

(\Rightarrow) Suppose s is graphic.

Claim s^* is graphic.

s graphic implies \exists simple graph $G = (V, E)$ with degree sequence s , where $V = \{v_1, \dots, v_n\}$ where $\delta(v_i) = d_i$.

Let $G^* = G - v_1 = (V^*, E^*)$ where $V^* = \{v_2, \dots, v_n\}$ and $E^* = \{e \in E \mid e \text{ is not adjacent to } v_1\}$

Note G^* is a simple graph.

Case 1: Suppose v_1 is adjacent to v_2, \dots, v_{d_1+1} and all edges adjacent to v_1 . Removing v_1 reduces the degree of v_2, \dots, v_{d_1+1} by one. Thus G^* has degree sequence s^* . Thus s^* is graphic if case 1 holds.

Case 2: Suppose $\exists j \in \{2, \dots, v_{d_1+1}\}$ such that v_1 is NOT adjacent to v_j .

Then since $\delta(v_1) = d_1$, $\exists \ell > d_1 + 1$ such that v_1 is adjacent to v_ℓ .

See chalkboard for rest of proof.

More elegant proof of s is graphic implies s^* is graphic.

(\Rightarrow) Suppose s is graphic.

Claim s^* is graphic.

s graphic implies \exists simple graph $G = (V, E)$ with degree sequence s , where $V = \{v_1, \dots, v_n\}$ where $\delta(v_i) = d_i$.

Let $N(v_1) = \{w \in V \mid w > v_1, w \in E\}$

Let $A_G = N(v_1) \cap \{v_2, \dots, v_{d_1+1}\}$

Note $0 \leq |A_G| \leq d_1$.

Among all simple graphs G with degree sequence s , choose one such that $|A_G|$ is as large as possible.

Let $G^* = G - v_1 = (V^*, E^*)$ where $V^* = \{v_2, \dots, v_n\}$ and $E^* = \{e \in E \mid e \text{ is not adjacent to } v_i\}$

Note G^* is a simple graph.

Case 1: Suppose $|A_G| = d_1$. Then v_1 is adjacent to v_2, \dots, v_{d_1+1}

Then G^* has degree sequence s^* . Thus s^* is graphic. ■