

A L^AT_EX and ZSL Input Notations

A.1 Paragraphs

A.1.1 Axiom Box

$$\left| \begin{array}{l} D_1; \dots; D_m \\ \hline P_1; \dots; P_n \end{array} \right.$$

L^AT_EX input:

```
\begin{axdef}
  D_1; ... ; D_m
\where
  P_1; ... ; P_n
\end{axdef}
```

ZSL input – text style:

```
global
  D1; ... ; Dm
axiom
  P1; ... ; Pn
end axiom
```

ZSL input – box style:

```
| D1; ... ; Dm
|-----
| P1; ... ; Pn
```

$$\left| D_1; \dots; D_m \right.$$

L^AT_EX input:

```
\begin{axdef}
  D_1; ... ; D_m
\end{axdef}
```

ZSL input – text style:

```
global
  D1; ... ; Dm
end global
```

ZSL input – box style:

```
| D1; ... ; Dm
```

$P_1; \dots; P_n$
L^AT_EX input:

```
\begin{zed}
  P_1; ... ; P_n
\end{zed}
```

ZSL input

```
axiom
  P1; ... ; Pn
end axiom
```

A.1.2 Schema Box

S
$D_1; \dots; D_m$
$P_1; \dots; P_n$

L^AT_EX input:

```
\begin{schema}{S}
  D_1; ... ; D_m
\where
  P_1; ... ; P_n
\end{schema}
```

ZSL input – text style:

```
schema S
  D1; ... ; Dm
where
  P1; ... ; Pn
end schema
```

ZSL input – box style:

```
----- S -----
|  D1; ... ; Dm
|-----
|  P1; ... ; Pn
|-----
```

A.1.3 Generic Schema Box

$S[X_1, \dots, X_k]$
$D_1; \dots; D_m$
$P_1; \dots; P_n$

L^AT_EX input:

```
\begin{schema}{S[X_1, \dots X_k]}
  D_1; ... ; D_m
\where
  P_1; ... ; P_n
\end{schema}
```

ZSL input – text style:

```
schema S [X1, ... , Xk]
  D1; ... ; Dm
where
  P1; ... ; Pn
end schema
```

ZSL input – box style:

```
--- S [X1, ... , Xk] -----
| D1; ... ; Dm
|-----
| P1; ... ; Pn
|-----
```

A.1.4 Generic Box

$[X_1, \dots, X_k]$
$D_1; \dots; D_m$
$P_1; \dots; P_n$

L^AT_EX input zed:

```
\begin{gendef}[X_1, \dots, X_k]
  D_1; ... ; D_m
\where
  P_1; ... ; P_n
\end{gendef}
```

L^AT_EX input oz:

```
\begin{gendef}{X_1, \dots, X_k}
  D_1; ... ; D_m
\where
  P_1; ... ; P_n
\end{gendef}
```

ZSL input – text style:

```
generic [X1, ... , Xk]
  D1; ... ; Dm
where
  P1; ... ; Pn
end generic
```

ZSL input – box style:

```

=== [X1, ... , Xk] =====
| D1; ... ; Dm
|-----
| P1; ... ; Pn
-----
    
```

A.1.5 Schema Definition

$S \hat{=} [D P]$	L ^A T _E X zed S \defs [D P]	ZSL S ^= [D P] S is [D P]
	L ^A T _E X oz S \sdef [D P]	

A.1.6 Given Set

$[T_1, \dots, T_n]$	L ^A T _E X [T_1, ..., T_n]	ZSL [T1, ..., Tn]
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A.1.7 Equivalence Definition

$id == Exp$	L ^A T _E X zed id == Exp	ZSL id == Exp
	L ^A T _E X oz id \defs Exp	

A.1.8 Free Type Definition

$ \begin{array}{l} T ::= c_1 \dots c_m \\ d_1 \langle\langle E_1[T] \rangle\rangle \\ \dots \\ d_n \langle\langle E_n[T] \rangle\rangle \end{array} $	L ^A T _E X input zed: <pre> \begin{syntax} T & ::= & c_1 ... c_m \\ & & d_1 \lldata E_1[T] \rdata \\ & & ... \\ & & d_n \lldata E_n[T] \rdata \end{syntax} </pre>
--	---

\LaTeX input oz:

```
\begin{syntax}
T & \ddef & c_1 | ... | c_m \\
  & | & d_1 \lang E_1[T] \rang \\
  & | & ... \\
  & | & d_n \lang E_n[T] \rang
\end{syntax}
```

ZSL input:

```
T ::= c1 | ... | cm
      | d1 << E1[T] >>
      | ...
      | dn << En[T] >>
```

