

An “Infamous” Example

Below we have an initial algebra specification by Goguen et al (chap. 5 of Yeh book) to describe the (signed) integer ADT (with sum and product operations) given Boolean and Nat ADTs. The intuitive idea is to pair a natural number with a Boolean value that represents its sign, and then express signed arithmetic (assuming we already know unsigned arithmetic). However, it will be seen that the initial algebra of this ADT is *not* the usual algebra of the integers.

The specification

Signat signature

Pre-defined types: Boolean (with '='), and Nat ($\{0, 1, 2, \dots\}$) with operations as usual, including $+$, $\dot{-}$, $*$, \leq (note that $\dot{-}$ is “proper subtraction”, $n \dot{-} m$ yields 0 when $n \leq m$).

Operation signatures:

PAIR: Nat, Bool \rightarrow Signat

ABS: Signat \rightarrow Nat

SGN: Signat \rightarrow Bool

SUM: Signat, Signat \rightarrow Signat

PROD: Signat, Signat \rightarrow Signat

Semantic equations (for all $s \in \text{Signat}$, $n \in \text{Nat}$, $b \in \text{Bool}$)

1. PAIR (ABS (s) , SGN (s)) = s
2. ABS (PAIR (n,b)) = n
3. SGN (PAIR (n,b)) = b
4. SUM (s₁ , s₂) = **if** SGN (s₁)=SGN (s₂)
 then PAIR (ABS (s₁) +ABS (s₂) , SGN (s₁))
 else if ABS (s₁) \leq ABS (s₂)
 then PAIR (ABS (s₂) $\dot{-}$ ABS (s₁) , SGN (s₂))
 else PAIR (ABS (s₁) $\dot{-}$ ABS (s₂) , SGN (s₁))
5. PROD (s₁ ,s₂) = PAIR (ABS (s₁) *ABS (s₂) , SGN (s₁)=SGN (s₂))

The flaw

The fault to be found with this specification of the signed integers is that there are *two* “zeros” — PAIR(0,True) or +0, and PAIR(0,False) or -0. There are no equations that enable us to deduce two different pairs are equivalent, and so in the initial algebra view they are different. This might appear to be a minor oversight, but in fact having two distinct representations of zero causes numerous familiar identities to be invalidated.

For instance, for all x , $x+0 = x$ in the integers. But neither of the corresponding values in this specification has this property — note that $SUM(+0,-0) = -0$, and $SUM(-0,+0) = +0$ (use the equations in the specification on the term forms of these values to confirm this). Also $x*0 = 0$ in the integers, but in the specification $PROD(-5,+0) = -0$, and $PROD(-0,-0) = +0$. If it is really the system of signed natural numbers we seek to specify, the specification given does not qualify.

A defective repair

In a widely circulated IBM technical report that preceded publication, the authors “corrected” the problem in the SigNat ADT described above by proposing the single additional axiom

$$6. \text{PAIR}(0,\text{True}) = \text{PAIR}(0,\text{False})$$

to unify these two different equivalence classes (i.e., $+0 = -0$). While at first glance this seems like a simple and obvious solution, it is a huge blunder. If we add this axiom to those we already have for SigNat, then

$$\text{True} \equiv_3 \text{SGN}(\text{PAIR}(0,\text{True})) \equiv_6 \text{SGN}(\text{PAIR}(0,\text{False})) \equiv_3 \text{False!}$$

Hence an inconsistency in the pre-defined type Boolean has been introduced into the specification, and the original flawed “approximate specification” has been destroyed rather than repaired.

An actual repair

A suitable correction to the original flaw is to replace equation 3 by

$$3'. \text{SGN}(n,b) = \text{if } n=0 \text{ then True else } b$$

Then the two representations of zero are equivalent:

$$\text{PAIR}(0,\text{False}) \equiv_1 \text{PAIR}(\text{ABS}(\text{PAIR}(0,\text{False})), \text{SGN}(\text{PAIR}(0,\text{False}))) \equiv_2$$

$$\text{PAIR}(0, \text{SGN}(\text{PAIR}(0,\text{False}))) \equiv_3, \text{PAIR}(0, \text{True}).$$

So now the two zeros fall into the same equivalence class, but no inconsistency is introduced.