \rangle \\
&= \sum a_{rs} \sum c_{gh} \langle X^{g+1} Y^h, X^{r+k-n} Y^{s+d+1-k} \rangle.
\end{aligned}$$

Since the row degrees of the terms in the last sum are at most d , by the previous cases (for $j+k < d$ and $j+k = d$), the last double sum may be expressed as

$$\sum a_{rs} \langle X^{i+1} Y^{d-i}, X^{r+k-n} Y^{s+d+1-k} \rangle$$

relative to $\begin{pmatrix} M_d & B(d+1) \end{pmatrix}$. Since M_{d+1} is real symmetric, the latter sum may be expressed as

$$\begin{aligned}
& \sum a_{rs} \langle X^{r+k-n} Y^{s+d+1-k}, X^{i+1} Y^{d-i} \rangle \\
&= \sum a_{rs} \langle X^r Y^{s+d+1-k+(k-n)}, X^{i+1+(k-n)} Y^{d-i-(k-n)} \rangle \\
&= \sum a_{rs} \langle X^r Y^{s+d+1-n}, X^{i+1+k-n} Y^{d-i-k+n} \rangle,
\end{aligned}$$

and this agrees with (5.35). The proof is now complete. \square

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