A.1.9 Schema Expressions

	\LaTeX	ZSL
$\forall D P \bullet S$	<code>\forall D P @ S</code>	<code>forall D P @ S</code>
$\exists D P \bullet S$	oz only ▶ <code>\all D P \dot S</code>	
	<code>\exists D P @ S</code>	<code>exists D P @ S</code>
	oz only ▶ <code>\exi D P \dot S</code>	
$\exists_1 D P \bullet S$	<code>\existss_1 D P @ S</code>	<code>exists1 D P @ S</code>
	oz only ▶ <code>\exione D P \dot S</code>	
$[D P]$	<code>[D P]</code>	<code>[D P]</code>
ΔS	<code>\Delta S</code>	<code>Delta S</code>
ΞS	<code>\Xi S</code>	<code>Xi S</code>
$S[T_1, \dots, T_n]$	<code>S[T_1, \dots, T_n]</code>	<code>S[T1, \dots, Tn]</code>
$S[x_1/y_1, \dots, x_n/y_n]$	<code>S[x_1/y_1, \dots, x_n/y_n]</code>	<code>S[x1/y1, \dots, xn/yn]</code>
$\text{pre } S$	<code>\pre S</code>	<code>pre S</code>
$\neg S$	<code>\lnot S</code>	<code>not S</code>
$S_1 \wedge S_2$	<code>S_1 \land S_2</code>	<code>S1 and S2</code>
		<code>S1 /\ S2</code>
$S_1 \vee S_2$	<code>S_1 \lor S_2</code>	<code>S1 or S2</code>
		<code>S1 \/ S2</code>
$S_1 \Rightarrow S_2$	<code>S_1 \implies S_2</code>	<code>S1 implies S2</code>
	oz only ▶ <code>S_1 \imp S_2</code>	<code>S1 => S2</code>
$S_1 \Leftrightarrow S_2$	<code>S_1 \iff S_2</code>	<code>S1 iff S2</code>
		<code>S1 <=> S2</code>
$S_1 \upharpoonright S_2$	<code>S_1 \project S_2</code>	<code>S1 project S2</code>
		<code>S1 \ S2</code>
$S \setminus (v_1, \dots, v_n)$	<code>S \hide (v_1, \dots, v_n)</code>	<code>S hide (v1, \dots, vn)</code>
	oz only ▶ <code>S \zhide (v_1, \dots, v_n)</code>	<code>S \\ (v1, \dots, vn)</code>
$S_1 \circledast S_2$	<code>S_1 \semi S_2</code>	<code>S1 semi S2</code>
	oz only ▶ <code>S_1 \zcmp S_2</code>	<code>S1 // S2</code>
$S_1 \gg S_2$	<code>S_1 \pipe S_2</code>	<code>S1 pipe S2</code>
	oz only ▶ <code>S_1 \zpipe S_2</code>	

A.1.10 Predicates

	L ^A T _E X	ZSL
$\forall D P \bullet Q$	<code>\forall D P @ Q</code>	<code>forall D P @ Q</code>
	oz only ▶ <code>\all D P \dot S</code>	
$\exists D P \bullet Q$	<code>\exists D P @ Q</code>	<code>exists D P @ Q</code>
	oz only ▶ <code>\exi D P \dot S</code>	
$\exists_1 D P \bullet Q$	<code>\exists_1 D P @ Q</code>	<code>exists1 D P @ Q</code>
	oz only ▶ <code>\exione D P \dot S</code>	
let $v == e \bullet P$	<code>\zlet v==e @ P</code>	<code>let v==e @ P</code>
	oz only ▶ <code>\zlet v==e \dot P</code>	
$p \wedge q$	<code>p \land q</code>	<code>p and q</code>
$p \vee q$	<code>p \lor q</code>	<code>p /\ q</code>
$p \Rightarrow q$	<code>p \implies q</code>	<code>p or q</code>
	oz only ▶ <code>p \imp q</code>	<code>p \ / q</code>
$p \Leftrightarrow q$	<code>p \iff q</code>	<code>p implies q</code>
$\neg p$	<code>\lnot p</code>	<code>p => q</code>
<i>true</i>	<code>true</code>	<code>p iff q</code>
		<code>p <=> q</code>
<i>false</i>	<code>false</code>	<code>not p</code>
		<code>true</code>
		<code>TRUE</code>
		<code>false</code>
		<code>FALSE</code>

A.2 Expressions

A.2.1 Lambda Expression

	L ^A T _E X	ZSL
$\lambda D P \bullet E$	<code>\lambda D P @ E</code>	<code>lambda D P @ E</code>

