## From Coloring Maps to Avoiding Conflicts

# Nathaniel Dean, Robert M. Nehs, and Tong Wu 

Department of Mathematical Sciences
Texas Southern University
3100 Cleburne Avenue
Houston, TX 77004

## Map Coloring

Countries with a common boundary must have different colors.


## Four Color Problem

1852 letter by Augustus de Morgan to Sir William Hamilton:
Four colors are required. Do 4 colors suffice?


1976: Appel and Haken proved it using an intricate case analysis on a computer.

## Exercise: <br> Draw a map that requires four colors.



## 3-Coloring Maps

## Computer Science project by Malvika Rao (student), McGill U. http://www.cs.mcgill.ca/~rao/cs507/MapColoring.html



Welcome! Select a map or draw one.

## The Dual is a Planar Graph.



## Graph G=(V,E) consists of a set V of objects called vertices and a set E of unordered pairs of vertices.



## Vertex Coloring

- A $k$-coloring is a labeling $\mathrm{f}: \mathrm{V}(\mathrm{G}) \rightarrow\{1,2, \ldots, \mathrm{k}\}$.
- A $k$-coloring is proper if $x y \in E(G)$ implies $f(x) \neq f(y)$.
- $G$ is $k$-colorable if it has a proper k-coloring.
- The chromatic number $\chi(\mathrm{G})$ is the smallest $k$ such that G is k -colorable.



## Exercise: <br> Prove $\chi($ Moser Graph $)=4$.



## Party Problem

- People $P_{1}, P_{2}, \ldots, P_{n}$ meet for a party, but certain pairs are incompatible.
- Goal: Assign people to rooms so that no two people in the same room are incompatible.
- How many rooms are needed?


## Solution to the Party Problem

## Construct a conflict graph G.

- $V(G)=\left\{P_{1}, P_{2}, \ldots, P_{n}\right\}$.
- $P_{i}, P_{j} \in E(G)$ iff $P_{i}$ and $P_{j}$ are incompatible.
- The chromatic number $\chi(\mathrm{G})$ is the least number of rooms.



## Scheduling Problem

- Five different groups of students $\{1,2,3\}$, $\{6,7\},\{1,7,9\},\{4,6,8\},\{2,3,4\}$ must take exams in the following engineering courses $\mathrm{S}_{1}, \mathrm{~S}_{2}, \mathrm{~S}_{3}, \mathrm{~S}_{4}, \mathrm{~S}_{5}$, respectively.
- Goal: Schedule the exams using a minimum number of time periods.


## Solution to the Scheduling Problem

## Construct a conflict graph $G$.

- $\mathrm{V}(\mathrm{G})=\left\{\mathrm{S}_{1}, \mathrm{~S}_{2}, \mathrm{~S}_{3}, \mathrm{~S}_{4}, \mathrm{~S}_{5}\right\}$.
- $S_{i,}, S_{j} \in E(G)$ iff $S_{i} \cap S_{j} \neq \varnothing$.
- The chromatic number $\chi(\mathrm{G})$ is the minimum number of time periods.



## Future Work

- Allow student to enter any arbitrary map or (nonplanar) graph.
- Add a map library, and allow transferring between map and corresponding graph.
- Develop algorithms with student.
- Explore applications.


## Thanks!



Frank



David

- Directors:

Lang Moore
David Smith
Frank Wattenberg

- Funding Agencies


MAA