A.2.2 Definite Description

	L ^A T _E X	ZSL
$\mu D P \bullet E$	<code>\mu D P @ E</code>	<code>mu D P @ E</code>
	oz only ▶ <code>\mu D P \dot E</code>	<code>unique D P @ E</code>

A.2.3 Conditional expression

	\LaTeX	ZSL
if P then E_1 else E_2	<code>\zif P \zthen E_1 \zelse E_2</code>	<code>if P then E1 else E2</code>

A.2.4 Local definition

	\LaTeX	ZSL
let $v == e \bullet E$	<code>\zlet v==e @ E</code>	<code>let v==e @ E</code>
oz only ▶	<code>\zlet v==e \dot E</code>	

A.2.5 Sets

	\LaTeX	ZSL
$\{x_1, \dots, x_n\}$	<code>\{ x_1, \dots, x_n \}</code>	<code>{ x1, ..., xn }</code>
$\{D P \bullet E\}$	<code>\{ D P @ E \}</code>	<code>{ D P @ E }</code>
oz only ▶	<code>\{ D P \dot E \}</code>	
$S_1 \times S_2$	<code>S_1 \cross S_2</code>	<code>S1 & S2</code>
$S_1 = S_2$	<code>S_1 = S_2</code>	<code>S1 = S2</code>
$S_1 \neq S_2$	<code>S_1 \neq S_2</code>	<code>S1 /= S2</code>
$x \in S$	<code>x \in S</code>	<code>x in S</code>
oz only ▶	<code>x \mem S</code>	
$x \notin S$	<code>x \notin S</code>	<code>x notin S</code>
oz only ▶	<code>x \nem S</code>	
\emptyset	<code>\empty</code>	<code>{ }</code>
$S_1 \subset S_2$	<code>S_1 \subset S_2</code>	<code>S1 subset S2</code>
oz only ▶	<code>S_1 \psubs S_2</code>	
$S_1 \subseteq S_2$	<code>S_1 \subseteq S_2</code>	<code>S1 subseteq S2</code>
oz only ▶	<code>S_1 \subs S_2</code>	
$\mathbb{P} S$	<code>\power S</code>	<code>P S</code>
oz only ▶	<code>\pset S</code>	
$\mathbb{P}_1 S$	<code>\power_1 S</code>	<code>P1 S</code>
oz only ▶	<code>\psetone S</code>	
$\mathbb{F} S$	<code>\finset S</code>	<code>F S</code>
oz only ▶	<code>\fset S</code>	
$\mathbb{F}_1 S$	<code>\finset_1 S</code>	<code>F1 S</code>
oz only ▶	<code>\fsetone S</code>	
$S_1 \cup S_2$	<code>S_1 \cup S_2</code>	<code>S1 setunion S2</code>
oz only ▶	<code>S_1 \uni S_2</code>	<code>S1 S2</code>
$S_1 \cap S_2$	<code>S_1 \cap S_2</code>	<code>S1 setint S2</code>
oz only ▶	<code>S_1 \int S_2</code>	<code>S1 && S2</code>
$S_1 \setminus S_2$	<code>S_1 \setminus S_2</code>	<code>S1 setminus S2</code>
$\bigcup SS$	<code>\bigcup SS</code>	<code>S1 \ S2 Union SS</code>

$\bigcap SS$	<code>\bigcap SS</code>	Intersection SS
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A.2.6 Ordered Pairs

	\LaTeX	ZSL
$x \mapsto y$	<code>x \mapsto y</code>	<code>x mapsto y</code>
<i>first P</i>	<code>first P</code>	<code>first P</code>
<i>second P</i>	<code>second P</code>	<code>second P</code>

A.2.7 Relations

	\LaTeX	ZSL
$A \leftrightarrow B$	<code>A \rel B</code>	<code>A <-> B</code>
$x \mathrel{R} y$	<code>x \inrel{R} y</code>	<code>x _R_ y</code>
$\text{dom } R$	<code>\dom R</code>	<code>dom R</code>
$\text{ran } R$	<code>\ran R</code>	<code>ran R</code>
$\text{id } S$	<code>\id S</code>	<code>id S</code>
$R_1 \circ R_2$	<code>R_1 \comp R_2</code>	<code>R1 comp R2</code>
$R_1 \circ R_2$	<code>R_1 \fcmp R_2</code>	<code>R1 :> R2</code>
$R_1 \circ R_2$	<code>R_1 \circ R_2</code>	<code>R1 backcomp R2</code>
$R_1 \triangleleft R_2$	<code>R_1 \cmp R_2</code>	<code>R1 <: R2</code>
$R_1 \triangleleft R_2$	<code>R_1 \dres R_2</code>	<code>R1 dres R2</code>
$R_1 \triangleleft R_2$	<code>R_1 \ndres R_2</code>	<code>R1 < R2</code>
$R_1 \triangleleft R_2$	<code>R_1 \dsub R_2</code>	<code>R1 dsub R2</code>
$R_1 \triangleright R_2$	<code>R_1 \rres R_2</code>	<code>R1 <+ R2</code>
$R_1 \triangleright R_2$	<code>R_1 \nrres R_2</code>	<code>R1 rres R2</code>
$R_1 \triangleright R_2$	<code>R_1 \rsub R_2</code>	<code>R1 > R2</code>
$R_1 \oplus R_2$	<code>R_1 \oplus R_2</code>	<code>R1 rsub R2</code>
$R_1 \oplus R_2$	<code>R_1 \oplus R_2</code>	<code>R1 +> R2</code>
$R_1 \oplus R_2$	<code>R_1 \oplus R_2</code>	<code>R1 oplus R2</code>
$R(S)$	<code>R \lim S \ring</code>	<code>R1 += R2</code>
R^{-1}	<code>R \inv</code>	<code>R (S)</code>
R^*	<code>R \star</code>	<code>R~</code>
R^*	<code>R \rtcl</code>	<code>R inversion</code>
R^+	<code>R \plus</code>	<code>R^*</code>
R^+	<code>R \tcl</code>	<code>R rtclosure</code>
R^k	<code>R \bsup k \esup</code>	<code>R^+</code>
R^k		<code>R tclosure</code>
R^k		<code>R^(k)</code>

A.2.8 Functions

	L ^A T _E X	ZSL
$A \twoheadrightarrow B$	<code>A \pfun B</code>	<code>A ++> B</code> <code>A pfun B</code>
$A \rightarrow B$	<code>A \fun B</code>	<code>A --> B</code>
$A \dashrightarrow B$	oz only ▶ <code>A \tfun B</code>	<code>A fun B</code>
$A \twoheadleftarrow B$	<code>A \pinj B</code>	<code>A >> B</code> <code>A pinj B</code>
$A \rightarrowtail B$	<code>A \inj B</code>	<code>A >-> B</code>
$A \dashrightarrowtail B$	oz only ▶ <code>A \tinj B</code>	<code>A inj B</code>
$A \twoheadrightarrowtail B$	<code>A \psurj B</code>	<code>A +>> B</code>
$A \rightarrowtail B$	oz only ▶ <code>A \psur B</code>	<code>A psurj B</code>
$A \dashrightarrowtail B$	<code>A \surj B</code>	<code>A ->> B</code>
$A \twoheadrightarrowtail B$	oz only ▶ <code>A \tsur B</code>	<code>A surj B</code>
$A \rightarrowtail B$	<code>A \bij B</code>	<code>A >->> B</code> <code>A bij B</code>
$A \twoheadrightarrow B$	<code>A \ffun B</code>	<code>A ++> B</code> <code>A ffun B</code>
$A \twoheadrightarrowtail B$	<code>A \finj B</code>	<code>A >+>> B</code> <code>A finj B</code>

A.2.9 Numbers

	L ^A T _E X	ZSL
\mathbb{N}	<code>\nat</code>	<code>N</code> <code>Nat</code>
\mathbb{N}_1	<code>\nat_1</code>	<code>N1</code>
\mathbb{Z}	oz only ▶ <code>\natone</code> <code>\num</code>	<code>Nat1</code> <code>Z</code>
$n \dots m$	oz only ▶ <code>\integer</code> <code>n \upto m</code>	<code>Int</code> <code>n upto m</code> <code>n .. m</code>
$x + y$	<code>x + y</code>	<code>x + y</code>
$x - y$	<code>x - y</code>	<code>x - y</code>
$x * y$	<code>x * y</code>	<code>x * y</code>
$x = y$	<code>x = y</code>	<code>x = y</code>
$x \neq y$	<code>x \neq y</code>	<code>x /= y</code>
$x \operatorname{div} y$	<code>x \div y</code>	<code>x div y</code>
$x \operatorname{mod} y$	<code>x \mod y</code>	<code>x mod y</code>
$x < y$	<code>x < y</code>	<code>x < y</code>
$x \leq y$	<code>x \leq y</code>	<code>x <= y</code>
$x > y$	<code>x > y</code>	<code>x > y</code>
$x \geq y$	<code>x \geq y</code>	<code>x >= y</code>
$\operatorname{succ} x$	<code>succ x</code>	<code>succ x</code>

$\#S$	$\backslash\# S$	$\# S$
$\min S$	$\min\sim S$	$\min S$
$\max S$	$\max\sim S$	$\max S$

A.2.10 Sequences

	L ^A T _E X	ZSL
$\text{seq } X$	$\backslash\text{seq } X$	$\text{seq } X$
$\text{seq}_1 X$	$\backslash\text{seq}_1 X$	$\text{seq}_1 X$
	oz only ▶ $\backslash\text{seqone } X$	
$\text{iseq } X$	$\backslash\text{iseq } X$	$\text{iseq } X$
$\langle s_1, \dots, s_n \rangle$	$\backslash\langle s_1, \dots, s_n \rangle$	$\langle\langle s_1, \dots, s_n \rangle\rangle$
	$\backslash\langle s_1, \dots, s_n \rangle$	
	oz only ▶ $\backslash\text{lseq } s_1, \dots, s_n \backslash\text{rseq}$	
$s \hat{\ } t$	$s \backslash\text{cat } t$	$s \text{ concat } t$
$\text{head } s$	$\text{head}\sim s$	$s \hat{\ } t$
$\text{last } s$	$\text{last}\sim s$	$\text{head } s$
$\text{tail } s$	$\text{tail}\sim s$	$\text{last } s$
$\text{front } s$	$\text{front}\sim s$	$\text{tail } s$
$\text{rev } s$	$\text{rev}\sim s$	$\text{front } s$
$s X$	$s \backslash\text{filter } X$	$\text{rev } s$
	oz only ▶ $s \backslash\text{sres } X$	$s \text{ filter } X$
$X s$	$X \backslash\text{extract } s$	$s - X$
	oz only ▶ $X \backslash\text{ires } s$	$X \text{ extract } s$
\wedge / ss	$\backslash\text{dcat } ss$	$X - s$
$\text{disjoint } ss$	$\backslash\text{disjoint } ss$	\wedge / ss
$ss \text{ partition } S$	$ss \backslash\text{partition } S$	$\text{disjoint } ss$
$s_1 \text{ in } s_2$	$s_1 \backslash\text{subseq } s_2$	$ss \text{ partition } S$
	oz only ▶ $s_1 \backslash\text{inseq } s_2$	$s_1 \text{ subseq } s_2$
$s_1 \text{ prefix } s_2$	$s_1 \backslash\text{prefix } s_2$	
$s_1 \text{ suffix } s_2$	$s_1 \backslash\text{suffix } s_2$	$s_1 \text{ prefix } s_2$
$\text{squash } s$	$\text{squash}\sim s$	$s_1 \text{ suffix } s_2$
		$\text{squash } s$

A.2.11 Bags

	L ^A T _E X	ZSL
$\text{bag } X$	$\backslash\text{bag } X$	$\text{bag } X$
$\llbracket a_1, \dots, a_n \rrbracket$	$\backslash\text{lbag } a_1, \dots, a_n \backslash\text{rbag}$	$\llbracket [a_1, \dots, a_n] \rrbracket$
$x \in B$	$x \backslash\text{inbag } B$	$x \text{ inbag } B$
$\text{count } B$	$\text{count } B$	$\text{count } B$
$B_1 \sqsubseteq B_2$	$B_1 \backslash\text{subbag } B_2$	$B_1 \text{ subbag } B_2$
$B_1 \cup B_2$	$B_1 \backslash\text{bagdiff } B_2$	$B_1 \text{ bagdiff } B_2$

$n \otimes B$	<code>n \bagscale B</code>	<code>B1 -- B2</code>
$B \sharp x$	<code>B \bagcount x</code>	<code>n bagscale B</code>
$B_1 \uplus B_2$	<code>B_1 \uplus B_2</code>	<code>B bagcount x</code>
	oz only ► <code>B_1 \buni B_2</code>	<code>B1 bagunion B2</code>
$items\ s$	<code>items s</code>	<code>B1 ++ B2</code>
		<code>items s</code>

A.2.12 Binding

	\LaTeX	ZSL
θS	<code>\theta S</code>	<code>theta S</code>

A.2.13 Selection

	\LaTeX	ZSL
$S.x$	<code>S.x</code>	<code>S.x</code>

A.2.14 Operators

	\LaTeX	ZSL
$PreSym_$	<code>PreSym _</code>	<code>PreSym _</code>
$_InSym_$	<code>_ InSym _</code>	<code>_ InSym _</code>
$_PostSym$	<code>_ PostSym</code>	<code>_ PostSym</code>
$-(-)$	<code>_ \limg _ \ring</code>	<code>_ (-)</code>